

# Top Google Questions

[Download PDF](#)

- [Top Google Questions](#)
  - [1. Two Sum](#)
  - [3. Longest Substring Without Repeating Characters](#)
  - [4. Median of Two Sorted Arrays](#)
  - [5. Longest Palindromic Substring](#)
  - [6. ZigZag Conversion](#)
  - [8. String to Integer \(atoi\)](#)
  - [11. Container With Most Water](#)
  - [14. Longest Common Prefix](#)
  - [19. Remove Nth Node From End of List](#)
  - [23. Merge k Sorted Lists](#)
  - [24. Swap Nodes in Pairs](#)
  - [26. Remove Duplicates from Sorted Array](#)
  - [27. Remove Element](#)
  - [33. Search in Rotated Sorted Array](#)
  - [34. Find First and Last Position of Element in Sorted Array](#)
  - [35. Search Insert Position](#)
  - [38. Count and Say](#)
  - [41. First Missing Positive](#)
  - [46. Permutations](#)
  - [50. Pow\(x, n\)](#)
  - [51. N-Queens](#)
  - [53. Maximum Subarray](#)
  - [54. Spiral Matrix](#)
  - [55. Jump Game](#)
  - [59. Spiral Matrix II](#)
  - [60. Permutation Sequence](#)
  - [66. Plus One](#)
  - [67. Add Binary](#)
  - [69. Sqrt\(x\)](#)
  - [70. Climbing Stairs](#)
  - [72. Edit Distance](#)
  - [74. Search a 2D Matrix](#)
  - [75. Sort Colors](#)
  - [77. Combinations](#)

- 78. Subsets
- 79. Word Search
- 80. Remove Duplicates from Sorted Array II
- 81. Search in Rotated Sorted Array II
- 83. Remove Duplicates from Sorted List
- 84. Largest Rectangle in Histogram
- 88. Merge Sorted Array
- 94. Binary Tree Inorder Traversal
- 96. Unique Binary Search Trees
- 101. Symmetric Tree
- 102. Binary Tree Level Order Traversal
- 103. Binary Tree Zigzag Level Order Traversal
- 105. Construct Binary Tree from Preorder and Inorder Traversal
- 109. Convert Sorted List to Binary Search Tree
- 110. Balanced Binary Tree
- 114. Flatten Binary Tree to Linked List
- 118. Pascal's Triangle
- 119. Pascal's Triangle II
- 121. Best Time to Buy and Sell Stock
- 129. Sum Root to Leaf Numbers
- 130. Surrounded Regions
- 136. Single Number
- 137. Single Number II
- 144. Binary Tree Preorder Traversal
- 145. Binary Tree Postorder Traversal
- 146. LRU Cache
- 153. Find Minimum in Rotated Sorted Array
- 162. Find Peak Element
- 163. Missing Ranges
- 169. Majority Element
- 173. Binary Search Tree Iterator
- 174. Dungeon Game
- 198. House Robber
- 200. Number of Islands
- 205. Isomorphic Strings
- 207. Course Schedule
- 208. Implement Trie (Prefix Tree)
- 219. Contains Duplicate II
- 221. Maximal Square
- 222. Count Complete Tree Nodes
- 226. Invert Binary Tree
- 230. Kth Smallest Element in a BST
- 231. Power of Two
- 237. Delete Node in a Linked List
- 242. Valid Anagram
- 246. Strobogrammatic Number

- 249. Group Shifted Strings
- 250. Count Univalued Subtrees
- 251. Flatten 2D Vector
- 252. Meeting Rooms
- 253. Meeting Rooms II
- 257. Binary Tree Paths
- 265. Paint House II
- 267. Palindrome Permutation II
- 268. Missing Number
- 270. Closest Binary Search Tree Value
- 274. H-Index
- 276. Paint Fence
- 278. First Bad Version
- 285. Inorder Successor in BST
- 286. Walls and Gates
- 287. Find the Duplicate Number
- 290. Word Pattern
- 299. Bulls and Cows
- 303. Range Sum Query - Immutable
- 305. Number of Islands II
- 311. Sparse Matrix Multiplication
- 314. Binary Tree Vertical Order Traversal
- 322. Coin Change
- 323. Number of Connected Components in an Undirected Graph
- 326. Power of Three
- 328. Odd Even Linked List
- 344. Reverse String
- 345. Reverse Vowels of a String
- 347. Top K Frequent Elements
- 348. Design Tic-Tac-Toe
- 367. Valid Perfect Square
- 368. Largest Divisible Subset
- 374. Guess Number Higher or Lower
- 379. Design Phone Directory
- 380. Insert Delete GetRandom O(1)
- 384. Shuffle an Array
- 387. First Unique Character in a String
- 389. Find the Difference
- 392. Is Subsequence
- 402. Remove K Digits
- 404. Sum of Left Leaves
- 406. Queue Reconstruction by Height
- 409. Longest Palindrome
- 415. Add Strings
- 429. N-ary Tree Level Order Traversal
- 438. Find All Anagrams in a String

- 447. Number of Boomerangs
- 451. Sort Characters By Frequency
- 468. Validate IP Address
- 470. Implement Rand10() Using Rand7()
- 475. Heaters
- 482. License Key Formatting
- 498. Diagonal Traverse
- 510. Inorder Successor in BST II
- 518. Coin Change 2
- 525. Contiguous Array
- 528. Random Pick with Weight
- 535. Encode and Decode TinyURL
- 540. Single Element in a Sorted Array
- 542. 01 Matrix
- 551. Student Attendance Record I
- 567. Permutation in String
- 609. Find Duplicate File in System
- 616. Add Bold Tag in String
- 658. Find K Closest Elements
- 669. Trim a Binary Search Tree
- 684. Redundant Connection
- 679. 24 Game
- 693. Binary Number with Alternating Bits
- 694. Number of Distinct Islands
- 700. Search in a Binary Search Tree
- 702. Search in a Sorted Array of Unknown Size
- 705. Design HashSet
- 706. Design HashMap
- 711. Number of Distinct Islands II
- 717. 1-bit and 2-bit Characters
- 720. Longest Word in Dictionary
- 721. Accounts Merge
- 733. Flood Fill
- 734. Sentence Similarity
- 737. Sentence Similarity II
- 748. Shortest Completing Word
- 758. Bold Words in String
- 773. Sliding Puzzle
- 787. Cheapest Flights Within K Stops
- 819. Most Common Word
- 844. Backspace String Compare
- 868. Binary Gap
- 917. Reverse Only Letters
- 973. K Closest Points to Origin
- 986. Interval List Intersections
- 988. Smallest String Starting From Leaf

- 997. Find the Town Judge
- 1029. Two City Scheduling
- 1057. Campus Bikes
- 1065. Index Pairs of a String
- 1066. Campus Bikes II
- 1143. Longest Common Subsequence
- 1170. Compare Strings by Frequency of the Smallest Character
- 1277. Count Square Submatrices with All Ones
- 1287. Element Appearing More Than 25\% In Sorted Array
- 1314. Matrix Block Sum
- 1331. Rank Transform of an Array
- 1422. Maximum Score After Splitting a String
- 1424. Diagonal Traverse II
- 1438. Longest Continuous Subarray With Absolute Diff Less Than or Equal to Limit
- 1441. Build an Array With Stack Operations
- 1447. Simplified Fractions
- 1452. People Whose List of Favorite Companies Is Not a Subset of Another List
- 1461. Check If a String Contains All Binary Codes of Size K
- 1463. Cherry Pickup II
- 1471. The k Strongest Values in an Array
- Check If a String Is a Valid Sequence from Root to Leaves Path in a Binary Tree
- 1477. Find Two Non-overlapping Sub-arrays Each With Target Sum
- 1480. Running Sum of 1d Array
- 1482. Minimum Number of Days to Make m Bouquets
- 1483. Kth Ancestor of a Tree Node
- 1488. Avoid Flood in The City

## 1. Two Sum

### *Description*

Given an array of integers, return indices of the two numbers such that they add up to a specific target.

You may assume that each input would have exactly one solution, and you may not use the same element twice.

Example:

Given `nums = [2, 7, 11, 15]`, `target = 9`,

Because `nums[0] + nums[1] = 2 + 7 = 9`,  
return `[0, 1]`.

*Solution*

01/02/2020:

```
class Solution {
public:
    vector<int> twoSum(vector<int>& nums, int target) {
        unordered_map<int, int> m;
        for (int i = 0; i < nums.size(); ++i) {
            if (m.find(target - nums[i]) == m.end()) {
                m[nums[i]] = i;
            } else {
                return {m[target - nums[i]], i};
            }
        }
        return {-1, -1};
    }
};
```

### 3. Longest Substring Without Repeating Characters

*Description*

Given a string, find the length of the longest substring without repeating characters.

Example 1:

Input: "abcabcbb"

Output: 3

Explanation: The answer is "abc", with the length of 3.

Example 2:

Input: "bbbbbb"

Output: 1

Explanation: The answer is "b", with the length of 1.

Example 3:

Input: "pwwkew"

Output: 3

Explanation: The answer is "wke", with the length of 3.

Note that the answer must be a substring, "pwke" is a subsequence and not a substring.

*Solution*

05/21/2020:

```
class Solution {
public:
    int lengthOfLongestSubstring(string s) {
        unordered_map<char, int> seen;
        int ret = 0, slow = 0, n = s.size();
        for (int fast = 0; fast < n; ++fast) {
            if (seen.count(s[fast]) != 0) slow = max(slow, seen[s[fast]] + 1);
            seen[s[fast]] = fast;
            ret = max(ret, fast - slow + 1);
        }
        return ret;
    }
};
```

## 4. Median of Two Sorted Arrays

### *Description*

There are two sorted arrays `nums1` and `nums2` of size `m` and `n` respectively.

Find the median of the two sorted arrays. The overall run time complexity should be  $O(\log(m+n))$ .

You may assume `nums1` and `nums2` cannot be both empty.

Example 1:

`nums1 = [1, 3]`

`nums2 = [2]`

The median is 2.0

Example 2:

`nums1 = [1, 2]`

`nums2 = [3, 4]`

The median is  $(2 + 3)/2 = 2.5$

### *Solution*

01/02/2020:

```
class Solution {
```

```

public:
    double findMedianSortedArrays(vector<int>& nums1, vector<int>& nums2) {
        int n = nums1.size() + nums2.size();
        vector<int> nums(n, 0);
        auto it1 = nums1.begin();
        auto it2 = nums2.begin();
        auto it = nums.begin();
        for (; it != nums.end(); ++it) {
            *it = it2 == nums2.end() || (it1 != nums1.end() && *it1 < *it2) ? *it1++ :
*it2++;
        }
        if (n % 2 == 0) {
            return ((double) nums[n/2] + nums[n/2-1]) / 2;
        } else {
            return (double) nums[n/2];
        }
    }
};

```

## 5. Longest Palindromic Substring

### Description

Given a string *s*, find the longest palindromic substring in *s*. You may assume that the maximum length of *s* is 1000.

Example 1:

Input: "babad"

Output: "bab"

Note: "aba" is also a valid answer.

Example 2:

Input: "cbbd"

Output: "bb"

### Solution

01/13/2020 (Dynamic Programming):

```

class Solution {
public:
    string longestPalindrome(string s) {
        int n = s.size(), start, max_len = 0;
        if (n == 0) return "";
    }
};

```



```

vector<vector<bool>> dp(n, vector<bool>(n, false));
for (int i = 0; i < n; ++i) dp[i][i] = true;
for (int i = 0; i < n - 1; ++i) dp[i][i + 1] = s[i] == s[i + 1];
for (int i = n - 3; i >= 0; --i) {
    for (int j = i + 2; j < n; ++j) {
        dp[i][j] = dp[i + 1][j - 1] && s[i] == s[j];
    }
}
for (int i = 0; i < n; ++i) {
    for (int j = i; j < n; ++j) {
        if (dp[i][j] && j - i + 1 > max_len) {
            max_len = j - i + 1;
            start = i;
        }
    }
}
return s.substr(start, max_len);
}
};

```

01/13/2020 (Expand Around Center):

```

class Solution {
public:
    string longestPalindrome(string s) {
        int n = s.size(), start = 0, max_len = n > 0 ? 1 : 0;
        for(int i = 0; i < n; ++i) {
            for (int l = i - 1, r = i; l >= 0 && r < n && s[l] == s[r]; --l, ++r) {
                if (r - l + 1 > max_len) {
                    max_len = r - l + 1;
                    start = l;
                }
            }
            for (int l = i - 1, r = i + 1; l >= 0 && r < n && s[l] == s[r]; --l, ++r)
            {
                if (r - l + 1 > max_len) {
                    max_len = r - l + 1;
                    start = l;
                }
            }
        }
        return max_len == 0 ? "" : s.substr(start, max_len);
    }
};

```

## 6. ZigZag Conversion

### Description

The string "PAYPALISHIRING" is written in a zigzag pattern on a given number of rows like this: (you may want to display this pattern in a fixed font for better legibility)

```
P   A   H   N
A P L S I I G
Y   I   R
```

And then read line by line: "PAHNAPLSIIGYIR"

Write the code that will take a string and make this conversion given a number of rows:

```
string convert(string s, int numRows);
```

Example 1:

Input: s = "PAYPALISHIRING", numRows = 3

Output: "PAHNAPLSIIGYIR"

Example 2:

Input: s = "PAYPALISHIRING", numRows = 4

Output: "PINALSIGYAHRPI"

Explanation:

```
P       I       N
A   L S   I G
Y A   H R
P       I
```

### Solution

05/24/2020:

```
class Solution {
public:
    string convert(string s, int numRows) {
        if (s.empty()) return "";
        if (numRows <= 1) return s;
        int k = 0, increment = 1;
        vector<string> strs(numRows);
        for (int i = 0; i < (int)s.size(); ++i) {
            strs[k].push_back(s[i]);
            if (k == numRows - 1) {
                increment = -1;
            }
        }
    }
};
```

```

    } else if (k == 0) {
        increment = 1;
    }
    k += increment;
}
string ret;
for (auto& s : strs) ret += s;
return ret;
}
};

```

## 8. String to Integer (atoi)

### *Description*

Implement atoi which converts a string to an integer.

The function first discards as many whitespace characters as necessary until the first non-whitespace character is found. Then, starting from this character, takes an optional initial plus or minus sign followed by as many numerical digits as possible, and interprets them as a numerical value.

The string can contain additional characters after those that form the integral number, which are ignored and have no effect on the behavior of this function.

If the first sequence of non-whitespace characters in str is not a valid integral number, or if no such sequence exists because either str is empty or it contains only whitespace characters, no conversion is performed.

If no valid conversion could be performed, a zero value is returned.

Note:

Only the space character ' ' is considered as whitespace character.

Assume we are dealing with an environment which could only store integers within the 32-bit signed integer range:  $[-2^{31}, 2^{31} - 1]$ . If the numerical value is out of the range of representable values, `INT_MAX` ( $2^{31} - 1$ ) or `INT_MIN` ( $-2^{31}$ ) is returned.

Example 1:

Input: "42"

Output: 42

Example 2:

Input: " -42"

Output: -42

Explanation: The first non-whitespace character is '-', which is the minus sign.  
Then take as many numerical digits as possible, which gets 42.

Example 3:

Input: "4193 with words"

Output: 4193

Explanation: Conversion stops at digit '3' as the next character is not a numerical digit.

Example 4:

Input: "words and 987"

Output: 0

Explanation: The first non-whitespace character is 'w', which is not a numerical digit or a +/- sign. Therefore no valid conversion could be performed.

Example 5:

Input: "-91283472332"

Output: -2147483648

Explanation: The number "-91283472332" is out of the range of a 32-bit signed integer.

Therefore INT\_MIN (-231) is returned.

*Solution*

05/24/2020:

```
class Solution {
public:
    int myAtoi(string str) {
        long long ret = atol(str.c_str());
        ret = ret > INT_MAX ? INT_MAX : ret;
        ret = ret < INT_MIN ? INT_MIN : ret;
        return ret;
    }
};
```

```
class Solution {
public:
    int myAtoi(string str) {
        if (str.empty()) return 0;
        long long ret = 0;
        istringstream iss(str);
        iss >> ret;
        ret = ret < INT_MIN ? INT_MIN : ret;
        ret = ret > INT_MAX ? INT_MAX : ret;
        return ret;
    }
};
```

## 11. Container With Most Water

### *Description*

Given  $n$  non-negative integers  $a_1, a_2, \dots, a_n$ , where each represents a point at coordinate  $(i, a_i)$ .  $n$  vertical lines are drawn such that the two endpoints of line  $i$  is at  $(i, a_i)$  and  $(i, 0)$ . Find two lines, which together with x-axis forms a container, such that the container contains the most water.

Note: You may not slant the container and  $n$  is at least 2.

The above vertical lines are represented by array  $[1,8,6,2,5,4,8,3,7]$ . In this case, the max area of water (blue section) the container can contain is 49.

Example:

Input:  $[1,8,6,2,5,4,8,3,7]$

Output: 49

### *Solution*

05/24/2020:

```

class Solution {
public:
    int maxArea(vector<int>& height) {
        int left = 0, right = height.size() - 1, ret = 0;
        while (left < right) {
            ret = max(ret, min(height[left], height[right]) * (right - left));
            height[left] < height[right] ? left += 1 : right -= 1;
        }
        return ret;
    }
};

```

## 14. Longest Common Prefix

### Description

Write a function to find the longest common prefix string amongst an array of strings.

If there is no common prefix, return an empty string "".

Example 1:

Input: ["flower","flow","flight"]

Output: "fl"

Example 2:

Input: ["dog","racecar","car"]

Output: ""

Explanation: There is no common prefix among the input strings.

Note:

All given inputs are in lowercase letters a-z.

### Solution

05/24/2020:

```

class Solution {
public:
    string longestCommonPrefix(vector<string>& strs) {
        if (strs.empty() || strs[0].empty()) return "";
        int minLength = strs[0].size();
        for (auto& s : strs) minLength = min(minLength, (int)s.size());
        string ret;
    }
};

```

```

for (int i = 0; i < minLength; ++i) {
    char cur = strs[0][i];
    for (int j = 1; j < (int)strs.size(); ++j) {
        if (strs[j][i] != cur) {
            return ret;
        }
    }
    ret += cur;
}
return ret;
};

```

```

class Solution {
public:
    string longestCommonPrefix(vector<string>& strs) {
        if (strs.empty() || strs[0].empty()) return "";
        sort(strs.begin(), strs.end(), [](const string& s1, const string& s2) {
            if (s1.size() == s2.size()) return s1 < s2;
            return s1.size() < s2.size();
        });
        string ret;
        for (int k = (int)strs[0].size(); k >= 0; --k) {
            bool isCommonPrefix = true;
            string prefix = strs[0].substr(0, k);
            for (int i = 1; i < (int)strs.size(); ++i) {
                if (prefix != strs[i].substr(0, k)) {
                    isCommonPrefix = false;
                    break;
                }
            }
            if (isCommonPrefix) {
                ret = prefix;
                break;
            }
        }
        return ret;
    }
};

```

```

struct Node {
    bool isWord;
    vector<Node*> children;
    Node() { isWord = false; children.resize(26, nullptr); }
    ~Node() { for (auto& c : children) delete c; }
};

```

```

class Trie {
private:
    Node* root;

    Node* find(const string& word) {
        Node* cur = root;
        for (auto i = 0; i < (int)word.size(); ++i) {
            if (cur->children[word[i] - 'a'] == nullptr) {
                break;
            }
            cur = cur->children[word[i] - 'a'];
        }
        return cur;
    }

public:
    Trie() { root = new Node(); }

    void insert(const string& word) {
        Node* cur = root;
        for (auto i = 0; i < (int)word.size(); ++i) {
            if (cur->children[word[i] - 'a'] == nullptr) {
                cur->children[word[i] - 'a'] = new Node();
            }
            cur = cur->children[word[i] - 'a'];
        }
        cur->isWord = true;
    }

    bool contains(const string& word) {
        Node* cur = find(word);
        return cur && cur->isWord;
    }

    bool startsWith(const string& prefix) {
        Node* cur = find(prefix);
        return cur;
    }

    string commonPrefix() {
        string ret;
        Node* cur = root;
        while (true) {
            int cnt = 0;
            char ch = 'a';
            bool isUnique = false;
            for (int i = 0; i < 26; ++i) {
                if (cur->children[i] != nullptr) {
                    isUnique = true;
                }
            }
            if (!isUnique) break;
            ret += ch;
            cur = cur->children[ch - 'a'];
        }
        return ret;
    }
};

```



```

        if (++cnt > 1) {
            isUnique = false;
            break;
        }
        ch = i + 'a';
    }
}
if (isUnique) {
    ret += ch;
    cur = cur->children[ch - 'a'];
} else {
    break;
}
}
return ret;
}
};

class Solution {
public:
    string longestCommonPrefix(vector<string>& strs) {
        Trie t;
        int minLength = INT_MAX;
        for (auto& s : strs) {
            if (s.empty()) return "";
            t.insert(s);
            minLength = min(minLength, (int)s.size());
        }
        return t.commonPrefix().substr(0, minLength);
    }
};

```

## 19. Remove Nth Node From End of List

### *Description*

Given a linked list, remove the n-th node from the end of list and return its head.

Example:

Given linked list: 1->2->3->4->5, and n = 2.

After removing the second node from the end, the linked list becomes 1->2->3->5.

Note:

Given n will always be valid.

Follow up:

Could you do this in one pass?

*Solution*

05/23/2020:

```
/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
public:
    ListNode* removeNthFromEnd(ListNode* head, int n) {
        ListNode *pre = new ListNode(0, head), *slow = pre, *fast = pre;
        while (fast->next != nullptr && n-- > 0) fast = fast->next;
        while (fast->next != nullptr) {
            fast = fast->next;
            slow = slow->next;
        }
        slow->next = slow->next->next;
        return pre->next;
    }
};
```

## 23. Merge k Sorted Lists

*Description*

Merge k sorted linked lists and return it as one sorted list. Analyze and describe its complexity.

Example:

Input:

```
[
  1->4->5,
  1->3->4,
  2->6
]
```

Output: 1->1->2->3->4->4->5->6

*Solution*

01/02/2020:

```
/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode(int x) : val(x), next(NULL) {}
 * };
 */
class Solution {
public:
    ListNode* mergeKLists(vector<ListNode*>& lists) {
        vector<int> v;
        for (auto& l : lists) {
            vector<int> tmp = ListNode2vector(l);
            v.insert(v.begin(), tmp.begin(), tmp.end());
        }
        sort(v.begin(), v.end());
        return vector2ListNode(v);
    }

    vector<int> ListNode2vector(ListNode* list) {
        vector<int> v;
        for (; list != nullptr; list = list->next)
            v.push_back(list->val);
        return v;
    }

    ListNode* vector2ListNode(vector<int>& v) {
        ListNode *pre = new ListNode(0), *cur = pre;
        for (auto& n : v) {
            cur->next = new ListNode(n);
        }
    }
};
```

```

        cur = cur->next;
    }
    return pre->next;
}
};

```

Mimicking merge\_sort 01/02/2020:

```

/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode(int x) : val(x), next(NULL) {}
 * };
 */
class Solution {
public:
    ListNode* mergeKLists(vector<ListNode*>& lists) {
        int n = lists.size();
        if (n == 0) return nullptr;
        if (n == 1) return lists[0];
        vector<ListNode*> lists1(lists.begin(), lists.begin() + n / 2);
        vector<ListNode*> lists2(lists.begin() + n / 2, lists.end());
        ListNode* l1 = mergeKLists(lists1);
        ListNode* l2 = mergeKLists(lists2);
        if (l1 == nullptr) return l2;
        ListNode* ret = l1;
        while (l2 != nullptr) {
            if (l1->val > l2->val) swap(l1->val, l2->val);
            while(l1->next && l1->next->val < l2->val) l1 = l1->next;
            ListNode* tmp2 = l2->next;
            l2->next = l1->next;
            l1->next = l2;
            l2 = tmp2;
        }
        return ret;
    }
};

```

## 24. Swap Nodes in Pairs

*Description*

Given a linked list, swap every two adjacent nodes and return its head.

You may not modify the values in the list's nodes, only nodes itself may be changed.

Example:

Given 1->2->3->4, you should return the list as 2->1->4->3.

*Solution*

05/22/2020:

```
/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
public:
    ListNode* swapPairs(ListNode* head) {
        // pre -> A -> B -> C -> D -> null
        // pre -> B (pre->next = pre->next->next)
        // A -> C (A->next = A->next->next)
        // B -> A (B->next = A):
        // (B -> A -> C -> D -> null)
        ListNode* pre = new ListNode(0, head);
        ListNode* cur = pre;
        while (cur->next && cur->next->next) {
            ListNode* tmp = cur->next;
            cur->next = cur->next->next;
            tmp->next = tmp->next->next;
            cur->next->next = tmp;
            cur = tmp;
        }
        return pre->next;
    }
};
```

## 26. Remove Duplicates from Sorted Array

### Description

Given a sorted array `nums`, remove the duplicates in-place such that each element appear only once and return the new length.

Do not allocate extra space for another array, you must do this by modifying the input array in-place with  $O(1)$  extra memory.

Example 1:

Given `nums = [1,1,2]`,

Your function should return `length = 2`, with the first two elements of `nums` being 1 and 2 respectively.

It doesn't matter what you leave beyond the returned length.

Example 2:

Given `nums = [0,0,1,1,1,2,2,3,3,4]`,

Your function should return `length = 5`, with the first five elements of `nums` being modified to 0, 1, 2, 3, and 4 respectively.

It doesn't matter what values are set beyond the returned length.

Clarification:

Confused why the returned value is an integer but your answer is an array?

Note that the input array is passed in by reference, which means modification to the input array will be known to the caller as well.

Internally you can think of this:

```
// nums is passed in by reference. (i.e., without making a copy)
int len = removeDuplicates(nums);

// any modification to nums in your function would be known by the caller.
// using the length returned by your function, it prints the first len elements.
for (int i = 0; i < len; i++) {
    print(nums[i]);
}
```

### Solution

05/27/2020:

```

class Solution {
public:
    int removeDuplicates(vector<int>& nums) {
        if (nums.size() <= 1) return nums.size();
        int slow = 0, fast = 1, n = nums.size();
        for (; fast < n; ++fast) {
            while (nums[fast] == nums[fast - 1]) if (++fast >= n) break;
            if (fast < n) nums[++slow] = nums[fast];
        }
        return slow + 1;
    }
};

```

## 27. Remove Element

### *Description*

Given an array `nums` and a value `val`, remove all instances of that value in-place and return the new length.

Do not allocate extra space for another array, you must do this by modifying the input array in-place with  $O(1)$  extra memory.

The order of elements can be changed. It doesn't matter what you leave beyond the new length.

Example 1:

Given `nums = [3,2,2,3]`, `val = 3`,

Your function should return `length = 2`, with the first two elements of `nums` being `2`.

It doesn't matter what you leave beyond the returned length.

Example 2:

Given `nums = [0,1,2,2,3,0,4,2]`, `val = 2`,

Your function should return `length = 5`, with the first five elements of `nums` containing `0, 1, 3, 0, and 4`.

Note that the order of those five elements can be arbitrary.

It doesn't matter what values are set beyond the returned length.

Clarification:

Confused why the returned value is an integer but your answer is an array?

Note that the input array is passed in by reference, which means modification to the input array will be known to the caller as well.

Internally you can think of this:

```
// nums is passed in by reference. (i.e., without making a copy)
int len = removeElement(nums, val);

// any modification to nums in your function would be known by the caller.
// using the length returned by your function, it prints the first len elements.
for (int i = 0; i < len; i++) {
    print(nums[i]);
}
```

*Solution*

05/23/2020:

```
class Solution {
public:
    int removeElement(vector<int>& nums, int val) {
        int right = nums.size() - 1, n = nums.size(), cnt = n;
        for (int i = 0; i < n && i <= right; ++i) {
            while (right >= 0 && nums[right] == val) {
                --right;
                --cnt;
            }
            if (nums[i] == val && i <= right) {
                swap(nums[i], nums[right--]);
                --cnt;
            }
        }
        return cnt;
    }
};
```

## 33. Search in Rotated Sorted Array

*Description*

Suppose an array sorted in ascending order is rotated at some pivot unknown to you beforehand.



(i.e., [0,1,2,4,5,6,7] might become [4,5,6,7,0,1,2]).

You are given a target value to search. If found in the array return its index, otherwise return -1.

You may assume no duplicate exists in the array.

Your algorithm's runtime complexity must be in the order of  $O(\log n)$ .

Example 1:

Input: nums = [4,5,6,7,0,1,2], target = 0

Output: 4

Example 2:

Input: nums = [4,5,6,7,0,1,2], target = 3

Output: -1

*Solution*

04/28/2020:

```
class Solution {
public:
    int search(vector<int>& nums, int target) {
        int lo = 0, hi = nums.size() - 1;
        while (lo <= hi) {
            int mid = lo + (hi - lo) / 2;
            if (nums[mid] == target) {
                return mid;
            } else {
                if (nums[mid] < nums[lo]) {
                    if (target > nums[lo] && target > nums[hi]) {
                        hi = mid - 1;
                    } else {
                        lo = mid + 1;
                    }
                } else {
                    if (target > nums[lo] && target > nums[hi]) {
                        hi = mid - 1;
                    } else {
                        lo = mid + 1;
                    }
                }
            }
        }
        return -1;
    }
};
```

```
}  
};
```

## 34. Find First and Last Position of Element in Sorted Array

### Description

Given an array of integers `nums` sorted in ascending order, find the starting and ending position of a given target value.

Your algorithm's runtime complexity must be in the order of  $O(\log n)$ .

If the target is not found in the array, return `[-1, -1]`.

Example 1:

Input: `nums = [5,7,7,8,8,10]`, `target = 8`

Output: `[3,4]`

Example 2:

Input: `nums = [5,7,7,8,8,10]`, `target = 6`

Output: `[-1,-1]`

### Solution

04/28/2020:

```
class Solution {  
public:  
    vector<int> searchRange(vector<int>& nums, int target) {  
        if (nums.empty()) return {-1, -1};  
        int n = nums.size();  
        int lower = -1, upper = -1;  
  
        // to find lower_bound  
        int lo = 0, hi = n - 1;  
        while (lo < hi) {  
            int mid = lo + (hi - lo) / 2;  
            if (nums[mid] >= target) {  
                hi = mid;  
            } else {  
                lo = mid + 1;  
            }  
        }  
        if (nums[lo] < target) ++lo;  
    }  
};
```

```

lower = lo;
if (lower > hi) return {-1, -1};

// to find the upper bound
lo = 0, hi = n - 1;
while (lo < hi) {
    int mid = lo + (hi - lo) / 2;
    if (nums[mid] > target) {
        hi = mid;
    } else {
        lo = mid + 1;
    }
}
if (nums[hi] > target) --hi;
upper = hi;
if (lower > upper) return {-1, -1};
return {lower, upper};
}
};

```

```

class Solution {
public:
    vector<int> searchRange(vector<int>& nums, int target) {
        if (binary_search(nums.begin(), nums.end(), target)) {
            auto p = equal_range(nums.begin(), nums.end(), target);
            return {int(p.first - nums.begin()), int(p.second - nums.begin() - 1)};
        } else {
            return {-1, -1};
        }
    }
};

```

## 35. Search Insert Position

### Description

Given a sorted array and a target value, return the index if the target is found. If not, return the index where it would be if it were inserted in order.

You may assume no duplicates in the array.

Example 1:

Input: [1,3,5,6], 5

Output: 2

Example 2:

Input: [1,3,5,6], 2

Output: 1

Example 3:

Input: [1,3,5,6], 7

Output: 4

Example 4:

Input: [1,3,5,6], 0

Output: 0

*Solution*

04/23/2020:

```
class Solution {
public:
    int searchInsert(vector<int>& nums, int target) {
        int lo = 0, hi = nums.size() - 1;
        while (lo <= hi) {
            int mid = lo + (hi - lo) / 2;
            if (nums[mid] == target) {
                return mid;
            } else if (nums[mid] > target) {
                hi = mid - 1;
            } else {
                lo = mid + 1;
            }
        }
        return lo;
    }
};
```

## 38. Count and Say

*Description*

The count-and-say sequence is the sequence of integers with the first five terms as following:

1. 1
2. 11
3. 21

4. 1211
  5. 111221
- 1 is read off as "one 1" or 11.  
11 is read off as "two 1s" or 21.  
21 is read off as "one 2, then one 1" or 1211.

Given an integer  $n$  where  $1 \leq n \leq 30$ , generate the  $n$ th term of the count-and-say sequence. You can do so recursively, in other words from the previous member read off the digits, counting the number of digits in groups of the same digit.

Note: Each term of the sequence of integers will be represented as a string.

Example 1:

Input: 1  
Output: "1"  
Explanation: This is the base case.  
Example 2:

Input: 4  
Output: "1211"  
Explanation: For  $n = 3$  the term was "21" in which we have two groups "2" and "1", "2" can be read as "12" which means frequency = 1 and value = 2, the same way "1" is read as "11", so the answer is the concatenation of "12" and "11" which is "1211".

*Solution*

05/24/2020:

```
class Solution {
public:
    string countAndSay(int n) {
        if (n == 1) return "1";
        string s = countAndSay(n - 1);
        string ret;
        int cnt = 1;
        char cur = s[0];
        for (int i = 1; i < (int)s.size(); ++i) {
            if (s[i] == cur) {
                ++cnt;
            } else {
                ret += to_string(cnt) + cur;
                cur = s[i];
                cnt = 1;
            }
        }
    }
};
```

```
    }
    if (cnt > 0) {
        ret += to_string(cnt) + cur;
    }
    return ret;
}
};
```

## 41. First Missing Positive

### Description

Given an unsorted integer array, find the smallest missing positive integer.

Example 1:

Input: [1,2,0]

Output: 3

Example 2:

Input: [3,4,-1,1]

Output: 2

Example 3:

Input: [7,8,9,11,12]

Output: 1

Note:

Your algorithm should run in  $O(n)$  time and uses constant extra space.

### Solution

05/24/2020:

```
class Solution {
public:
    int firstMissingPositive(vector<int>& nums) {
        unordered_set<int> seen;
        for (auto& i : nums) {
            if (i > 0) {
                seen.insert(i);
            }
        }
        int i = 1, n = nums.size();
        while (i <= n) {
```

```

        if (seen.count(i) == 0) return i;
        ++i;
    }
    return i;
}
};

```

## 46. Permutations

### *Description*

Given a collection of distinct integers, return all possible permutations.

Example:

Input: [1,2,3]

Output:

```

[
  [1,2,3],
  [1,3,2],
  [2,1,3],
  [2,3,1],
  [3,1,2],
  [3,2,1]
]

```

### *Solution*

05/26/2020 (Using next\_permutation):

```

class Solution {
public:
    vector<vector<int>> permute(vector<int>& nums) {
        sort(nums.begin(), nums.end());
        vector<vector<int>> ret{nums};
        while (next_permutation(nums.begin(), nums.end())) ret.push_back(nums);
        return ret;
    }
};

```

05/26/2020 (Backtracking):

```

class Solution {
public:
    vector<vector<int>> permute(vector<int>& nums) {

```

```

    int n = nums.size(), numOfChosen = 0;
    vector<bool> chosen(n, false);
    vector<vector<int>> ret;
    vector<int> permutation;
    backtrack(nums, chosen, numOfChosen, permutation, ret);
    return ret;
}

void backtrack(vector<int>& nums, vector<bool>& chosen, int numOfChosen,
vector<int>& permutation, vector<vector<int>>& ret) {
    if (numOfChosen == (int)nums.size()) {
        ret.push_back(permutation);
        return;
    }
    for (int i = 0; i < (int)nums.size(); ++i) {
        if (chosen[i] == true) continue;
        chosen[i] = true;
        permutation.push_back(nums[i]);
        backtrack(nums, chosen, numOfChosen + 1, permutation, ret);
        chosen[i] = false;
        permutation.pop_back();
    }
}
};

```

05/05/2020:

```

class Solution {
public:
    vector<vector<int>> permute(vector<int>& nums) {
        int n = nums.size();
        if (n <= 1) return {nums};
        vector<vector<int>> ret;
        for (int i = 0; i < n; ++i) {
            int cur = nums[i];
            swap(nums[i], nums[n - 1]);
            nums.pop_back();
            vector<vector<int>> sub = permute(nums);
            for (auto& s : sub) {
                s.push_back(cur);
                ret.push_back(s);
            }
            nums.push_back(cur);
            swap(nums[i], nums[n - 1]);
        }
        return ret;
    }
};

```



---

## 50. Pow(x, n)

### Description

Implement `pow(x, n)`, which calculates  $x$  raised to the power  $n$  ( $x^n$ ).

Example 1:

Input: 2.00000, 10

Output: 1024.00000

Example 2:

Input: 2.10000, 3

Output: 9.26100

Example 3:

Input: 2.00000, -2

Output: 0.25000

Explanation:  $2^{-2} = 1/2^2 = 1/4 = 0.25$

Note:

$-100.0 < x < 100.0$

$n$  is a 32-bit signed integer, within the range  $[-2^{31}, 2^{31} - 1]$

### Solution

04/23/2020:

- [Discussion](#)
- Binary Exponentiation: given a positive power  $n$ , e.g.,  $n = 22$  (binary: `10110`) =  $16$  (binary: `10000`) +  $4$  (binary: `100`) +  $2$  (binary: `10`), then  $\text{pow}(x, n) = x^n = x^{22} = x^{(16 + 4 + 2)} = x^{16} * x^4 * x^2$ . For negative powers, we just flip the sign first, and once we calculate the value and return its reciprocal.
- Time complexity:  $O(n)$ ;
- Space complexity:  $O(1)$ .

```

class Solution {
public:
    double myPow(double x, int n) {
        long double ret = 1.0, pow_x = x;
        for (long m = n >= 0 ? (long)n : -1 * (long)n; m != 0; m >>= 1) {
            if (m & 1) ret *= pow_x;
            pow_x = pow_x * pow_x;
        }
        return n >= 0 ? ret : 1.0L / ret;
    }
};

```

## 51. N-Queens

### *Description*

The n-queens puzzle is the problem of placing n queens on an n×n chessboard such that no two queens attack each other.

Given an integer n, return all distinct solutions to the n-queens puzzle.

Each solution contains a distinct board configuration of the n-queens' placement, where 'Q' and '.' both indicate a queen and an empty space respectively.

Example:

Input: 4

Output: [

```

[".Q..", // Solution 1
 "...Q",
 "Q...",
 "..Q."],

```

```

["..Q.", // Solution 2
 "Q...",
 "...Q",
 ".Q.."]

```

]

Explanation: There exist two distinct solutions to the 4-queens puzzle as shown above.

### *Solution*

05/27/2020:

```
class Solution {
public:
    unordered_set<string> seen;
    vector<vector<string>> solveNQueens(int n) {
        vector<vector<string>> ret;
        vector<string> board(n, string(n, '.'));
        vector<bool> col(n, false);
        vector<bool> diag1(2 * n - 1, false);
        vector<bool> diag2(2 * n - 1, false);
        backtrack(0, n, board, col, diag1, diag2, ret);
        return ret;
    }

    void backtrack(int r, int n, vector<string>& board, vector<bool>& col,
vector<bool>& diag1, vector<bool>& diag2, vector<vector<string>>& ret) {
        if (r == n) {
            ret.push_back(board);
            return;
        }
        for (int c = 0; c < n; ++c) {
            if (!col[c] && !diag1[r + c] && !diag2[r - c + n - 1]) {
                board[r][c] = 'Q';
                col[c] = diag1[r + c] = diag2[r - c + n - 1] = true;
                backtrack(r + 1, n, board, col, diag1, diag2, ret);
                col[c] = diag1[r + c] = diag2[r - c + n - 1] = false;
                board[r][c] = '.';
            }
        }
    }
};
```

## 53. Maximum Subarray

*Description*

Given an integer array `nums`, find the contiguous subarray (containing at least one number) which has the largest sum and return its sum.

Example:

Input: `[-2,1,-3,4,-1,2,1,-5,4]`,

Output: 6

Explanation: `[4,-1,2,1]` has the largest sum = 6.

Follow up:

If you have figured out the  $O(n)$  solution, try coding another solution using the divide and conquer approach, which is more subtle.

*Solution*

01/13/2020 (Dynamic Programming):

```
class Solution {
public:
    int maxSubArray(vector<int>& nums) {
        int n = nums.size(), max_sum = nums[0];
        for (int i = 1; i < n; ++i) {
            if (nums[i - 1] > 0) nums[i] += nums[i - 1];
            max_sum = max(max_sum, nums[i]);
        }
        return max_sum;
    }
};
```

## 54. Spiral Matrix

*Description*

54. Spiral Matrix

Medium

1984

526

Add to List

Share

Given a matrix of  $m \times n$  elements ( $m$  rows,  $n$  columns), return all elements of the matrix in spiral order.

Example 1:

Input:

```
[
  [ 1, 2, 3 ],
  [ 4, 5, 6 ],
  [ 7, 8, 9 ]
]
```

Output: [1,2,3,6,9,8,7,4,5]

Example 2:

Input:

```
[
  [1, 2, 3, 4],
  [5, 6, 7, 8],
  [9,10,11,12]
]
```

Output: [1,2,3,4,8,12,11,10,9,5,6,7]

*Solution*

04/23/2020:

```
class Solution {
public:
    vector<int> spiralOrder(vector<vector<int>>& matrix) {
        int m = matrix.size();
        if (m == 0) return {};
        int n = matrix[0].size();
        if (n == 0) return {};
        vector<vector<int>> dir{ {0, 1}, {1, 0}, {0, -1}, {-1, 0} };
        vector<int> ret;
        int N = m * n, d = 0, i = 0, j = -1;
        while (ret.size() < N) {
            int ni = i + dir[d][0], nj = j + dir[d][1];
            while (ni < 0 || ni >= m || nj < 0 || nj >= n || matrix[ni][nj] ==
INT_MIN) {
                d = (d + 1) % 4;
                ni = i + dir[d][0], nj = j + dir[d][1];
            }
            ret.push_back(matrix[ni][nj]);
            matrix[ni][nj] = INT_MIN;
            i = ni;
            j = nj;
        }
        return ret;
    }
}
```

```
};
```

## 55. Jump Game

### Description

Given an array of non-negative integers, you are initially positioned at the first index of the array.

Each element in the array represents your maximum jump length at that position.

Determine if you are able to reach the last index.

Example 1:

Input: [2,3,1,1,4]

Output: true

Explanation: Jump 1 step from index 0 to 1, then 3 steps to the last index.

Example 2:

Input: [3,2,1,0,4]

Output: false

Explanation: You will always arrive at index 3 no matter what. Its maximum jump length is 0, which makes it impossible to reach the last index.

### Solution

04/24/2020:

```
class Solution {
public:
    bool canJump(vector<int>& nums) {
        int ret = 0, n = nums.size();
        for (int i = 0; i < n; ++i) {
            if (i <= ret) {
                ret = max(i + nums[i], ret);
            } else {
                break;
            }
        }
        return ret >= n - 1;
    }
};
```

## 59. Spiral Matrix II

### Description

Given a positive integer  $n$ , generate a square matrix filled with elements from 1 to  $n^2$  in spiral order.

Example:

Input: 3

Output:

```
[
  [ 1, 2, 3 ],
  [ 8, 9, 4 ],
  [ 7, 6, 5 ]
]
```

### Solution

### Discussion:

```
class Solution {
public:
    vector<vector<int>> generateMatrix(int n) {
        if (n <= 0) return {};
        vector<vector<int>> dir{ {0, 1}, {1, 0}, {0, -1}, {-1, 0} };
        vector<vector<int>> matrix(n, vector<int>(n, 0));
        int i = 0, j = -1, d = 0, N = n * n, cnt = 0;
        while (cnt < N) {
            int ni = i + dir[d][0], nj = j + dir[d][1];
            while (ni < 0 || ni >= n || nj < 0 || nj >= n || matrix[ni][nj] != 0) {
                d = (d + 1) % 4;
                ni = i + dir[d][0];
                nj = j + dir[d][1];
            }
            i = ni;
            j = nj;
            matrix[i][j] = ++cnt;
        }
        return matrix;
    }
};
```

## 60. Permutation Sequence

## Description

The set  $[1,2,3,\dots,n]$  contains a total of  $n!$  unique permutations.

By listing and labeling all of the permutations in order, we get the following sequence for  $n = 3$ :

"123"

"132"

"213"

"231"

"312"

"321"

Given  $n$  and  $k$ , return the  $k$ th permutation sequence.

Note:

Given  $n$  will be between 1 and 9 inclusive.

Given  $k$  will be between 1 and  $n!$  inclusive.

Example 1:

Input:  $n = 3, k = 3$

Output: "213"

Example 2:

Input:  $n = 4, k = 9$

Output: "2314"

## Solution

05/26/2020:

```
class Solution {
public:
    string getPermutation(int n, int k) {
        string s;
        for (int i = 0; i < n; ++i) s += to_string(i + 1);
        while (--k > 0 && next_permutation(s.begin(), s.end()));
        return s;
    }
};
```

## 66. Plus One

## Description



Given a non-empty array of digits representing a non-negative integer, plus one to the integer.

The digits are stored such that the most significant digit is at the head of the list, and each element in the array contain a single digit.

You may assume the integer does not contain any leading zero, except the number 0 itself.

Example 1:

Input: [1,2,3]

Output: [1,2,4]

Explanation: The array represents the integer 123.

Example 2:

Input: [4,3,2,1]

Output: [4,3,2,2]

Explanation: The array represents the integer 4321.

*Solution*

05/23/2020:

```
class Solution {
public:
    vector<int> plusOne(vector<int>& digits) {
        if (digits.empty()) return {1};
        int carry = 1, n = digits.size();
        for (int i = n - 1; i >= 0; --i) {
            carry += digits[i];
            digits[i] = carry % 10;
            carry /= 10;
        }
        if (carry > 0) digits.insert(digits.begin(), carry);
        return digits;
    }
};
```

## 67. Add Binary

*Description*

Given two binary strings, return their sum (also a binary string).

The input strings are both non-empty and contains only characters 1 or 0.

Example 1:

Input: a = "11", b = "1"

Output: "100"

Example 2:

Input: a = "1010", b = "1011"

Output: "10101"

Constraints:

Each string consists only of '0' or '1' characters.

$1 \leq a.length, b.length \leq 10^4$

Each string is either "0" or doesn't contain any leading zero.

*Solution*

05/12/2020:

```
class Solution {
public:
    string addBinary(string a, string b) {
        int i = a.size() - 1, j = b.size() - 1;
        int carry = 0;
        string ret;
        for (; i >= 0 || j >= 0; --i, --j) {
            if (i >= 0) carry += a[i] - '0';
            if (j >= 0) carry += b[j] - '0';
            ret += (carry & 1) + '0';
            carry >>= 1;
        }
        if (carry > 0) ret += carry + '0';
        reverse(ret.begin(), ret.end());
        return ret;
    }
};
```

## 69. Sqrt(x)

*Description*

Implement int sqrt(int x).

Compute and return the square root of  $x$ , where  $x$  is guaranteed to be a non-negative integer.

Since the return type is an integer, the decimal digits are truncated and only the integer part of the result is returned.

Example 1:

Input: 4

Output: 2

Example 2:

Input: 8

Output: 2

Explanation: The square root of 8 is 2.82842..., and since the decimal part is truncated, 2 is returned.

*Solution*

04/28/2020:

```
class Solution {
public:
    int mySqrt(int x) {
        int lo = 0, hi = x;
        while (lo <= hi) {
            long long mid = lo + (hi - lo) / 2;
            long long sq = mid * mid;
            if (sq == x) {
                return mid;
            } else if (sq > x) {
                hi = mid - 1;
            } else {
                lo = mid + 1;
            }
        }
        return hi;
    }
};
```

## 70. Climbing Stairs

*Description*

You are climbing a stair case. It takes n steps to reach to the top.

Each time you can either climb 1 or 2 steps. In how many distinct ways can you climb to the top?

Note: Given n will be a positive integer.

Example 1:

Input: 2

Output: 2

Explanation: There are two ways to climb to the top.

1. 1 step + 1 step
2. 2 steps

Example 2:

Input: 3

Output: 3

Explanation: There are three ways to climb to the top.

1. 1 step + 1 step + 1 step
2. 1 step + 2 steps
3. 2 steps + 1 step

*Solution*

01/14/2020 (Dynamic Programming):

```
class Solution {
public:
    int climbStairs(int n) {
        int s1 = 1, s2 = 2;
        for (int i = 2; i < n; ++i) {
            s1 = s1 + s2;
            swap(s1, s2);
        }
        return n >= 2 ? s2 : s1;
    }
};
```

## 72. Edit Distance

*Description*

Given two words word1 and word2, find the minimum number of operations required to convert word1 to word2.

You have the following 3 operations permitted on a word:

Insert a character  
Delete a character  
Replace a character

Example 1:

Input: word1 = "horse", word2 = "ros"

Output: 3

Explanation:

horse -> rorse (replace 'h' with 'r')

rorse -> rose (remove 'r')

rose -> ros (remove 'e')

Example 2:

Input: word1 = "intention", word2 = "execution"

Output: 5

Explanation:

intention -> inention (remove 't')

inention -> enention (replace 'i' with 'e')

enention -> exention (replace 'n' with 'x')

exention -> exection (replace 'n' with 'c')

exection -> execution (insert 'u')

*Solution*

05/31/2020:

```
class Solution {
public:
    int minDistance(string word1, string word2) {
        int m = word1.size(), n = word2.size();
        if (m == 0) return n;
        if (n == 0) return m;
        const int INF = 1e9 + 5;
        vector<vector<int>> dp(m + 1, vector(n + 1, INF));
        dp[0][0] = 0;
        // dp[i][j]: the edit distance between word1[0..i] and word2[0..j]
        // dp[i][j] = dp[i - 1][j - 1]      if word1[i] == word2[j]: no operations
        // needed
        // dp[i][j] = min(
        //     if word1[i] != word2[j]
        //     dp[i - 1][j - 1] + 1,  replace word1[i] by word2[j]
        //     dp[i - 1][j] + 1,      delete character word1[i]
        //     dp[i][j - 1] + 1      delete character word2[j]
        // )
        for (int i = 1; i <= m; ++i) dp[i][0] = dp[i - 1][0] + 1;
        for (int j = 1; j <= n; ++j) dp[0][j] = dp[0][j - 1] + 1;
    }
};
```

```

for (int i = 1; i <= m; ++i) {
    for (int j = 1; j <= n; ++j) {
        if (word1[i - 1] == word2[j - 1]) {
            dp[i][j] = dp[i - 1][j - 1];
        } else {
            dp[i][j] = min(dp[i - 1][j], min(dp[i][j - 1], dp[i - 1][j - 1])) + 1;
        }
    }
}
return dp[m][n];
};

```

```

class Solution {
public:
    int minDistance(string word1, string word2) {
        int m = word1.size(), n = word2.size();
        vector<vector<int>> dp(m + 1, vector<int>(n + 1, INT_MAX));
        for (int i = 0; i <= m; ++i) {
            for (int j = 0; j <= n; ++j) {
                if (i == 0) {
                    dp[i][j] = j;
                } else if (j == 0) {
                    dp[i][j] = i;
                } else if (word1[i - 1] == word2[j - 1]) {
                    dp[i][j] = dp[i - 1][j - 1];
                } else {
                    dp[i][j] = min({dp[i - 1][j - 1], dp[i - 1][j], dp[i][j - 1]}) + 1;
                }
            }
        }
        return dp[m][n];
    }
};

```

## 74. Search a 2D Matrix

### Description

Write an efficient algorithm that searches for a value in an  $m \times n$  matrix. This matrix has the following properties:

Integers in each row are sorted from left to right.

The first integer of each row is greater than the last integer of the previous row.

Example 1:

Input:

```
matrix = [  
  [1, 3, 5, 7],  
  [10, 11, 16, 20],  
  [23, 30, 34, 50]  
]
```

target = 3

Output: true

Example 2:

Input:

```
matrix = [  
  [1, 3, 5, 7],  
  [10, 11, 16, 20],  
  [23, 30, 34, 50]  
]
```

target = 13

Output: false

*Solution*

05/23/2020:

```
class Solution {  
public:  
    bool searchMatrix(vector<vector<int>>& matrix, int target) {  
        if (matrix.empty() || matrix[0].empty()) return false;  
        int m = matrix.size(), n = matrix[0].size();  
        int lo = 0, hi = m * n - 1;  
        while (lo <= hi) {  
            int mid = lo + (hi - lo) / 2;  
            int i = mid / n, j = mid % n;  
            if (matrix[i][j] == target) {  
                return true;  
            } else if (matrix[i][j] < target) {  
                lo = mid + 1;  
            } else {  
                hi = mid - 1;  
            }  
        }  
        return false;  
    }  
};
```

## 75. Sort Colors

### Description

Given an array with  $n$  objects colored red, white or blue, sort them in-place so that objects of the same color are adjacent, with the colors in the order red, white and blue.

Here, we will use the integers 0, 1, and 2 to represent the color red, white, and blue respectively.

Note: You are not suppose to use the library's sort function for this problem.

Example:

Input: [2,0,2,1,1,0]

Output: [0,0,1,1,2,2]

Follow up:

A rather straight forward solution is a two-pass algorithm using counting sort. First, iterate the array counting number of 0's, 1's, and 2's, then overwrite array with total number of 0's, then 1's and followed by 2's.

Could you come up with a one-pass algorithm using only constant space?

### Solution

05/12/2020:

```
class Solution {
public:
    void sortColors(vector<int>& nums) {
        sort(nums.begin(), nums.end());
    }
};
```

## 77. Combinations

### Description

Given two integers  $n$  and  $k$ , return all possible combinations of  $k$  numbers out of  $1 \dots n$ .

Example:



Input: n = 4, k = 2

Output:

```
[
  [2,4],
  [3,4],
  [2,3],
  [1,2],
  [1,3],
  [1,4],
]
```

*Solution*

05/27/2020:

```
class Solution {
public:
    vector<vector<int>> combine(int n, int k) {
        vector<int> combination;
        vector<vector<int>> ret;
        vector<bool> chosen(n + 1, false);
        backtrack(1, n, k, chosen, combination, ret);
        return ret;
    }

    void backtrack(int j, int n, int k, vector<bool>& chosen, vector<int>&
combination, vector<vector<int>>& ret) {
        if (combination.size() == k) ret.push_back(combination);
        for (int i = j; i <= n; ++i) {
            if (chosen[i] == true || (!combination.empty() && i < combination.back()))
                continue;
            chosen[i] = true;
            combination.push_back(i);
            backtrack(j + 1, n, k, chosen, combination, ret);
            chosen[i] = false;
            combination.pop_back();
        }
    }
};
```

## 78. Subsets

*Description*

Given a set of distinct integers, `nums`, return all possible subsets (the power set).

Note: The solution set must not contain duplicate subsets.

Example:

Input: `nums = [1,2,3]`

Output:

```
[
  [3],
  [1],
  [2],
  [1,2,3],
  [1,3],
  [2,3],
  [1,2],
  []
]
```

*Solution*

05/26/2020:

```
class Solution {
public:
    vector<int> subset;
    vector<vector<int>> s;
    int n;
    vector<vector<int>> subsets(vector<int>& nums) {
        n = nums.size();
        search(0);
        vector<vector<int>> ret;
        for (auto& ss : s) {
            vector<int> tmp;
            for (auto& i : ss) {
                tmp.push_back(nums[i]);
            }
            ret.push_back(tmp);
        }
        return ret;
    }

    void search(int k) {
        if (k == n) {
            s.push_back(subset);
        } else {
            subset.push_back(k);
        }
    }
};
```

```

        search(k + 1);
        subset.pop_back();
        search(k + 1);
    }
}
};

```

```

class Solution {
public:
    vector<vector<int>> subsets(vector<int>& nums) {
        vector<vector<int>> ret;
        vector<int> subset;
        backtrack(0, nums, subset, ret);
        return ret;
    }

    void backtrack(int k, vector<int>& nums, vector<int>& subset,
vector<vector<int>>& ret) {
        int n = nums.size();
        if (k == n) {
            ret.push_back(subset);
            return;
        }
        subset.push_back(nums[k]);
        backtrack(k + 1, nums, subset, ret);
        subset.pop_back();
        backtrack(k + 1, nums, subset, ret);
    }
};

```

```

class Solution {
public:
    vector<vector<int>> subsets(vector<int>& nums) {
        vector<vector<int>> ret{ {} };
        for (auto& n : nums) {
            int sz = ret.size();
            for (int i = 0; i < sz; ++i) {
                ret.push_back(ret[i]);
                ret.back().push_back(n);
            }
        }
        return ret;
    }
};

```

```

public:

```

```

vector<vector<int>> subsets(vector<int>& nums) {
    vector<vector<int>> ret;
    vector<int> subset;
    backtrack(0, nums, subset, ret);
    return ret;
}

void backtrack(int k, vector<int>& nums, vector<int>& subset,
vector<vector<int>>& ret) {
    ret.push_back(subset);
    for (int i = k; i < (int)nums.size(); ++i) {
        subset.push_back(nums[i]);
        backtrack(i + 1, nums, subset, ret);
        subset.pop_back();
    }
}
};

```

## 79. Word Search

### *Description*

Given a 2D board and a word, find if the word exists in the grid.

The word can be constructed from letters of sequentially adjacent cell, where "adjacent" cells are those horizontally or vertically neighboring. The same letter cell may not be used more than once.

Example:

```

board =
[
  ['A','B','C','E'],
  ['S','F','C','S'],
  ['A','D','E','E']
]

```

Given word = "ABCCED", return true.

Given word = "SEE", return true.

Given word = "ABCB", return false.

Constraints:

board and word consists only of lowercase and uppercase English letters.  
 $1 \leq \text{board.length} \leq 200$

```
1 <= board[i].length <= 200
1 <= word.length <= 10^3
```

*Solution*

05/27/2020:

```
class Solution {
public:
    int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
    bool exist(vector<vector<char>>& board, string word) {
        if (word.empty()) return true;
        if (board.empty() || board[0].empty()) return false;
        int m = board.size(), n = board[0].size(), k = word.size();
        for (int i = 0; i < m; ++i)
            for (int j = 0; j < n; ++j)
                if (backtrack(board, word, i, j, 0))
                    return true;
        return false;
    }

    bool backtrack(vector<vector<char>>& board, string word, int i, int j, int k)
    {
        int m = board.size(), n = board[0].size(), sz = word.size();
        if (board[i][j] != word[k]) return false;
        if (k == sz - 1) return true;
        board[i][j] = '#';
        for (int d = 0; d < 4; ++d) {
            int ni = i + dir[d][0], nj = j + dir[d][1];
            if (ni < 0 || ni >= m || nj < 0 || nj >= n) continue;
            if (backtrack(board, word, ni, nj, k + 1)) return true;
        }
        board[i][j] = word[k];
        return false;
    }
};
```

## 80. Remove Duplicates from Sorted Array II

*Description*

Given a sorted array `nums`, remove the duplicates in-place such that duplicates appeared at most twice and return the new length.

Do not allocate extra space for another array, you must do this by modifying the input array in-place with  $O(1)$  extra memory.

Example 1:

Given `nums = [1,1,1,2,2,3]`,

Your function should return `length = 5`, with the first five elements of `nums` being `1, 1, 2, 2` and `3` respectively.

It doesn't matter what you leave beyond the returned length.

Example 2:

Given `nums = [0,0,1,1,1,1,2,3,3]`,

Your function should return `length = 7`, with the first seven elements of `nums` being modified to `0, 0, 1, 1, 2, 3` and `3` respectively.

It doesn't matter what values are set beyond the returned length.

Clarification:

Confused why the returned value is an integer but your answer is an array?

Note that the input array is passed in by reference, which means modification to the input array will be known to the caller as well.

Internally you can think of this:

```
// nums is passed in by reference. (i.e., without making a copy)
int len = removeDuplicates(nums);

// any modification to nums in your function would be known by the caller.
// using the length returned by your function, it prints the first len elements.
for (int i = 0; i < len; i++) {
    print(nums[i]);
}
```

*Solution*

05/27/2020:

```
class Solution {
public:
    int removeDuplicates(vector<int>& nums) {
        if (nums.size() <= 2) return nums.size();
        int slow = 0, fast = 1, n = nums.size(), cnt = 1;
        for (; fast < n; ++fast) {
            while (fast < n && nums[fast] == nums[fast - 1]) {
```

```

    if (cnt <= 1) {
        ++cnt;
        nums[++slow] = nums[fast];
    }
    ++fast;
}
if (fast < n && nums[fast] != nums[slow]) {
    cnt = 1;
    nums[++slow] = nums[fast];
}
}
return slow + 1;
}
};

```

## 81. Search in Rotated Sorted Array II

### Description

Suppose an array sorted in ascending order is rotated at some pivot unknown to you beforehand.

(i.e.,  $[0,0,1,2,2,5,6]$  might become  $[2,5,6,0,0,1,2]$ ).

You are given a target value to search. If found in the array return true, otherwise return false.

Example 1:

Input:  $nums = [2,5,6,0,0,1,2]$ ,  $target = 0$

Output: true

Example 2:

Input:  $nums = [2,5,6,0,0,1,2]$ ,  $target = 3$

Output: false

Follow up:

This is a follow up problem to Search in Rotated Sorted Array, where  $nums$  may contain duplicates.

Would this affect the run-time complexity? How and why?

### Solution

05/27/2020:

```

class Solution {
public:
    bool search(vector<int>& nums, int target) {
        sort(nums.begin(), nums.end());
        return binary_search(nums.begin(), nums.end(), target);
    }
};

```

## 83. Remove Duplicates from Sorted List

### Description

Given a sorted linked list, delete all duplicates such that each element appear only once.

Example 1:

Input: 1->1->2

Output: 1->2

Example 2:

Input: 1->1->2->3->3

Output: 1->2->3

### Solution

05/27/2020:

```

/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
public:
    ListNode* deleteDuplicates(ListNode* head) {
        if (head == nullptr || head->next == nullptr) return head;
        ListNode *slow = head, *fast = head->next;
        while (fast->next) {
            while (fast && fast->val == slow->val && fast->next) fast = fast->next;

```



```

    if (fast && fast->val != slow->val) {
        slow->next = fast;
        slow = slow->next;
    }
}
if (slow->val == fast->val) {
    slow->next = fast->next;
}
return head;
}
};

```

## 84. Largest Rectangle in Histogram

### Description

Given  $n$  non-negative integers representing the histogram's bar height where the width of each bar is 1, find the area of largest rectangle in the histogram.

Above is a histogram where width of each bar is 1, given height = [2,1,5,6,2,3].

The largest rectangle is shown in the shaded area, which has area = 10 unit.

### Solution

01/14/2020 (Stack):

```

class Solution {
public:
    // Approach 1: Brute Force
    // int largestRectangleArea(vector<int>& heights) {
    //     int max_area = INT_MIN;
    //     for (int i = 0; i < heights.size(); ++i) {
    //         for (int j = i; j < heights.size(); ++j) {
    //             int min_height = *min_element(heights.begin() + i, heights.begin() +
j + 1);
    //             max_area = max(max_area, min_height * (j - i + 1));
    //         }
    //     }
    //     return max_area;
    // }

    // Approach 2: Better Brute Force
    // int largestRectangleArea(vector<int>& heights) {
    //     int n = heights.size(), max_area = n > 0 ? heights[0] : 0;

```

```

// for (int i = 0; i < heights.size(); ++i) {
//     int l = i, r = i;
//     for (; l >= 0 && heights[l] >= heights[i]; --l);
//     for (; r < heights.size() && heights[r] >= heights[i]; ++r);
//     max_area = max(max_area, heights[i] * (--r - ++l + 1));
// }
// return max_area;
// }

// Approach 3: Divide and Conquer
// int maxRectangleArea(vector<int>& heights, int left, int right) {
//     if (left > right) return 0;
//     int min_height_index = left;
//     for (int i = left; i <= right; ++i) {
//         if (heights[min_height_index] > heights[i])
//             min_height_index = i;
//     }
//     return max(max(maxRectangleArea(heights, left, min_height_index - 1),
maxRectangleArea(heights, min_height_index + 1, right)),
heights[min_height_index] * (right - left + 1));
// }
//
// int largestRectangleArea(vector<int>& heights) {
//     return maxRectangleArea(heights, 0, heights.size() - 1);
// }

// Approach 4: Stack
int largestRectangleArea(vector<int>& heights) {
    stack<int> s;
    s.push(-1);
    int max_area = 0;
    for (int i = 0; i < heights.size(); ++i) {
        while (s.top() != -1 && heights[s.top()] >= heights[i]) {
            int k = s.top();
            s.pop();
            max_area = max(max_area, heights[k] * (i - s.top() - 1));
        }
        s.push(i);
    }
    while (s.top() != -1) {
        int k = s.top();
        s.pop();
        max_area = max(max_area, heights[k] * ((int)heights.size() - s.top() -
1));
    }
    return max_area;
}
};

```

## 88. Merge Sorted Array

### Description

Given two sorted integer arrays `nums1` and `nums2`, merge `nums2` into `nums1` as one sorted array.

Note:

The number of elements initialized in `nums1` and `nums2` are `m` and `n` respectively. You may assume that `nums1` has enough space (size that is greater or equal to `m + n`) to hold additional elements from `nums2`.

Example:

Input:

`nums1 = [1,2,3,0,0,0]`, `m = 3`

`nums2 = [2,5,6]`, `n = 3`

Output: `[1,2,2,3,5,6]`

### Solution

05/23/2020:

```
class Solution {
public:
    void merge(vector<int>& nums1, int m, vector<int>& nums2, int n) {
        copy(nums2.begin(), nums2.end(), nums1.begin() + m);
        inplace_merge(nums1.begin(), nums1.begin() + m, nums1.end());
    }
};
```

## 94. Binary Tree Inorder Traversal

### Description

Given a binary tree, return the inorder traversal of its nodes' values.

Example:

Input: [1,null,2,3]

```
  1
   \
    2
   /
  3
```

Output: [1,3,2]

Follow up: Recursive solution is trivial, could you do it iteratively?

*Solution*

05/20/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
class Solution {
public:
    vector<int> inorderTraversal(TreeNode* root) {
        if (!root) return {};
        stack<TreeNode*> st;
        st.push(root);
        unordered_set<TreeNode*> visited;
        vector<int> ret;
        while (!st.empty()) {
            TreeNode* cur = st.top(); st.pop();
            if (cur->left && visited.count(cur->left) == 0) {
                if (cur->right) st.push(cur->right);
                st.push(cur);
                st.push(cur->left);
            } else if (!cur->left) {
                if (cur->right) st.push(cur->right);
                ret.push_back(cur->val);
                visited.insert(cur);
            } else {
                ret.push_back(cur->val);
                visited.insert(cur);
            }
        }
    }
};
```

```

    }
}
return ret;
}
};

```

```

class Solution {
public:
    vector<int> inorderTraversal(TreeNode* root) {
        if (!root) return {};
        vector<int> ret = inorderTraversal(root->left);
        ret.push_back(root->val);
        vector<int> right = inorderTraversal(root->right);
        ret.insert(ret.end(), right.begin(), right.end());
        return ret;
    }
};

```

```

class Solution {
public:
    vector<int> inorderTraversal(TreeNode* root) {
        if (!root) return {};
        vector<int> ret;
        TreeNode* cur = root;
        stack<TreeNode*> st;
        while (cur || !st.empty()) {
            while (cur) {
                st.push(cur);
                cur = cur->left;
            }
            cur = st.top(); st.pop();
            ret.push_back(cur->val);
            cur = cur->right;
        }
        return ret;
    }
};

```

## 96. Unique Binary Search Trees

### *Description*

Given  $n$ , how many structurally unique BST's (binary search trees) that store values  $1 \dots n$ ?

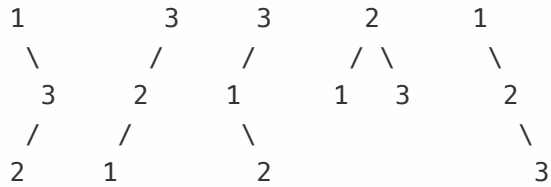
Example:

Input: 3

Output: 5

Explanation:

Given  $n = 3$ , there are a total of 5 unique BST's:



*Solution*

06/24/2020:

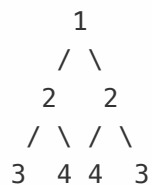
```
class Solution {
public:
    int numTrees(int n) {
        long C = 1;
        for (int i = 0; i < n; ++i) C = C * 2 * (2 * i + 1) / (i + 2);
        return C;
    }
};
```

## 101. Symmetric Tree

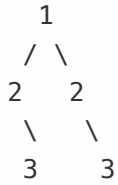
*Description*

Given a binary tree, check whether it is a mirror of itself (ie, symmetric around its center).

For example, this binary tree [1,2,2,3,4,4,3] is symmetric:



But the following [1,2,2,null,3,null,3] is not:



Follow up: Solve it both recursively and iteratively.

*Solution*

**Discussion:** This problem is very similar to <https://leetcode.com/problems/same-tree/>. Non-recursive DFS:

```

class Solution {
public:
    bool isSymmetric(TreeNode* root) {
        stack<pair<TreeNode*, TreeNode*>> st;
        st.emplace(root, root);
        while (!st.empty()) {
            pair<TreeNode*, TreeNode*> cur = st.top(); st.pop();
            if (cur.first == nullptr && cur.second == nullptr) continue;
            if (cur.first != nullptr && cur.second == nullptr) return false;
            if (cur.first == nullptr && cur.second != nullptr) return false;
            if (cur.first->val != cur.second->val) return false;
            st.emplace(cur.first->left, cur.second->right);
            st.emplace(cur.first->right, cur.second->left);
        }
        return true;
    }
};

```

Recursive DFS:

```

class Solution {
public:
    bool isSymmetric(TreeNode* root) {
        if (!root) return true;
        return isSubtreeSymmetric(root->left, root->right);
    }

    bool isSubtreeSymmetric(TreeNode* left, TreeNode* right) {
        if (!left && !right) return true;
        if ((!left && right) || (left && !right)) return false;
        return left->val == right->val && isSubtreeSymmetric(left->left, right->right) && isSubtreeSymmetric(left->right, right->left);
    }
};

```

## 102. Binary Tree Level Order Traversal

### Description

Given a binary tree, return the level order traversal of its nodes' values. (ie, from left to right, level by level).

For example:

Given binary tree [3,9,20,null,null,15,7],

```

    3
   / \
  9  20
   / \
  15  7

```

return its level order traversal as:

```

[
  [3],
  [9,20],
  [15,7]
]

```

### Solution

04/26/2020:

```

/**
 * Definition for a binary tree node.
 * struct TreeNode {

```



```

*   int val;
*   TreeNode *left;
*   TreeNode *right;
*   TreeNode(int x) : val(x), left(NULL), right(NULL) {}
* };
*/
class Solution {
public:
    vector<vector<int>> levelOrder(TreeNode* root) {
        if (root == nullptr) return {};
        vector<vector<int>> ret;
        queue<TreeNode*> q;
        q.push(root);
        while (!q.empty()) {
            int n = q.size();
            vector<int> level;
            for (int i = 0; i < n; ++i) {
                TreeNode* cur = q.front(); q.pop();
                level.push_back(cur->val);
                if (cur->left) q.push(cur->left);
                if (cur->right) q.push(cur->right);
            }
            ret.push_back(level);
        }
        return ret;
    }
};

```

## 103. Binary Tree Zigzag Level Order Traversal

### *Description*

Given a binary tree, return the zigzag level order traversal of its nodes' values. (ie, from left to right, then right to left for the next level and alternate between).

For example:

Given binary tree [3,9,20,null,null,15,7],

```

    3
   / \
  9  20
   / \
  15  7

```

return its zigzag level order traversal as:

```
[
  [3],
  [20,9],
  [15,7]
]
```

*Solution*

[Discussion](#) 04/26/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
class Solution {
public:
    vector<vector<int>> zigzagLevelOrder(TreeNode* root) {
        vector<vector<int>> ret;
        if (root == nullptr) return ret;
        queue<pair<TreeNode*, int>> q;
        q.emplace(root, 0);
        while (!q.empty()) {
            int n = q.size();
            vector<int> level;
            int depth = q.front().second;
            for (int i = 0; i < n; ++i) {
                pair<TreeNode*, int> cur = q.front(); q.pop();
                level.push_back(cur.first->val);
                if (cur.first->left) q.emplace(cur.first->left, cur.second + 1);
                if (cur.first->right) q.emplace(cur.first->right, cur.second + 1);
            }
            if (depth % 2 != 0) reverse(level.begin(), level.end());
            ret.push_back(level);
        }
        return ret;
    }
};
```

Better:

```
class Solution {
public:
```

```

vector<vector<int>> zigzagLevelOrder(TreeNode* root) {
    vector<vector<int>> ret;
    if (root == nullptr) return ret;
    queue<TreeNode*> q;
    q.push(root);
    for (int depth = 0; !q.empty(); ++depth) {
        int n = q.size();
        vector<int> level;
        for (int j = 0; j < n; ++j) {
            TreeNode* cur = q.front(); q.pop();
            level.push_back(cur->val);
            if (cur->left) q.push(cur->left);
            if (cur->right) q.push(cur->right);
        }
        if (depth % 2 != 0) reverse(level.begin(), level.end());
        ret.push_back(level);
    }
    return ret;
}
};

```

## 105. Construct Binary Tree from Preorder and Inorder Traversal

### Description

Given preorder and inorder traversal of a tree, construct the binary tree.

Note:

You may assume that duplicates do not exist in the tree.

For example, given

preorder = [3,9,20,15,7]

inorder = [9,3,15,20,7]

Return the following binary tree:

```

    3
   / \
  9  20
   / \
  15  7

```

### Solution

05/20/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
 * };
 */
class Solution {
public:
    TreeNode* buildTree(vector<int>& preorder, vector<int>& inorder, int p_start
= 0, int p_stop = INT_MAX, int i_start = 0, int i_stop = INT_MAX) {
        if (p_stop == INT_MAX) p_stop = i_stop = inorder.size() - 1;
        if (i_start > i_stop) return nullptr;
        int i = i_start;
        for (; i <= i_stop; ++i)
            if (inorder[i] == preorder[p_start])
                break;
        TreeNode* root = new TreeNode(preorder[p_start]);
        root->left = buildTree(preorder, inorder, p_start + 1, p_start + 1 + (i -
1) - i_start, i_start, i - 1);
        root->right = buildTree(preorder, inorder, p_start + 1 + (i - 1) - i_start
+ 1, p_stop, i + 1, i_stop);
        return root;
    }
};
```

## 109. Convert Sorted List to Binary Search Tree

### *Description*

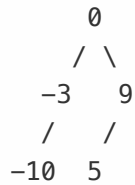
Given a singly linked list where elements are sorted in ascending order, convert it to a height balanced BST.

For this problem, a height-balanced binary tree is defined as a binary tree in which the depth of the two subtrees of every node never differ by more than 1.

Example:

Given the sorted linked list: [-10,-3,0,5,9],

One possible answer is: [0,-3,9,-10,null,5], which represents the following height balanced BST:



*Solution*

05/27/2020:

```
/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
 * };
 */
class Solution {
public:
    TreeNode* sortedListToBST(ListNode* head) {
        if (head == nullptr) return nullptr;
        if (head->next == nullptr) return new TreeNode(head->val);
        ListNode *slow = head, *fast = head, *preSlow = head;
        while (fast && fast->next) {
            preSlow = slow;
            slow = slow->next;
            fast = fast->next->next;
        }
    }
};
```

```

preSlow->next = nullptr;
TreeNode* root = new TreeNode(slow->val);
root->left = sortedListToBST(head);
root->right = sortedListToBST(slow->next);
return root;
}
};

```

## 110. Balanced Binary Tree

### *Description*

Given a binary tree, determine if it is height-balanced.

For this problem, a height-balanced binary tree is defined as:

a binary tree in which the left and right subtrees of every node differ in height by no more than 1.

Example 1:

Given the following tree [3,9,20,null,null,15,7]:

```

  3
 / \
9  20
 / \
15  7

```

Return true.

Example 2:

Given the following tree [1,2,2,3,3,null,null,4,4]:

```

  1
 / \
 2  2
 / \
 3  3
 / \
 4  4

```

Return false.

Solution

05/27/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
 * };
 */
class Solution {
public:
    bool isBalanced(TreeNode* root) {
        if (root == nullptr) return true;
        return abs(depth(root->left) - depth(root->right)) <= 1 && isBalanced(root->left) && isBalanced(root->right);
    }

    int depth(TreeNode* root) {
        if (root == nullptr) return 0;
        return 1 + max(depth(root->left), depth(root->right));
    }
};
```

## 114. Flatten Binary Tree to Linked List

*Description*

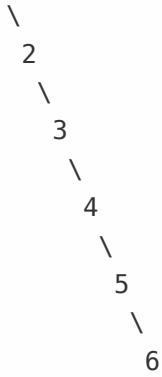
Given a binary tree, flatten it to a linked list in-place.

For example, given the following tree:

```
    1
   / \
  2   5
 / \   \
3  4   6
```

The flattened tree should look like:

```
1
```



*Solution*

05/27/2020:

```

/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
 *     right(right) {}
 * };
 */
class Solution {
public:
    void flatten(TreeNode* root) {
        if (root == nullptr) return;
        flatten(root->left);
        flatten(root->right);
        if (root->left) {
            TreeNode* cur = root->left;
            while (cur->right) cur = cur->right;
            cur->right = root->right;
            root->right = root->left;
            root->left = nullptr;
        }
    }
};

```

```

class Solution {
public:
    void flatten(TreeNode* root) {
        if (root == nullptr) return;

```



```

TreeNode *list = new TreeNode(0), *l = list;
stack<TreeNode*> st;
st.push(root);
while (!st.empty()) {
    TreeNode* cur = st.top(); st.pop();
    if (cur->right) st.push(cur->right);
    if (cur->left) st.push(cur->left);
    l->right = cur;
    cur->left = nullptr;
    l = l->right;
}
root = list->right;
}
};

```

## 118. Pascal's Triangle

### *Description*

Given a non-negative integer numRows, generate the first numRows of Pascal's triangle.

In Pascal's triangle, each number is the sum of the two numbers directly above it.

Example:

Input: 5

Output:

```

[
  [1],
  [1,1],
  [1,2,1],
  [1,3,3,1],
  [1,4,6,4,1]
]

```

### *Solution*

05/27/2020:

```

class Solution {
public:
    vector<vector<int>> generate(int numRows) {
        if (numRows < 0) return {};
        vector<vector<int>> ret(numRows);
        for (int i = 0; i < numRows; ++i) {
            ret[i].resize(i + 1, 1);
            for (int j = 1; j < i; ++j)
                ret[i][j] = ret[i - 1][j] + ret[i - 1][j - 1];
        }
        return ret;
    }
};

```

## 119. Pascal's Triangle II

### Description

Given a non-negative index  $k$  where  $k \leq 33$ , return the  $k$ th index row of the Pascal's triangle.

Note that the row index starts from 0.

In Pascal's triangle, each number is the sum of the two numbers directly above it.

Example:

Input: 3

Output: [1,3,3,1]

Follow up:

Could you optimize your algorithm to use only  $O(k)$  extra space?

### Solution

05/27/2020:

```

class Solution {
public:
    vector<int> getRow(int rowIndex) {
        if (rowIndex <= 1) return vector<int>(rowIndex + 1, 1);
        vector<int> ret(2, 1);
        for (int i = 2; i <= rowIndex; ++i) {

```

```

    vector<int> nextRow(i + 1, 1);
    for (int j = 1; j < i; ++j)
        nextRow[j] = ret[j - 1] + ret[j];
    ret = nextRow;
}
return ret;
}
};

```

```

class Solution {
public:
    vector<int> getRow(int rowIndex) {
        int n = rowIndex + 1;
        vector<vector<int>> ret(n);
        for (int i = 0; i < n; ++i) {
            ret[i].resize(i + 1, 1);
            for (int j = 1; j < i; ++j) {
                ret[i][j] = ret[i - 1][j - 1] + ret[i - 1][j];
            }
        }
        return ret.back();
    }
};

```

## 121. Best Time to Buy and Sell Stock

### Description

Say you have an array for which the  $i$ th element is the price of a given stock on day  $i$ .

If you were only permitted to complete at most one transaction (i.e., buy one and sell one share of the stock), design an algorithm to find the maximum profit.

Note that you cannot sell a stock before you buy one.

Example 1:

Input: [7,1,5,3,6,4]

Output: 5

Explanation: Buy on day 2 (price = 1) and sell on day 5 (price = 6), profit = 6 - 1 = 5.

Not 7 - 1 = 6, as selling price needs to be larger than buying price.

Example 2:

Input: [7,6,4,3,1]

Output: 0

Explanation: In this case, no transaction is done, i.e. max profit = 0.

*Solution*

01/13/2020 (Dynamic Programming):

```
class Solution {
public:
    int maxProfit(vector<int>& prices) {
        int n = prices.size();
        if (n <= 1) return 0;
        vector<int> diff(n - 1, 0);
        for (int i = 0; i < n - 1; ++i) {
            diff[i] = prices[i + 1] - prices[i];
        }
        int max_sum = max(0, diff[0]);
        for (int i = 1; i < n - 1; ++i) {
            if (diff[i - 1] > 0) diff[i] += diff[i - 1];
            max_sum = max(diff[i], max_sum);
        }
        return max_sum;
    }
};
```

## 129. Sum Root to Leaf Numbers

*Description*

Given a binary tree containing digits from 0-9 only, each root-to-leaf path could represent a number.

An example is the root-to-leaf path 1->2->3 which represents the number 123.

Find the total sum of all root-to-leaf numbers.

Note: A leaf is a node with no children.

Example:

Input: [1,2,3]

```
  1
 / \
```

2 3

Output: 25

Explanation:

The root-to-leaf path 1->2 represents the number 12.

The root-to-leaf path 1->3 represents the number 13.

Therefore,  $sum = 12 + 13 = 25$ .

Example 2:

Input: [4,9,0,5,1]

```
    4
   / \
  9  0
 / \
5  1
```

Output: 1026

Explanation:

The root-to-leaf path 4->9->5 represents the number 495.

The root-to-leaf path 4->9->1 represents the number 491.

The root-to-leaf path 4->0 represents the number 40.

Therefore,  $sum = 495 + 491 + 40 = 1026$ .

*Solution*

06/26/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
 *     right(right) {}
 * };
 */
class Solution {
public:
    int sumNumbers(TreeNode* root) {
        long long ret = 0, num = 0;
        backtrack(num, ret, root);
        return ret;
    }

    void backtrack(long long& num, long long& ret, TreeNode* root) {
        if (!root) return;
        if (!root->left && !root->right) {
```

```

        ret += num * 10 + root->val;
        return;
    }
    num = num * 10 + root->val;
    backtrack(num, ret, root->left);
    backtrack(num, ret, root->right);
    num /= 10;
}
};

```

## 130. Surrounded Regions

### Description

Given a 2D board containing 'X' and 'O' (the letter O), capture all regions surrounded by 'X'.

A region is captured by flipping all 'O's into 'X's in that surrounded region.

Example:

```

X X X X
X O O X
X X O X
X O X X

```

After running your function, the board should be:

```

X X X X
X X X X
X X X X
X O X X

```

Explanation:

Surrounded regions shouldn't be on the border, which means that any 'O' on the border of the board are not flipped to 'X'. Any 'O' that is not on the border and it is not connected to an 'O' on the border will be flipped to 'X'. Two cells are connected if they are adjacent cells connected horizontally or vertically.

### Solution

05/08/2020 [Discussion](#):

The idea is borrowed from one of the assignments of the course [Algorithms \(Part 1\)](#):

1. We assume there is a virtual cell (labeled as  $m * n$ ) that the node on the boundary of the board is

automatically connected to it.

2. We only need to deal with the cells that are '0': merge these neighboring '0' cells.
3. Traverse '0' cells which are stored in visited: if the cell is connected to the virtual cell do nothing; otherwise, change it to 'X'.

```
class UnionFind {
private:
    vector<int> id;
    vector<int> sz;

public:
    UnionFind (int n) {
        id.resize(n);
        iota(id.begin(), id.end(), 0);
        sz.resize(n, 1);
    }

    int find(int x) {
        if (x == id[x]) return x;
        return id[x] = find(id[x]);
    }

    bool connected(int x, int y) { return find(x) == find(y); }

    bool merge(int x, int y) {
        int i = find(x), j = find(y);
        if (i == j) return false;
        if (sz[j] > sz[i]) {
            swap(i, j);
            swap(sz[i], sz[j]);
        }
        id[j] = i;
        sz[i] += sz[j];
        return true;
    }
};

class Solution {
public:
    void solve(vector<vector<char>>& board) {
        if (board.empty() || board[0].empty()) return;
        int m = board.size(), n = board[0].size(), bound = m * n;
        UnionFind uf(m * n + 1);
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        unordered_set<int> visited;
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (board[i][j] == 'X' || visited.count(i * n + j)) continue;
            }
        }
    }
};
```

```

        if (i == 0 || i == m - 1 || j == 0 || j == n - 1) uf.merge(i * n + j,
bound);
        for (int d = 0; d < 4; ++d) {
            int ni = i + dir[d][0], nj = j + dir[d][1];
            if (ni >= 0 && ni < m && nj >= 0 && nj < n && board[ni][nj] == '0' &&
visited.count(i * n + j) == 0) {
                uf.merge(i * n + j, ni * n + nj);
            }
        }
        visited.insert(i * n + j);
    }
}
for (auto& p : visited) {
    if (!uf.connected(bound, p)) {
        int i = p / n, j = p % n;
        board[i][j] = 'X';
    }
}
};

```

Update: Inspired by the above, we can actually use a set uncaptured to store all the cells that are connected to the bounds. We can always use BFS which starts from the '0' cells on the boundary to find the inner cells which are connected to the boundary.

```

class Solution {
public:
    void solve(vector<vector<char>>& board) {
        if (board.empty() || board[0].empty()) return;
        int m = board.size(), n = board[0].size();
        unordered_set<int> uncaptured;
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (i == 0 || i == m - 1 || j == 0 || j == n - 1) {
                    if (board[i][j] == 'X') continue;
                    queue<int> q;
                    q.push(i * n + j);
                    while (!q.empty()) {
                        int sz = q.size();
                        for (int s = 0; s < sz; ++s) {
                            int cur = q.front(); q.pop();
                            if (uncaptured.count(cur) > 0) continue;
                            for (int d = 0; d < 4; ++d) {
                                int ni = cur / n + dir[d][0], nj = cur % n + dir[d][1];
                                if (ni >= 0 && ni < m && nj >= 0 && nj < n && board[ni][nj] ==
'0') {
                                    q.push(ni * n + nj);
                                }
                            }
                        }
                    }
                }
            }
        }
    }
};

```





```
class Solution {
public:
    int singleNumber(vector<int>& nums) {
        int ret = 0;
        for (auto& i : nums) ret ^= i;
        return ret;
    }
};
```

## 137. Single Number II

### *Description*

Given a non-empty array of integers, every element appears three times except for one, which appears exactly once. Find that single one.

Note:

Your algorithm should have a linear runtime complexity. Could you implement it without using extra memory?

Example 1:

Input: [2,2,3,2]

Output: 3

Example 2:

Input: [0,1,0,1,0,1,99]

Output: 99

### *Solution*

05/27/2020:

```

class Solution {
public:
    int singleNumber(vector<int>& nums) {
        unordered_map<int, int> cnt;
        for (auto& i : nums) ++cnt[i];
        for (auto& [k, v] : cnt)
            if (v == 1)
                return k;
        return 0;
    }
};

```

## 144. Binary Tree Preorder Traversal

### Description

Given a binary tree, return the preorder traversal of its nodes' values.

Example:

Input: [1,null,2,3]

```

  1
   \
    2
   /
  3

```

Output: [1,2,3]

Follow up: Recursive solution is trivial, could you do it iteratively?

### Solution

05/20/2020:

```

/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
class Solution {
public:

```

```

vector<int> preorderTraversal(TreeNode* root) {
    if (root == nullptr) return {};
    vector<int> ret;
    stack<TreeNode*> st;
    st.push(root);
    while (!st.empty()) {
        TreeNode* cur = st.top(); st.pop();
        ret.push_back(cur->val);
        if (cur->right) st.push(cur->right);
        if (cur->left) st.push(cur->left);
    }
    return ret;
}
};

```

```

class Solution {
public:
    vector<int> preorderTraversal(TreeNode* root) {
        if (root == nullptr) return {};
        vector<int> ret;
        ret.push_back(root->val);
        vector<int> l = preorderTraversal(root->left);
        vector<int> r = preorderTraversal(root->right);
        ret.insert(ret.end(), l.begin(), l.end());
        ret.insert(ret.end(), r.begin(), r.end());
        return ret;
    }
};

```

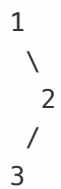
## 145. Binary Tree Postorder Traversal

*Description*

Given a binary tree, return the postorder traversal of its nodes' values.

Example:

Input: [1,null,2,3]



Output: [3,2,1]

Follow up: Recursive solution is trivial, could you do it iteratively?

*Solution*

05/20/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
class Solution {
public:
    vector<int> postorderTraversal(TreeNode* root) {
        vector<int> ret;
        if (root == nullptr) return ret;
        stack<TreeNode*> st;
        st.push(root);
        while (!st.empty()) {
            TreeNode* cur = st.top(); st.pop();
            ret.push_back(cur->val);
            if (cur->left != nullptr) st.push(cur->left);
            if (cur->right != nullptr) st.push(cur->right);
        }
        reverse(ret.begin(), ret.end());
        return ret;
    }
};
```

```
class Solution {
public:
```

```

vector<int> postorderTraversal(TreeNode* root) {
    if (!root) return {};
    stack<TreeNode*> st;
    st.push(root);
    vector<int> ret;
    unordered_set<TreeNode*> visited;
    while (!st.empty()) {
        TreeNode* cur = st.top(); st.pop();
        if ((!cur->left || visited.count(cur->left)) && (!cur->right ||
visited.count(cur->right))) {
            ret.push_back(cur->val);
            visited.insert(cur);
        } else {
            st.push(cur);
            if (cur->right) st.push(cur->right);
            if (cur->left) st.push(cur->left);
        }
    }
    return ret;
}
};

```

```

class Solution {
public:
    vector<int> postorderTraversal(TreeNode* root) {
        if (!root) return {};
        vector<int> ret = postorderTraversal(root->left);
        vector<int> right = postorderTraversal(root->right);
        ret.insert(ret.end(), right.begin(), right.end());
        ret.push_back(root->val);
        return ret;
    }
};

```

## 146. LRU Cache

### Description

Design and implement a data structure for Least Recently Used (LRU) cache. It should support the following operations: get and put.

get(key) – Get the value (will always be positive) of the key if the key exists in the cache, otherwise return -1.

put(key, value) – Set or insert the value if the key is not already present. When the cache reached its capacity, it should invalidate the least recently used item before inserting a new item.

The cache is initialized with a positive capacity.

Follow up:

Could you do both operations in  $O(1)$  time complexity?

Example:

```
LRUCache cache = new LRUCache( 2 /* capacity */ );

cache.put(1, 1);
cache.put(2, 2);
cache.get(1);      // returns 1
cache.put(3, 3);   // evicts key 2
cache.get(2);      // returns -1 (not found)
cache.put(4, 4);   // evicts key 1
cache.get(1);      // returns -1 (not found)
cache.get(3);      // returns 3
cache.get(4);      // returns 4
```

*Solution*

**Discussion:**

- Use an `unordered_map<int, int>` cache to store the key-value pairs.
- Use an `unordered_map<int, int>` priority to store the current rank for a certain key.
- Use a max-heap (`priority_queue`) to store the priorities of key's. The higher rank of the key, the earlier that the key will be removed from the cache.
- When we call `get(key)`, we update the rank for the current key by `priority[key] = rank` and then push this pair `{priority[key], key}` to the `priority_queue`. Therefore, the new pair would invalidate the original `{rank, key}` in the `priority_queue`. So inside the `put` method, we use a `while` loop to invalidate all those invalid pairs until the first valid pair with the highest rank.
- Time complexity: `get`:  $O(1)$ ; `put`:  $O(m)$ , where  $m$  is the number of operations of `get`.
- Space complexity:  $O(n + m)$ , where  $n$  is the number of key-value pairs and  $m$  is the number of operations of `get`.

```
class LRUCache {
private:
    unordered_map<int, int> cache;
    unordered_map<int, int> priority;
    priority_queue<pair<int, int>, vector<pair<int, int>>, less<pair<int, int>>>
    pq; // max-heap
    int capacity, rank;

public:
```

```

LRUCache(int capacity) {
    this->capacity = capacity;
    rank = INT_MAX;
    cache.clear();
    priority.clear();
}

int get(int key) {
    if (cache.count(key) == 0) {
        return -1;
    } else {
        priority[key] = rank--;
        priority.emplace(key, priority[key]);
        pq.emplace(priority[key], key);
        return cache[key];
    }
}

void put(int key, int value) {
    cache[key] = value;
    priority[key] = rank--;
    pq.emplace(priority[key], key);
    while (cache.size() > capacity) {
        pair<int, int> top = pq.top();
        while (!pq.empty() && priority[top.second] != top.first) {
            pq.pop();
            top = pq.top();
        }
        pq.pop();
        cache.erase(top.second);
        priority.erase(top.second);
    }
}
};

```

## 153. Find Minimum in Rotated Sorted Array

### *Description*

Suppose an array sorted in ascending order is rotated at some pivot unknown to you beforehand.

(i.e.,  $[0,1,2,4,5,6,7]$  might become  $[4,5,6,7,0,1,2]$ ).

Find the minimum element.



You may assume no duplicate exists in the array.

Example 1:

Input: [3,4,5,1,2]

Output: 1

Example 2:

Input: [4,5,6,7,0,1,2]

Output: 0

*Solution*

04/28/2020:

```
class Solution {
public:
    int findMin(vector<int>& nums) {
        if (nums.empty()) return -1;
        int n = nums.size();
        int lo = 0, hi = n - 1;
        while (lo < hi) {
            int mid = lo + (hi - lo) / 2;
            if (nums[mid] > nums[hi]) {
                lo = mid + 1;
            } else {
                hi = mid;
            }
        }
        return nums[hi];
    }
};
```

## 162. Find Peak Element

*Description*

A peak element is an element that is greater than its neighbors.

Given an input array `nums`, where `nums[i] ≠ nums[i+1]`, find a peak element and return its index.

The array may contain multiple peaks, in that case return the index to any one of the peaks is fine.

You may imagine that  $\text{nums}[-1] = \text{nums}[n] = -\infty$ .

Example 1:

Input:  $\text{nums} = [1,2,3,1]$

Output: 2

Explanation: 3 is a peak element and your function should return the index number 2.

Example 2:

Input:  $\text{nums} = [1,2,1,3,5,6,4]$

Output: 1 or 5

Explanation: Your function can return either index number 1 where the peak element is 2,

or index number 5 where the peak element is 6.

Note:

Your solution should be in logarithmic complexity.

*Solution*

04/28/2020:

```
class Solution {
public:
    int findPeakElement(vector<int>& nums) {
        if (nums.empty()) return -1;
        int lo = 0, hi = nums.size() - 1;
        while (lo < hi) {
            int mid = lo + (hi - lo) / 2;
            if (mid - 1 >= lo && nums[mid] < nums[mid - 1]) {
                hi = mid - 1;
            } else if (mid + 1 <= hi && nums[mid] < nums[mid + 1]) {
                lo = mid + 1;
            } else {
                return mid;
            }
        }
        return lo;
    }
};
```

## 163. Missing Ranges

## Description

Given a sorted integer array `nums`, where the range of elements are in the inclusive range `[lower, upper]`, return its missing ranges.

Example:

Input: `nums = [0, 1, 3, 50, 75]`, `lower = 0` and `upper = 99`,  
Output: `["2", "4->49", "51->74", "76->99"]`

## Solution

05/24/2020:

```
class Solution {
public:
    vector<string> findMissingRanges(vector<int>& nums, int lower, int upper) {
        if (nums.empty()) {
            string s = lower == upper ? to_string(lower) : to_string(lower) + "->" +
to_string(upper);
            return {s};
        }
        long long cur = nums.back(); nums.pop_back();
        vector<string> ranges;
        string s = cur + 1 < upper ? to_string(cur + 1) + "->" + to_string(upper) :
to_string(cur + 1);
        if (cur - 1 >= lower) ranges = findMissingRanges(nums, lower, cur - 1);
        if (cur + 1 <= upper) ranges.push_back(s);
        return ranges;
    }
};
```

```
class Solution {
public:
    vector<string> findMissingRanges(vector<int>& nums, int lower, int upper) {
        if (lower > upper) return {};
        if (nums.empty()) {
            if (lower == upper)
                return { to_string(lower) };
            else
                return { to_string(lower) + "->" + to_string(upper) };
        }
        long long cur = nums.back(); nums.pop_back();
        vector<string> ranges;
        if (cur - 1 >= INT_MIN) {
            ranges = findMissingRanges(nums, lower, cur - 1);
        }
    }
};
```

```

    if (cur + 1 < upper) {
        ranges.push_back(to_string(cur + 1) + "->" + to_string(upper));
    } else if (cur + 1 == upper) {
        ranges.push_back(to_string(cur + 1));
    }
    return ranges;
}
};

```

## 169. Majority Element

### Description

Given an array of size  $n$ , find the majority element. The majority element is the element that appears more than  $\lfloor n/2 \rfloor$  times.

You may assume that the array is non-empty and the majority element always exist in the array.

Example 1:

Input: [3,2,3]

Output: 3

Example 2:

Input: [2,2,1,1,1,2,2]

Output: 2

### Solution

05/06/2020:

```

class Solution {
public:
    int majorityElement(vector<int>& nums) {
        int n = nums.size();
        if (n == 0) return -1;
        unordered_map<int, int> s;
        int major = nums[0], cnt = 1;
        for (auto& i : nums) {
            if (++s[i] > cnt) {
                cnt = s[i];
                major = i;
            }
        }
    }
};

```

```
        return cnt > n / 2 ? major : -1;
    }
};
```

## 173. Binary Search Tree Iterator

### Description

Implement an iterator over a binary search tree (BST). Your iterator will be initialized with the root node of a BST.

Calling `next()` will return the next smallest number in the BST.

Example:

```
BSTIterator iterator = new BSTIterator(root);
iterator.next();    // return 3
iterator.next();    // return 7
iterator.hasNext(); // return true
iterator.next();    // return 9
iterator.hasNext(); // return true
iterator.next();    // return 15
iterator.hasNext(); // return true
iterator.next();    // return 20
iterator.hasNext(); // return false
```

Note:

`next()` and `hasNext()` should run in average  $O(1)$  time and uses  $O(h)$  memory, where  $h$  is the height of the tree.

You may assume that `next()` call will always be valid, that is, there will be at least a next smallest number in the BST when `next()` is called.

### Solution

05/24/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
```

```

*   int val;
*   TreeNode *left;
*   TreeNode *right;
*   TreeNode() : val(0), left(nullptr), right(nullptr) {}
*   TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
*   TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
* };
*/
class BSTIterator {
private:
    stack<TreeNode*> st;
    TreeNode* cur;

public:
    BSTIterator(TreeNode* root) {
        cur = root;
    }

    /** @return the next smallest number */
    int next() {
        while (cur) {
            st.push(cur);
            cur = cur->left;
        }
        cur = st.top(); st.pop();
        int ret = cur->val;
        cur = cur->right;
        return ret;
    }

    /** @return whether we have a next smallest number */
    bool hasNext() {
        return cur || !st.empty();
    }
};

/**
 * Your BSTIterator object will be instantiated and called as such:
 * BSTIterator* obj = new BSTIterator(root);
 * int param_1 = obj->next();
 * bool param_2 = obj->hasNext();
 */

```

```

class BSTIterator {
private:
    vector<TreeNode*> inorder;

```

```

vector<TreeNode*> inorderTraversal(TreeNode* root) {
    if (root == nullptr) return {};
    vector<TreeNode*> leftTraversal = inorderTraversal(root->left);
    vector<TreeNode*> rightTraversal = inorderTraversal(root->right);
    leftTraversal.push_back(root);
    leftTraversal.insert(leftTraversal.end(), rightTraversal.begin(),
rightTraversal.end());
    return leftTraversal;
}

public:
    BSTIterator(TreeNode* root) {
        inorder = inorderTraversal(root);
        reverse(inorder.begin(), inorder.end());
    }

    /** @return the next smallest number */
    int next() {
        TreeNode* ret = inorder.back();
        inorder.pop_back();
        return ret->val;
    }

    /** @return whether we have a next smallest number */
    bool hasNext() {
        return !inorder.empty();
    }
};

```

## 174. Dungeon Game

### *Description*

The demons had captured the princess (P) and imprisoned her in the bottom-right corner of a dungeon. The dungeon consists of  $M \times N$  rooms laid out in a 2D grid. Our valiant knight (K) was initially positioned in the top-left room and must fight his way through the dungeon to rescue the princess.

The knight has an initial health point represented by a positive integer. If at any point his health point drops to 0 or below, he dies immediately.

Some of the rooms are guarded by demons, so the knight loses health (negative integers) upon entering these rooms; other rooms are either empty (0's) or contain magic orbs that increase the knight's health (positive integers).

In order to reach the princess as quickly as possible, the knight decides to move only rightward or downward in each step.

Write a function to determine the knight's minimum initial health so that he is able to rescue the princess.

For example, given the dungeon below, the initial health of the knight must be at least 7 if he follows the optimal path RIGHT-> RIGHT -> DOWN -> DOWN.

```
-2 (K)  -3  3
-5  -10  1
10  30  -5 (P)
```

Note:

The knight's health has no upper bound.

Any room can contain threats or power-ups, even the first room the knight enters and the bottom-right room where the princess is imprisoned.

*Solution*

06/21/2020:

```
class Solution {
public:
    const int inf = numeric_limits<int>::max();
    vector<vector<int>> dp;
    int rows, cols;

    int getMinHealth(int curCell, int nextRow, int nextCol) {
        if (nextRow >= rows || nextCol >= cols) return inf;
        int nextCell = dp[nextRow][nextCol];
        return max(1, nextCell - curCell);
    }

    int calculateMinimumHP(vector<vector<int>>& dungeon) {
        rows = dungeon.size();
        cols = dungeon[0].size();
        dp.assign(rows, vector<int>(cols, inf));
        int curCell, rightHealth, downHealth, nextHealth, minHealth;
        for (int row = rows - 1; row >= 0; --row) {
            for (int col = cols - 1; col >= 0; --col) {
                curCell = dungeon[row][col];
                rightHealth = getMinHealth(curCell, row, col + 1);
                downHealth = getMinHealth(curCell, row + 1, col);
```



```

    nextHealth = min(rightHealth, downHealth);
    if (nextHealth != inf) {
        minHealth = nextHealth;
    } else {
        minHealth = curCell >= 0 ? 1 : 1 - curCell;
    }
    dp[row][col] = minHealth;
}
}
return dp[0][0];
}
};

```

## 198. House Robber

### Description

You are a professional robber planning to rob houses along a street. Each house has a certain amount of money stashed, the only constraint stopping you from robbing each of them is that adjacent houses have security system connected and it will automatically contact the police if two adjacent houses were broken into on the same night.

Given a list of non-negative integers representing the amount of money of each house, determine the maximum amount of money you can rob tonight without alerting the police.

Example 1:

Input: [1,2,3,1]

Output: 4

Explanation: Rob house 1 (money = 1) and then rob house 3 (money = 3).

Total amount you can rob = 1 + 3 = 4.

Example 2:

Input: [2,7,9,3,1]

Output: 12

Explanation: Rob house 1 (money = 2), rob house 3 (money = 9) and rob house 5 (money = 1).

Total amount you can rob = 2 + 9 + 1 = 12.

### Solution

01/14/2020 (Dynamic Programming):

```

class Solution {
public:
    int rob(vector<int>& nums) {
        if (nums.size() >= 2) nums[1] = max(nums[0], nums[1]);
        for (int i = 2; i < nums.size(); ++i)
            nums[i] = max(nums[i - 1], nums[i - 2] + nums[i]);
        return nums.size() > 0 ? nums.back() : 0;
    }
};

```

## 200. Number of Islands

### *Description*

Given a 2d grid map of '1's (land) and '0's (water), count the number of islands. An island is surrounded by water and is formed by connecting adjacent lands horizontally or vertically. You may assume all four edges of the grid are all surrounded by water.

Example 1:

Input:  
11110  
11010  
11000  
00000

Output: 1  
Example 2:

Input:  
11000  
11000  
00100  
00011

Output: 3

### *Solution*

05/08/2020 (Union-Find):

```

class UnionFind {
private:
    vector<int> id;

```

```

vector<int> sz;

public:
UnionFind(int n) {
    id.resize(n);
    iota(id.begin(), id.end(), 0);
    sz.resize(n, 1);
}

int find(int x) {
    if (x == id[x]) return x;
    return id[x] = find(id[x]);
}

bool merge(int x, int y) {
    int i = find(x), j = find(y);
    if (i == j) return false;
    if (sz[i] > sz[j]) {
        id[j] = i;
        sz[i] += sz[j];
    } else {
        id[i] = j;
        sz[j] += sz[i];
    }
    return true;
}
};

class Solution {
public:
int numIslands(vector<vector<char>>& grid) {
    if (grid.empty() || grid[0].empty()) return 0;
    int m = grid.size(), n = grid[0].size();
    UnionFind uf(m * n);
    int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
    for (int i = 0; i < m; ++i) {
        for (int j = 0; j < n; ++j) {
            if (grid[i][j] == '0') continue;
            for (int d = 0; d < 4; ++d) {
                int ni = i + dir[d][0], nj = j + dir[d][1];
                if (ni >= 0 && ni < m && nj >= 0 && nj < n && grid[ni][nj] == '1') {
                    uf.merge(i * n + j, ni * n + nj);
                }
            }
        }
    }
    unordered_set<int> components;
    for (int i = 0; i < m; ++i)
        for (int j = 0; j < n; ++j)

```

```

        if (grid[i][j] == '1')
            components.insert(uf.find(i * n + j));
    return components.size();
}
};

```

```

class Solution {
public:
    int numIslands(vector<vector<char>>& grid) {
        int m = grid.size();
        if (m == 0) return 0;
        int n = grid[0].size(), ret = 0;
        stack<pair<int, int>> st;
        int d[4][2] = { {0, -1}, {0, 1}, {-1, 0}, {1, 0} };
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (grid[i][j] == '1') {
                    ++ret;
                    st.emplace(i, j);
                    while (!st.empty()) {
                        pair<int, int> cur = st.top(); st.pop();
                        grid[cur.first][cur.second] = '0';
                        for (int k = 0; k < 4; ++k) {
                            int ni = cur.first + d[k][0], nj = cur.second + d[k][1];
                            if (ni > -1 && ni < m && nj > -1 && nj < n && grid[ni][nj] == '1')
                                st.emplace(ni, nj);
                        }
                    }
                }
            }
        }
        return ret;
    }
};

```

Iterative DFS:

```

class Solution {
public:
    int numIslands(vector<vector<char>>& grid) {
        int m = grid.size();
        if (m == 0) return 0;
        int n = grid[0].size(), ret = 0;
        stack<pair<int, int>> st;
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (grid[i][j] == '1') {

```

```

        ++ret;
        st.emplace(i, j);
        while (!st.empty()) {
            pair<int, int> cur = st.top(); st.pop();
            int ni = cur.first, nj = cur.second;
            grid[ni][nj] = '0';
            if (ni + 1 < m && grid[ni + 1][nj] == '1') st.emplace(ni + 1, nj);
            if (ni - 1 >= 0 && grid[ni - 1][nj] == '1') st.emplace(ni - 1, nj);
            if (nj + 1 < n && grid[ni][nj + 1] == '1') st.emplace(ni, nj + 1);
            if (nj - 1 >= 0 && grid[ni][nj - 1] == '1') st.emplace(ni, nj - 1);
        }
    }
}
return ret;
}
};

```

Recursive DFS:

```

class Solution {
public:
    int numIslands(vector<vector<char>>& grid) {
        int m = grid.size();
        if (m == 0) return 0;
        int n = grid[0].size();
        int ret = 0;
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (grid[i][j] == '1') {
                    ++ret;
                    dfs(grid, i, j, m, n);
                }
            }
        }
        return ret;
    }

    void dfs(vector<vector<char>>& grid, int i, int j, int m, int n) {
        if (i >= 0 && i < m && j >= 0 && j < n && grid[i][j] == '1') {
            grid[i][j] = '0';
            dfs(grid, i - 1, j, m, n);
            dfs(grid, i + 1, j, m, n);
            dfs(grid, i, j - 1, m, n);
            dfs(grid, i, j + 1, m, n);
        }
    }
};

```

---

## 205. Isomorphic Strings

### Description

Given two strings *s* and *t*, determine if they are isomorphic.

Two strings are isomorphic if the characters in *s* can be replaced to get *t*.

All occurrences of a character must be replaced with another character while preserving the order of characters. No two characters may map to the same character but a character may map to itself.

Example 1:

Input: *s* = "egg", *t* = "add"

Output: true

Example 2:

Input: *s* = "foo", *t* = "bar"

Output: false

Example 3:

Input: *s* = "paper", *t* = "title"

Output: true

Note:

You may assume both *s* and *t* have the same length.

### Solution

05/20/2020:

```
class Solution {
public:
    bool isIsomorphic(string s, string t) {
        unordered_map<char, unordered_set<char>> mpST, mpTS;
        for (int i = 0; i < (int)s.size(); ++i) {
            mpST[s[i]].insert(t[i]);
            mpTS[t[i]].insert(s[i]);
        }
        for (auto& m : mpST)
            if (m.second.size() > 1)
                return false;
        for (auto& m : mpTS)
            if (m.second.size() > 1)
                return false;
    }
};
```

```
    return true;
}
};
```

## 207. Course Schedule

### *Description*

There are a total of `numCourses` courses you have to take, labeled from `0` to `numCourses-1`.

Some courses may have prerequisites, for example to take course `0` you have to first take course `1`, which is expressed as a pair: `[0,1]`

Given the total number of courses and a list of prerequisite pairs, is it possible for you to finish all courses?

Example 1:

Input: `numCourses = 2, prerequisites = [[1,0]]`

Output: `true`

Explanation: There are a total of 2 courses to take.

To take course 1 you should have finished course 0. So it is possible.

Example 2:

Input: `numCourses = 2, prerequisites = [[1,0],[0,1]]`

Output: `false`

Explanation: There are a total of 2 courses to take.

To take course 1 you should have finished course 0, and to take course 0 you should also have finished course 1. So it is impossible.

Constraints:

The input `prerequisites` is a graph represented by a list of edges, not adjacency matrices. Read more about how a graph is represented.

You may assume that there are no duplicate edges in the input `prerequisites`.

$1 \leq \text{numCourses} \leq 10^5$

### *Solution*

05/29/2020:

```
class Solution {
public:
    bool canFinish(int numCourses, vector<vector<int>>& prerequisites) {
        vector<vector<int>> adj(numCourses);
        for (auto& p : prerequisites) adj[p[1]].push_back(p[0]);
        vector<bool> path(numCourses);
        for (int i = 0; i < numCourses; ++i)
            if (isCyclic(i, adj, path))
                return false;
        return true;
    }

    bool isCyclic(int i, vector<vector<int>>& adj, vector<bool>& path) {
        if (path[i]) return true;
        if (adj[i].empty()) return false;
        path[i] = true;
        bool ret = false;
        for (auto& j : adj[i]) {
            ret = isCyclic(j, adj, path);
            if (ret) break;
        }
        path[i] = false;
        return ret;
    }
};
```

## 208. Implement Trie (Prefix Tree)

### *Description*

Implement a trie with insert, search, and startsWith methods.

Example:

```
Trie trie = new Trie();
```

```
trie.insert("apple");
trie.search("apple"); // returns true
trie.search("app"); // returns false
trie.startsWith("app"); // returns true
trie.insert("app");
trie.search("app"); // returns true
```

Note:



You may assume that all inputs are consist of lowercase letters a-z.  
All inputs are guaranteed to be non-empty strings.

*Solution*

05/13/2020:

```
class Trie {
public:
    /** Initialize your data structure here. */
    Trie() {
        root = new Node();
    }

    /** Inserts a word into the trie. */
    void insert(string word) {
        Node* cur = root;
        for (auto& c : word) {
            if (!cur->mp[c])
                cur->mp[c] = new Node();
            cur = cur->mp[c];
        }
        cur->isWord = true;
    }

    /** Returns if the word is in the trie. */
    bool search(string word) {
        Node* cur = find(word);
        return cur != nullptr && cur->isWord;
    }

    /** Returns if there is any word in the trie that starts with the given
    prefix. */
    bool startsWith(string prefix) {
        Node* cur = find(prefix);
        return cur != nullptr;
    }

private:
    struct Node {
        bool isWord;
        unordered_map<char, Node*> mp;
        Node() : isWord(false) {}
    };

    Node* root;
};
```

```

Node* find(const string& word) {
    Node* cur = root;
    for (auto& c : word) {
        cur = cur->mp[c];
        if (!cur) break;
    }
    return cur;
}
};

```

```

/**
 * Your Trie object will be instantiated and called as such:
 * Trie* obj = new Trie();
 * obj->insert(word);
 * bool param_2 = obj->search(word);
 * bool param_3 = obj->startsWith(prefix);
 */

```

```

class Trie {
private:
    struct Node {
        bool isWord;
        vector<Node*> children;
        Node() { isWord = false; children.resize(26, nullptr); }
        ~Node() { for(auto& c : children) delete c; }
    };

    Node* root;

    Node* find(const string& word) {
        Node* cur = root;
        for (auto& c : word) {
            cur = cur->children[c - 'a'];
            if (!cur) break;
        }
        return cur;
    }

public:
    /** Initialize your data structure here. */
    Trie() {
        root = new Node();
    }

    /** Inserts a word into the trie. */
    void insert(string word) {
        Node* cur = root;
        for (auto& c : word) {

```

```

        if (!cur->children[c - 'a'])
            cur->children[c - 'a'] = new Node();
        cur = cur->children[c - 'a'];
    }
    cur->isWord = true;
}

/** Returns if the word is in the trie. */
bool search(string word) {
    Node* cur = find(word);
    return cur && cur->isWord;
}

/** Returns if there is any word in the trie that starts with the given
prefix. */
bool startsWith(string prefix) {
    Node* cur = find(prefix);
    return cur;
}
};

/**
 * Your Trie object will be instantiated and called as such:
 * Trie* obj = new Trie();
 * obj->insert(word);
 * bool param_2 = obj->search(word);
 * bool param_3 = obj->startsWith(prefix);
 */

```

```

struct Node {
    char character;
    bool isWord;
    vector<Node*> children;
    Node(char c, bool i) : character(c), isWord(i) {}
    ~Node() {
        for (auto& c : children) delete c;
    }
};

class Trie {
private:
    Node *root;

public:
    /** Initialize your data structure here. */
    Trie() {
        root = new Node('r', false);
    }
}

```

```

/** Inserts a word into the trie. */
void insert(string word) {
    int n = word.size();
    Node* cur = root;
    for (int i = 0; i < n; ++i) {
        bool foundCharacter = false;
        for (auto& c : cur->children) {
            if (word[i] == c->character) {
                cur = c;
                foundCharacter = true;
                break;
            }
        }
        if (!foundCharacter) {
            Node* c = new Node(word[i], false);
            cur->children.push_back(c);
            cur = c;
        }
    }
    cur->isWord = true;
}

```

```

/** Returns if the word is in the trie. */
bool search(string word) {
    Node* cur = root;
    int n = word.size();
    for (int i = 0; i < n; ++i) {
        bool foundCharacter = false;
        for (auto& c : cur->children) {
            if (c->character == word[i]) {
                foundCharacter = true;
                cur = c;
                break;
            }
        }
        if (!foundCharacter) return false;
    }
    return cur->isWord;
}

```

/\*\* Returns if there is any word in the trie that starts with the given prefix. \*/

```

bool startsWith(string prefix) {
    Node* cur = root;
    int n = prefix.size();
    for (int i = 0; i < n; ++i) {
        bool foundCharacter = false;
        for (auto& c : cur->children) {

```

```

        if (c->character == prefix[i]) {
            foundCharacter = true;
            cur = c;
            break;
        }
    }
    if (!foundCharacter) return false;
}
return true;
}
};

/**
 * Your Trie object will be instantiated and called as such:
 * Trie* obj = new Trie();
 * obj->insert(word);
 * bool param_2 = obj->search(word);
 * bool param_3 = obj->startsWith(prefix);
 */

```

## 219. Contains Duplicate II

### Description

Given an array of integers and an integer  $k$ , find out whether there are two distinct indices  $i$  and  $j$  in the array such that  $\text{nums}[i] = \text{nums}[j]$  and the absolute difference between  $i$  and  $j$  is at most  $k$ .

Example 1:

Input:  $\text{nums} = [1,2,3,1]$ ,  $k = 3$

Output: true

Example 2:

Input:  $\text{nums} = [1,0,1,1]$ ,  $k = 1$

Output: true

Example 3:

Input:  $\text{nums} = [1,2,3,1,2,3]$ ,  $k = 2$

Output: false

### Solution

05/20/2020:

```

class Solution {
public:
    bool containsNearbyDuplicate(vector<int>& nums, int k) {
        unordered_map<int, int> mp;
        int n = nums.size();
        for (int i = 0; i < n; ++i) {
            if (mp.count(nums[i]) > 0) {
                if (i - mp[nums[i]] <= k) {
                    return true;
                }
            }
            mp[nums[i]] = i;
        }
        return false;
    }
};

```

```

class Solution {
public:
    bool containsNearbyDuplicate(vector<int>& nums, int k) {
        unordered_set<int> visited;
        for (int i = 0; i < (int)nums.size(); ++i) {
            if (visited.count(nums[i]) > 0) return true;
            visited.insert(nums[i]);
            if (i - k >= 0) visited.erase(nums[i - k]);
        }
        return false;
    }
};

```

## 221. Maximal Square

*Description*

Given a 2D binary matrix filled with 0's and 1's, find the largest square containing only 1's and return its area.

Example:

Input:

```
1 0 1 0 0
1 0 1 1 1
1 1 1 1 1
1 0 0 1 0
```

Output: 4

*Solution*

04/27/2020 (Dynamic Programming):

```
class Solution {
public:
    int maximalSquare(vector<vector<char>>& matrix) {
        int m = matrix.size();
        if (m == 0) return 0;
        int n = matrix[0].size();
        int ret = 0;
        vector<vector<int>> dp(m + 1, vector<int>(n + 1, 0));
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                dp[i + 1][j + 1] = matrix[i][j] - '0';
                if (dp[i + 1][j + 1] > 0) {
                    dp[i + 1][j + 1] = min(dp[i][j], min(dp[i][j + 1], dp[i + 1][j])) + 1;
                    ret = max(ret, dp[i + 1][j + 1]);
                }
            }
        }
        return ret * ret;
    }
};
```

More space-efficient Dynamic Programming:

```
class Solution {
public:
    int maximalSquare(vector<vector<char>>& matrix) {
        if (matrix.size() == 0 || matrix[0].size() == 0) return 0;
        int m = matrix.size(), n = matrix[0].size();
        int ret = 0;
```

```

vector<int> dp(n + 1, 0);
for (int i = 0; i < m; ++i) {
    vector<int> tmp(dp);
    for (int j = 0; j < n; ++j) {
        dp[j + 1] = matrix[i][j] - '0';
        if (dp[j + 1] > 0) {
            dp[j + 1] = min(tmp[j], min(tmp[j + 1], dp[j])) + 1;
            ret = max(ret, dp[j + 1]);
        }
    }
}
return ret * ret;
};

```

## 222. Count Complete Tree Nodes

### Description

Given a complete binary tree, count the number of nodes.

Note:

Definition of a complete binary tree from Wikipedia:

In a complete binary tree every level, except possibly the last, is completely filled, and all nodes in the last level are as far left as possible. It can have between 1 and  $2^h$  nodes inclusive at the last level  $h$ .

Example:

Input:

```

    1
   / \
  2   3
 / \ /
4  5 6

```

Output: 6

### Solution

06/23/2020:

```

/**
 * Definition for a binary tree node.

```



```

* struct TreeNode {
*     int val;
*     TreeNode *left;
*     TreeNode *right;
*     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
* };
*/
class Solution {
public:
    int countNodes(TreeNode* root) {
        if (!root) return 0;
        queue<TreeNode*> q;
        q.push(root);
        int ret = 0;
        while (!q.empty()) {
            int sz = q.size();
            ret += sz;
            for (int i = 0; i < sz; ++i) {
                TreeNode* cur = q.front(); q.pop();
                if (cur->right) q.push(cur->right);
                if (cur->left) q.push(cur->left);
            }
        }
        return ret;
    }
};

```

## 226. Invert Binary Tree

### *Description*

Invert a binary tree.

Example:

Input:

```

      4
     / \
    2   7
   / \ / \
  1  3 6  9

```

Output:

```

      4
     / \

```

```
    7    2
   / \  / \
  9  6 3  1
```

Trivia:

This problem was inspired by this original tweet by Max Howell:

Google: 90% of our engineers use the software you wrote (Homebrew), but you can't invert a binary tree on a whiteboard so f\*\*\* off.

*Solution*

06/01/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
 * };
 */
class Solution {
public:
    TreeNode* invertTree(TreeNode* root) {
        if (root == nullptr) return nullptr;
        swap(root->left, root->right);
        invertTree(root->left);
        invertTree(root->right);
        return root;
    }
};
```

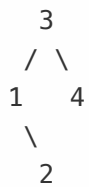
## 230. Kth Smallest Element in a BST

*Description*

Given a binary search tree, write a function `kthSmallest` to find the `k`th smallest element in it.

Example 1:

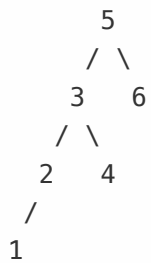
Input: root = [3,1,4,null,2], k = 1



Output: 1

Example 2:

Input: root = [5,3,6,2,4,null,null,1], k = 3



Output: 3

Follow up:

What if the BST is modified (insert/delete operations) often and you need to find the kth smallest frequently? How would you optimize the kthSmallest routine?

Constraints:

The number of elements of the BST is between 1 to  $10^4$ .

You may assume k is always valid,  $1 \leq k \leq$  BST's total elements.

*Solution*

05/20/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
 * };
 */
```

```

class Solution {
public:
    int kthSmallest(TreeNode* root, int k) {
        stack<TreeNode*> st;
        TreeNode* cur = root;
        while (cur || !st.empty()) {
            while (cur) {
                st.push(cur);
                cur = cur->left;
            }
            cur = st.top(); st.pop();
            if (--k == 0) return cur->val;
            cur = cur->right;
        }
        return 0;
    }
};

```

```

class Solution {
public:
    int kthSmallest(TreeNode* root, int k) {
        stack<TreeNode*> st;
        st.push(root);
        unordered_set<TreeNode*> visited;
        while (!st.empty()) {
            TreeNode* cur = st.top(); st.pop();
            if (cur->left && visited.count(cur->left) == 0) {
                if (cur->right) st.push(cur->right);
                st.push(cur);
                st.push(cur->left);
            } else {
                if (!cur->left && cur->right) st.push(cur->right);
                visited.insert(cur);
                if (--k == 0) return cur->val;
            }
        }
        return 0;
    }
};

```

```

class Solution {
public:
    int n, ans;
    int kthSmallest(TreeNode* root, int k) {
        n = k;
        inorder(root);
        return ans;
    }
};

```

```

}

void inorder(TreeNode* root) {
    if (!root) return;
    inorder(root->left);
    if (--n == 0) ans = root->val;
    inorder(root->right);
}
};

```

```

class Solution {
public:
    int kthSmallest(TreeNode* root, int k) {
        vector<int> list = BST2vector(root);
        return list[k - 1];
    }

    vector<int> BST2vector(TreeNode* root) {
        if (!root) return {};
        vector<int> ret = BST2vector(root->left);
        ret.push_back(root->val);
        vector<int> right = BST2vector(root->right);
        ret.insert(ret.end(), right.begin(), right.end());
        return ret;
    }
};

```

## 231. Power of Two

### *Description*

Given an integer, write a function to determine if it is a power of two.

Example 1:

Input: 1

Output: true

Explanation:  $2^0 = 1$

Example 2:

Input: 16

Output: true

Explanation:  $2^4 = 16$

Example 3:

Input: 218  
Output: false

*Solution*

06/07/2020:

```
class Solution {  
public:  
    bool isPowerOfTwo(int n) {  
        return n > 0 && __builtin_popcount(n) == 1;  
    }  
};
```

## 237. Delete Node in a Linked List

*Description*

Write a function to delete a node (except the tail) in a singly linked list, given only access to that node.

Given linked list -- head = [4,5,1,9], which looks like following:

Example 1:

Input: head = [4,5,1,9], node = 5

Output: [4,1,9]

Explanation: You are given the second node with value 5, the linked list should become 4 -> 1 -> 9 after calling your function.

Example 2:

Input: head = [4,5,1,9], node = 1

Output: [4,5,9]

Explanation: You are given the third node with value 1, the linked list should become 4 -> 5 -> 9 after calling your function.

Note:

The linked list will have at least two elements.  
All of the nodes' values will be unique.

The given node will not be the tail and it will always be a valid node of the linked list.  
Do not return anything from your function.

*Solution*

05/10/2020:

```
/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode(int x) : val(x), next(NULL) {}
 * };
 */
class Solution {
public:
    void deleteNode(ListNode* node) {
        if (node == nullptr) return;
        if (node->next == nullptr) {
            node = node->next;
        } else {
            node->val = node->next->val;
            node->next = node->next->next;
            node = node->next;
        }
    }
};
```

06/02/2020:

```
class Solution {
public:
    void deleteNode(ListNode* node) {
        *node = *(node->next);
    }
};
```

```
class Solution {
public:
    void deleteNode(ListNode* node) {
        node->val = node->next->val;
        node->next = node->next->next;
    }
};
```

## 242. Valid Anagram

### Description

Given two strings *s* and *t* , write a function to determine if *t* is an anagram of *s*.

Example 1:

Input: *s* = "anagram", *t* = "nagaram"

Output: true

Example 2:

Input: *s* = "rat", *t* = "car"

Output: false

Note:

You may assume the string contains only lowercase alphabets.

Follow up:

What if the inputs contain unicode characters? How would you adapt your solution to such case?

### Solution

05/17/2020:

```
class Solution {
public:
    bool isAnagram(string s, string t) {
        vector<int> cntS(26, 0), cntT(26, 0);
        for (int i = 0; i < (int)s.size(); ++i) ++cntS[s[i] - 'a'];
        for (int i = 0; i < (int)t.size(); ++i) ++cntT[t[i] - 'a'];
        for (int i = 0; i < 26; ++i) {
            if (cntS[i] != cntT[i]) {
                return false;
            }
        }
        return true;
    }
};
```



```
class Solution {
public:
    bool isAnagram(string s, string t) {
        sort(s.begin(), s.end());
        sort(t.begin(), t.end());
        return s == t;
    }
};
```

## 246. Strobogrammatic Number

### *Description*

A strobogrammatic number is a number that looks the same when rotated 180 degrees (looked at upside down).

Write a function to determine if a number is strobogrammatic. The number is represented as a string.

Example 1:

Input: "69"

Output: true

Example 2:

Input: "88"

Output: true

Example 3:

Input: "962"

Output: false

### *Solution*

05/18/2020:

```

class Solution {
public:
    bool isStrobogrammatic(string num) {
        unordered_map<char, char> mp{ {'0', '0'}, {'1', '1'}, {'6', '9'}, {'8',
'8'}, {'9', '6'} };
        string stro = "";
        for (auto& n : num) {
            if (mp.count(n) == 0) return false;
            stro = mp[n] + stro;
        }
        return stro == num;
    }
};

```

## 249. Group Shifted Strings

### Description

Given a string, we can "shift" each of its letter to its successive letter, for example: "abc" -> "bcd". We can keep "shifting" which forms the sequence:

"abc" -> "bcd" -> ... -> "xyz"

Given a list of strings which contains only lowercase alphabets, group all strings that belong to the same shifting sequence.

Example:

Input: ["abc", "bcd", "acef", "xyz", "az", "ba", "a", "z"],

Output:

```

[
  ["abc","bcd","xyz"],
  ["az","ba"],
  ["acef"],
  ["a","z"]
]

```

### Solution

05/25/2020:

```

class Solution {
public:
    vector<vector<string>> groupStrings(vector<string>& strings) {
        vector<vector<string>> ret;
        unordered_map<string, vector<string>> shiftedStrings;

```

```

    for (auto& s : strings) shiftedStrings[base(s)].push_back(s);
    for (auto& [key, val] : shiftedStrings) ret.push_back(val);
    return ret;
}

string base(const string& s) {
    string ts;
    for (auto& c : s) ts += (c - s[0] + 26) % 26 + 'a';
    return ts;
}
};

```

```

class Solution {
public:
    vector<vector<string>> groupStrings(vector<string>& strings) {
        unordered_map<int, vector<string>> shiftedStrings;
        for (auto& s : strings) shiftedStrings[hash(s)].push_back(s);
        vector<vector<string>> ret;
        for (auto& m : shiftedStrings) ret.push_back(m.second);
        return ret;
    }

    int hash(const string& s) {
        long long h = 1;
        if (s.empty()) return h;
        int offset = s[0] - 'a';
        string ts;
        const int MOD = 1e9 + 7;
        for (auto& c : s) ts += (c - 'a' - offset + 26) % 26 + 'a';
        for (auto& c : ts) h = (h * 31 + c - 'a') % MOD;
        return h;
    }
};

```

## 250. Count Univalve Subtrees

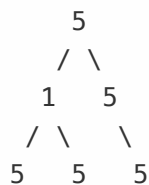
### Description

Given a binary tree, count the number of uni-value subtrees.

A Uni-value subtree means all nodes of the subtree have the same value.

Example :

Input: root = [5,1,5,5,5,null,5]



Output: 4

*Solution*

05/25/2020:

```

/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
 * };
 */
class Solution {
public:
    int countUnivalSubtrees(TreeNode* root) {
        pair<int, bool> ret = dfs(root);
        return ret.first;
    }

    pair<int, bool> dfs(TreeNode* root) {
        if (root == nullptr) return {0, true};
        if (root->left == nullptr && root->right == nullptr) return {1, true};
        pair<int, bool> countLeft = dfs(root->left);
        pair<int, bool> countRight = dfs(root->right);
        bool isUniqueTree = true;
        if (countLeft.first > 0) {
            isUniqueTree = isUniqueTree && root->val == root->left->val &&
countLeft.second;
        }
        if (countRight.first > 0) {
            isUniqueTree = isUniqueTree && root->val == root->right->val &&
countRight.second;
        }
        return {isUniqueTree + countLeft.first + countRight.first, isUniqueTree};
    }
}

```

```
};
```

```
class Solution {
public:
    int countUnivalSubtrees(TreeNode* root) {
        int cnt = 0;
        dfs(root, cnt);
        return cnt;
    }

    bool dfs(TreeNode* root, int& cnt) {
        if (root == nullptr) return true;
        if (root->left == nullptr && root->right == nullptr) {
            ++cnt;
            return true;
        }
        bool isUniqueTree = true;
        bool checkLeft = dfs(root->left, cnt);
        bool checkRight = dfs(root->right, cnt);
        if (root->left) isUniqueTree = isUniqueTree && root->val == root->left->val
&& checkLeft;
        if (root->right) isUniqueTree = isUniqueTree && root->val == root->right-
>val && checkRight;
        if (isUniqueTree) ++cnt;
        return isUniqueTree;
    }
};
```

## 251. Flatten 2D Vector

### *Description*

Design and implement an iterator to flatten a 2d vector. It should support the following operations: next and hasNext.

Example:

```
Vector2D iterator = new Vector2D([[1,2],[3],[4]]);
```

```
iterator.next(); // return 1
iterator.next(); // return 2
iterator.next(); // return 3
iterator.hasNext(); // return true
```

```
iterator.hasNext(); // return true
iterator.next(); // return 4
iterator.hasNext(); // return false
```

#### Notes:

Please remember to RESET your class variables declared in Vector2D, as static/class variables are persisted across multiple test cases. Please see here for more details.

You may assume that next() call will always be valid, that is, there will be at least a next element in the 2d vector when next() is called.

#### Follow up:

As an added challenge, try to code it using only iterators in C++ or iterators in Java.

#### *Solution*

05/25/2020:

```
class Vector2D {
private:
    vector<vector<int>> v;
    int i, j;

public:
    Vector2D(vector<vector<int>>& v) {
        this->v = v;
        i = j = 0;
        while (i < (int)v.size() && v[i].empty()) ++i;
    }

    int next() {
        int ret = v[i][j];
        if (i < (int)v.size() && ++j >= (int)v[i].size()) {
            j = 0, ++i;
            while (i < (int)v.size() && v[i].empty()) ++i;
        }
        return ret;
    }

    bool hasNext() {
        return i < (int)v.size() && j < (int)v[i].size();
    }
};
```

```
/**
 * Your Vector2D object will be instantiated and called as such:
 * Vector2D* obj = new Vector2D(v);
 * int param_1 = obj->next();
 * bool param_2 = obj->hasNext();
 */
```

```
class Vector2D {
private:
    vector<int> flattenedVector;

public:
    Vector2D(vector<vector<int>>& v) {
        flattenedVector.clear();
        for (auto& nums : v) {
            for (auto& i : nums) {
                flattenedVector.push_back(i);
            }
        }
        reverse(flattenedVector.begin(), flattenedVector.end());
    }

    int next() {
        int ret = flattenedVector.back();
        flattenedVector.pop_back();
        return ret;
    }

    bool hasNext() {
        return !flattenedVector.empty();
    }
};
```

## 252. Meeting Rooms

*Description*

Given an array of meeting time intervals consisting of start and end times  $[[s_1, e_1], [s_2, e_2], \dots]$  ( $s_i < e_i$ ), determine if a person could attend all meetings.

Example 1:

Input:  $[[0,30], [5,10], [15,20]]$

Output: false

Example 2:

Input:  $[[7,10], [2,4]]$

Output: true

NOTE: input types have been changed on April 15, 2019. Please reset to default code definition to get new method signature.

*Solution*

05/26/2020:

```
class Solution {
public:
    bool canAttendMeetings(vector<vector<int>>& intervals) {
        if (intervals.empty()) return true;
        sort(intervals.begin(), intervals.end());
        for (int i = 1; i < (int)intervals.size(); ++i)
            if (intervals[i - 1][1] > intervals[i][0])
                return false;
        return true;
    }
};
```

## 253. Meeting Rooms II

*Description*



Given an array of meeting time intervals consisting of start and end times  $[[s_1, e_1], [s_2, e_2], \dots]$  ( $s_i < e_i$ ), find the minimum number of conference rooms required.

Example 1:

Input:  $[[0, 30], [5, 10], [15, 20]]$

Output: 2

Example 2:

Input:  $[[7, 10], [2, 4]]$

Output: 1

NOTE: input types have been changed on April 15, 2019. Please reset to default code definition to get new method signature.

*Solution*

05/27/2020:

```
class Solution {
public:
    int minMeetingRooms(vector<vector<int>>& intervals) {
        vector<int> startTime, endTime;
        for (auto& i : intervals) {
            startTime.push_back(i[0]);
            endTime.push_back(i[1]);
        }
        sort(startTime.begin(), startTime.end());
        sort(endTime.begin(), endTime.end());
        int ret = 0, cnt = 0;
        auto startIt = startTime.begin();
        auto endIt = endTime.begin();
        for (; startIt != startTime.end() || endIt != endTime.end(); ) {
            if (endIt == endTime.end() || (startIt != startTime.end() && *startIt <
*endIt)) {
                ret = max(ret, ++cnt);
                ++startIt;
            } else {
                --cnt;
                ++endIt;
            }
        }
        return ret;
    }
};
```

## 257. Binary Tree Paths

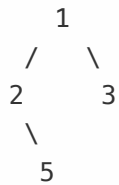
### Description

Given a binary tree, return all root-to-leaf paths.

Note: A leaf is a node with no children.

Example:

Input:



Output: ["1->2->5", "1->3"]

Explanation: All root-to-leaf paths are: 1->2->5, 1->3

### Solution

05/26/2020:

```
class Solution {
public:
    vector<string> binaryTreePaths(TreeNode* root) {
        string path;
        vector<string> ret;
        backtrack(root, path, ret);
        return ret;
    }

    void backtrack(TreeNode* root, const string& path, vector<string>& ret) {
        if (root == nullptr) return;
        if (root->left == nullptr && root->right == nullptr) {
            path.empty() ? ret.push_back(to_string(root->val)) : ret.push_back(path +
"->" + to_string(root->val));
            return;
        }
        string addend = (path.empty() ? "" : "->") + to_string(root->val);
        backtrack(root->left, path + addend, ret);
        backtrack(root->right, path + addend, ret);
    }
};
```

```

class Solution {
public:
    vector<string> binaryTreePaths(TreeNode* root) {
        if (root == nullptr) return {};
        stack<pair<TreeNode*, string>> st;
        st.emplace(root, "");
        vector<string> ret;
        while (!st.empty()) {
            pair<TreeNode*, string> cur = st.top(); st.pop();
            cur.second += "->" + to_string(cur.first->val);
            if (cur.first->left == nullptr && cur.first->right == nullptr)
                ret.push_back(cur.second.substr(2, cur.second.size() - 2));
            if (cur.first->right != nullptr) st.emplace(cur.first->right, cur.second);
            if (cur.first->left != nullptr) st.emplace(cur.first->left, cur.second);
        }
        return ret;
    }
};

```

```

class Solution {
public:
    vector<string> binaryTreePaths(TreeNode* root) {
        vector<string> ret;
        string p = "";
        dfs(root, p, ret);
        return ret;
    }

    void dfs(TreeNode* root, string p, vector<string>& ret) {
        if (root == nullptr) return;
        p += "->" + to_string(root->val);
        if (root->left != nullptr) dfs(root->left, p, ret);
        if (root->right != nullptr) dfs(root->right, p, ret);
        if (root->left == nullptr && root->right == nullptr)
            ret.push_back(p.substr(2, (int)p.size() - 2));
    }
};

```

## 265. Paint House II

*Description*

There are a row of  $n$  houses, each house can be painted with one of the  $k$  colors. The cost of painting each house with a certain color is different. You have to paint all the houses such that no two adjacent houses have the same color.

The cost of painting each house with a certain color is represented by a  $n \times k$  cost matrix. For example,  $\text{costs}[0][0]$  is the cost of painting house 0 with color 0;  $\text{costs}[1][2]$  is the cost of painting house 1 with color 2, and so on... Find the minimum cost to paint all houses.

Note:

All costs are positive integers.

Example:

Input:  $[[1,5,3],[2,9,4]]$

Output: 5

Explanation: Paint house 0 into color 0, paint house 1 into color 2. Minimum cost:  $1 + 4 = 5$ ;

Or paint house 0 into color 2, paint house 1 into color 0. Minimum cost:  $3 + 2 = 5$ .

Follow up:

Could you solve it in  $O(nk)$  runtime?

*Solution*

05/28/2020:

```
class Solution {
public:
    int minCostII(vector<vector<int>>& costs) {
        if (costs.empty() || costs[0].empty()) return 0;
        int m = costs.size(), n = costs[0].size();
        for (int i = 1; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                int prevMinCost = INT_MAX;
                for (int k = 0; k < n; ++k)
                    if (k != j)
                        prevMinCost = min(prevMinCost, costs[i - 1][k]);
                costs[i][j] += prevMinCost;
            }
        }
        return *min_element(costs.back().begin(), costs.back().end());
    }
};
```

## 267. Palindrome Permutation II

### Description

Given a string *s*, return all the palindromic permutations (without duplicates) of it. Return an empty list if no palindromic permutation could be form.

Example 1:

Input: "aabb"

Output: ["abba", "baab"]

Example 2:

Input: "abc"

Output: []

### Solution

05/28/2020:

```
class Solution {
public:
    unordered_map<char, int> cnt;
    vector<string> generatePalindromes(string s) {
        if (!isPalindromic(s)) return {};
        unordered_set<string> permutation;
        string prefix, oddChar;
        for (auto& [k, v] : cnt) {
            prefix += string(v / 2, k);
            if (v % 2 == 1) oddChar = string(1, k);
        }
        sort(prefix.begin(), prefix.end());
        do {
            string suffix(prefix);
            reverse(suffix.begin(), suffix.end());
            permutation.insert(prefix + oddChar + suffix);
        } while (next_permutation(prefix.begin(), prefix.end()));
        return vector<string>(permutation.begin(), permutation.end());
    }

    bool isPalindromic(const string& s) {
        int numOfOdds = 0;
        for (auto& c : s) ++cnt[c];
        for (auto& [k, v] : cnt)
            if (v % 2 == 1)
                if (++numOfOdds > 1)
                    return false;
    }
};
```

```
    return true;
}
};
```

## 268. Missing Number

### *Description*

Given an array containing  $n$  distinct numbers taken from  $0, 1, 2, \dots, n$ , find the one that is missing from the array.

Example 1:

Input: [3,0,1]

Output: 2

Example 2:

Input: [9,6,4,2,3,5,7,0,1]

Output: 8

Note:

Your algorithm should run in linear runtime complexity. Could you implement it using only constant extra space complexity?

### *Solution*

05/27/2020:

```
class Solution {
public:
    int missingNumber(vector<int>& nums) {
        sort(nums.begin(), nums.end());
        int n = nums.size();
        for (int i = 0; i < n; ++i) {
            if (nums[i] != i) {
                return i;
            }
        }
        return n;
    }
};
```

```

class Solution {
public:
    int missingNumber(vector<int>& nums) {
        unordered_set<int> seen;
        for (auto& i : nums) seen.insert(i);
        for (int i = 0; i <= (int)nums.size(); ++i)
            if (seen.count(i) == 0)
                return i;
        return -1;
    }
};

```

```

class Solution {
public:
    int missingNumber(vector<int>& nums) {
        int n = nums.size();
        bitset<200000000> s;
        for (auto& i : nums) s[i] = 1;
        for (int i = 0; i <= n; ++i)
            if (s[i] == 0)
                return i;
        return -1;
    }
};

```

## 270. Closest Binary Search Tree Value

### *Description*

Given a non-empty binary search tree and a target value, find the value in the BST that is closest to the target.

**Note:**

Given target value is a floating point.

You are guaranteed to have only one unique value in the BST that is closest to the target.

Example:

Input: root = [4,2,5,1,3], target = 3.714286

```

    4
   / \
  2   5

```

```
/ \  
1  3
```

Output: 4

*Solution*

04/28/2020:

```
/**  
 * Definition for a binary tree node.  
 * struct TreeNode {  
 *     int val;  
 *     TreeNode *left;  
 *     TreeNode *right;  
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}  
 * };  
 */  
class Solution {  
public:  
    int closestValue(TreeNode* root, double target) {  
        if (root == nullptr) return INT_MAX;  
        int n = root->val;  
        int n_left = closestValue(root->left, target);  
        int n_right = closestValue(root->right, target);  
        n = !root->left || fabs(n - target) < fabs(n_left - target) ? n : n_left;  
        n = !root->right || fabs(n - target) < fabs(n_right - target) ? n : n_right;  
        return n;  
    }  
};
```

```
class Solution {  
public:  
    int closestValue(TreeNode* root, double target) {  
        int n = root->val;  
        double d = fabs(root->val - target);  
        if (root->left && target < root->val) {  
            int n_left = closestValue(root->left, target);  
            double d_left = fabs(n_left - target);  
            if (d_left < d) {  
                n = n_left;  
                d = d_left;  
            }  
        }  
        if (root->right && target > root->val) {  
            int n_right = closestValue(root->right, target);  
            double d_right = fabs(n_right - target);  
            if (d_right < d) {  
                n = n_right;  
                d = d_right;  
            }  
        }  
        return n;  
    }  
};
```



```

        if (d_right < d) {
            n = n_right;
            d = d_right;
        }
    }
    return n;
}
};

```

```

class Solution {
public:
    int closestValue(TreeNode* root, double target) {
        int n = root->val;
        double d = fabs(root->val - target);
        if (root->left && target < root->val) {
            int n_left = closestValue(root->left, target);
            double d_left = fabs(n_left - target);
            if (d_left < d) {
                n = n_left;
                d = d_left;
            }
        }
        if (root->right && target > root->val) {
            int n_right = closestValue(root->right, target);
            double d_right = fabs(n_right - target);
            if (d_right < d) {
                n = n_right;
                d = d_right;
            }
        }
        return n;
    }
};

```

## 274. H-Index

*Description*

Given an array of citations (each citation is a non-negative integer) of a researcher, write a function to compute the researcher's h-index.

According to the definition of h-index on Wikipedia: "A scientist has index h if h of his/her N papers have at least h citations each, and the other N - h papers have no more than h citations each."

Example:

Input: citations = [3,0,6,1,5]

Output: 3

Explanation: [3,0,6,1,5] means the researcher has 5 papers in total and each of them had

received 3, 0, 6, 1, 5 citations respectively.

Since the researcher has 3 papers with at least 3 citations each and the remaining

two with no more than 3 citations each, her h-index is 3.

Note: If there are several possible values for h, the maximum one is taken as the h-index.

*Solution*

05/30/2020:

```
class Solution {
public:
    int hIndex(vector<int>& citations) {
        sort(citations.rbegin(), citations.rend());
        int lo = 0, hi = citations.size() - 1;
        while (lo <= hi) {
            int mid = lo + (hi - lo) / 2;
            if (citations[mid] >= mid + 1) {
                lo = mid + 1;
            } else {
                hi = mid - 1;
            }
        }
        return lo;
    }
};
```

## 276. Paint Fence

---

*Description*

There is a fence with  $n$  posts, each post can be painted with one of the  $k$  colors.

You have to paint all the posts such that no more than two adjacent fence posts have the same color.

Return the total number of ways you can paint the fence.

Note:

$n$  and  $k$  are non-negative integers.

Example:

Input:  $n = 3, k = 2$

Output: 6

Explanation: Take  $c_1$  as color 1,  $c_2$  as color 2. All possible ways are:

	post1	post2	post3
1	c1	c1	c2
2	c1	c2	c1
3	c1	c2	c2
4	c2	c1	c1
5	c2	c1	c2
6	c2	c2	c1

*Solution*

01/14/2020 (Dynamic Programming):

```
class Solution {
public:
    int numWays(int n, int k) {
        if (n == 0 || k == 0) return 0;
        vector<vector<int>> dp(n, vector<int>(2, 0));
        dp[0][0] = k;
        for (int i = 1; i < n; ++i) {
            dp[i][0] = (dp[i - 1][0] + dp[i - 1][1]) * (k - 1);
            dp[i][1] = dp[i - 1][0];
        }
        return dp.back()[0] + dp.back()[1];
    }
};
```

01/14/2020: (Dynamic Programming, Improve Space Complexity):

```
class Solution {
```

```

public:
    int numWays(int n, int k) {
        if (n == 0 || k == 0) return 0;
        vector<int> dp(2, 0);
        dp[0] = k;
        for (int i = 1; i < n; ++i) {
            int tmp = dp[0];
            dp[0] = (dp[0] + dp[1]) * (k - 1);
            dp[1] = tmp;
        }
        return dp[0] + dp[1];
    }
};

```

## 278. First Bad Version

### Description

You are a product manager and currently leading a team to develop a new product. Unfortunately, the latest version of your product fails the quality check. Since each version is developed based on the previous version, all the versions after a bad version are also bad.

Suppose you have  $n$  versions  $[1, 2, \dots, n]$  and you want to find out the first bad one, which causes all the following ones to be bad.

You are given an API `bool isBadVersion(version)` which will return whether version is bad. Implement a function to find the first bad version. You should minimize the number of calls to the API.

Example:

Given  $n = 5$ , and version = 4 is the first bad version.

```

call isBadVersion(3) -> false
call isBadVersion(5) -> true
call isBadVersion(4) -> true

```

Then 4 is the first bad version.

### Solution

04/23/2020:

```

// The API isBadVersion is defined for you.

```

```

// bool isBadVersion(int version);

class Solution {
public:
    int firstBadVersion(int n) {
        int lo = 1, hi = n;
        while (lo <= hi) {
            int mid = lo + (hi - lo) / 2;
            if (isBadVersion(mid) == true) {
                hi = mid - 1;
            } else {
                lo = mid + 1;
            }
        }
        return lo;
    }
};

```

## 285. Inorder Successor in BST

### Description

Given a binary search tree and a node in it, find the in-order successor of that node in the BST.

The successor of a node  $p$  is the node with the smallest key greater than  $p.val$ .

Example 1:

Input: root = [2,1,3], p = 1

Output: 2

Explanation: 1's in-order successor node is 2. Note that both  $p$  and the return value is of `TreeNode` type.

Example 2:

Input: root = [5,3,6,2,4,null,null,1], p = 6

Output: null

Explanation: There is no in-order successor of the current node, so the answer is null.

Note:

If the given node has no in-order successor in the tree, return null.  
It's guaranteed that the values of the tree are unique.

*Solution*

05/24/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
class Solution {
public:
    TreeNode* inorderSuccessor(TreeNode* root, TreeNode* p) {
        stack<TreeNode*> st;
        TreeNode* cur = root;
        while (cur || !st.empty()) {
            while (cur) {
                st.push(cur);
                cur = cur->left;
            }
            cur = st.top(); st.pop();
            if (cur->val > p->val) return cur;
            cur = cur->right;
        }
        return nullptr;
    }
};
```

```
class Solution {
public:
    TreeNode* inorderSuccessor(TreeNode* root, TreeNode* p) {
        vector<TreeNode*> inorder = inorderTraversal(root);
        for (auto& i : inorder) {
            if (i->val > p->val) {
                return i;
            }
        }
        return nullptr;
    }
};
```

```

vector<TreeNode*> inorderTraversal(TreeNode* root) {
    if (root == nullptr) return {};
    vector<TreeNode*> leftTraversal = inorderTraversal(root->left);
    vector<TreeNode*> rightTraversal = inorderTraversal(root->right);
    leftTraversal.push_back(root);
    leftTraversal.insert(leftTraversal.end(), rightTraversal.begin(),
rightTraversal.end());
    return leftTraversal;
}
};

```

## 286. Walls and Gates

### Description

You are given a  $m \times n$  2D grid initialized with these three possible values.

-1 - A wall or an obstacle.

0 - A gate.

INF - Infinity means an empty room. We use the value  $2^{31} - 1 = 2147483647$  to represent INF as you may assume that the distance to a gate is less than 2147483647.

Fill each empty room with the distance to its nearest gate. If it is impossible to reach a gate, it should be filled with INF.

Example:

Given the 2D grid:

```

INF  -1  0  INF
INF INF INF  -1
INF  -1 INF  -1
  0  -1 INF INF

```

After running your function, the 2D grid should be:

```

 3  -1  0  1
 2  2  1  -1
 1  -1  2  -1
 0  -1  3  4

```

### Solution

05/08/2020 (DFS):

```

class Solution {

```

```

public:
    void wallsAndGates(vector<vector<int>>& rooms) {
        if (rooms.empty() || rooms[0].empty()) return;
        for (int i = 0; i < rooms.size(); ++i) {
            for (int j = 0; j < rooms[0].size(); ++j) {
                if (rooms[i][j] == 0) {
                    dfs(rooms, i, j, rooms[i][j]);
                }
            }
        }
    }

    void dfs(vector<vector<int>>& rooms, int i, int j, int d) {
        int m = rooms.size(), n = rooms[0].size();
        if (!(i >= 0 && i < m && j >= 0 && j < n) || rooms[i][j] < d) return;
        rooms[i][j] = min(rooms[i][j], d);
        dfs(rooms, i - 1, j, d + 1);
        dfs(rooms, i + 1, j, d + 1);
        dfs(rooms, i, j - 1, d + 1);
        dfs(rooms, i, j + 1, d + 1);
    }
};

```

05/08/2020 (BFS):

```

class Solution {
public:
    void wallsAndGates(vector<vector<int>>& rooms) {

        int m = rooms.size(), n = rooms[0].size();
        queue<pair<int, int>> q;
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (rooms[i][j] == 0) {
                    q.emplace(i, j);
                }
            }
        }

        const int INF = numeric_limits<int>::max();
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        while (!q.empty()) {
            int sz = q.size();
            for (int k = 0; k < sz; ++k) {
                pair<int, int> cur = q.front(); q.pop();
                int i = cur.first, j = cur.second;
                for (int d = 0; d < 4; ++d) {

```



```

int ni = i + dir[d][0];
int nj = j + dir[d][1];
if (ni >= 0 && ni < m && nj >= 0 && nj < n && rooms[ni][nj] == INF) {
    rooms[ni][nj] = rooms[i][j] + 1;
    q.emplace(ni, nj);
}
}
}
}
};

```

## 287. Find the Duplicate Number

### Description

Given an array `nums` containing  $n + 1$  integers where each integer is between 1 and  $n$  (inclusive), prove that at least one duplicate number must exist. Assume that there is only one duplicate number, find the duplicate one.

Example 1:

Input: [1,3,4,2,2]

Output: 2

Example 2:

Input: [3,1,3,4,2]

Output: 3

Note:

You must not modify the array (assume the array is read only).

You must use only constant,  $O(1)$  extra space.

Your runtime complexity should be less than  $O(n^2)$ .

There is only one duplicate number in the array, but it could be repeated more than once.

### Solution

04/28/2020:

Use linear search on sorted array:

```

class Solution {
public:
    int findDuplicate(vector<int>& nums) {
        sort(nums.begin(), nums.end());
        for (int i = 1; i < (int)nums.size(); ++i) {
            if (nums[i] == nums[i - 1]) {
                return nums[i];
            }
        }
        return -1;
    }
};

```

Use binary search on sorted array:

```

class Solution {
public:
    int findDuplicate(vector<int>& nums) {
        sort(nums.begin(), nums.end());
        for (int i = 0; i < nums.size(); ++i) {
            auto p = equal_range(nums.begin() + i, nums.end(), nums[i]);
            if (p.second - p.first >= 2) return *p.first;
        }
        return -1;
    }
};

```

Use hash map:

```

class Solution {
public:
    int findDuplicate(vector<int>& nums) {
        unordered_map<int, int> mp;
        for (auto& n : nums) {
            if (++mp[n] > 1) {
                return n;
            }
        }
        return -1;
    }
};

```

## Description

Given a pattern and a string str, find if str follows the same pattern.

Here follow means a full match, such that there is a bijection between a letter in pattern and a non-empty word in str.

Example 1:

Input: pattern = "abba", str = "dog cat cat dog"

Output: true

Example 2:

Input: pattern = "abba", str = "dog cat cat fish"

Output: false

Example 3:

Input: pattern = "aaaa", str = "dog cat cat dog"

Output: false

Example 4:

Input: pattern = "abba", str = "dog dog dog dog"

Output: false

Notes:

You may assume pattern contains only lowercase letters, and str contains lowercase letters that may be separated by a single space.

## Solution

05/20/2020:

```
class Solution {
public:
    bool wordPattern(string pattern, string str) {
        istringstream iss(str);
        string s;
        vector<string> words;
        while (iss >> s) words.push_back(s);
        if (words.size() != pattern.size()) return false;
        unordered_map<char, unordered_set<string>> mp1;
        unordered_map<string, unordered_set<char>> mp2;
        for (int i = 0; i < (int)pattern.size(); ++i) {
            mp1[pattern[i]].insert(words[i]);
            mp2[words[i]].insert(pattern[i]);
        }
        for (auto& m : mp1) if (m.second.size() > 1) return false;
        for (auto& m : mp2) if (m.second.size() > 1) return false;
        return true;
    }
};
```

```
}  
};
```

## 299. Bulls and Cows

### Description

You are playing the following Bulls and Cows game with your friend: You write down a number and ask your friend to guess what the number is. Each time your friend makes a guess, you provide a hint that indicates how many digits in said guess match your secret number exactly in both digit and position (called "bulls") and how many digits match the secret number but locate in the wrong position (called "cows"). Your friend will use successive guesses and hints to eventually derive the secret number.

Write a function to return a hint according to the secret number and friend's guess, use A to indicate the bulls and B to indicate the cows.

Please note that both secret number and friend's guess may contain duplicate digits.

Example 1:

Input: secret = "1807", guess = "7810"

Output: "1A3B"

Explanation: 1 bull and 3 cows. The bull is 8, the cows are 0, 1 and 7.

Example 2:

Input: secret = "1123", guess = "0111"

Output: "1A1B"

Explanation: The 1st 1 in friend's guess is a bull, the 2nd or 3rd 1 is a cow.

Note: You may assume that the secret number and your friend's guess only contain digits, and their lengths are always equal.

### Solution

02/05/2020:

```
class Solution {  
public:  
    string getHint(string secret, string guess) {
```

```

vector<int> arr1(10, 0), arr2(10, 0);
for (auto& n : secret) {
    ++arr1[n - '0'];
}
for (auto& n : guess) {
    ++arr2[n - '0'];
}
int match = 0, bull = 0;
for (int i = 0; i < 10; ++i) {
    match += min(arr1[i], arr2[i]);
}
for (int i = 0; i < (int)secret.size(); ++i) {
    if (secret[i] == guess[i]) {
        ++bull;
    }
}
string ret;
ret += to_string(bull) + "A" + to_string(match - bull) + "B";
return ret;
}
};

```

## 303. Range Sum Query - Immutable

### *Description*

Given an integer array `nums`, find the sum of the elements between indices `i` and `j` ( $i \leq j$ ), inclusive.

Example:

Given `nums = [-2, 0, 3, -5, 2, -1]`

`sumRange(0, 2) -> 1`

`sumRange(2, 5) -> -1`

`sumRange(0, 5) -> -3`

Note:

You may assume that the array does not change.

There are many calls to `sumRange` function.

### *Solution*

01/14/2020 (Dynamic Programming):

```

class NumArray {
public:
    vector<int> nums;
    NumArray(vector<int>& nums) {
        for (int i = 1; i < nums.size(); ++i) nums[i] += nums[i - 1];
        this->nums = nums;
    }

    int sumRange(int i, int j) {
        return i > 0 ? nums[j] - nums[i - 1] : nums[j];
    }
};

/**
 * Your NumArray object will be instantiated and called as such:
 * NumArray* obj = new NumArray(nums);
 * int param_1 = obj->sumRange(i,j);
 */

```

## 305. Number of Islands II

### *Description*

A 2d grid map of  $m$  rows and  $n$  columns is initially filled with water. We may perform an `addLand` operation which turns the water at position  $(row, col)$  into a land. Given a list of positions to operate, count the number of islands after each `addLand` operation. An island is surrounded by water and is formed by connecting adjacent lands horizontally or vertically. You may assume all four edges of the grid are all surrounded by water.

Example:

Input:  $m = 3, n = 3, positions = [[0,0], [0,1], [1,2], [2,1]]$

Output:  $[1,1,2,3]$

Explanation:

Initially, the 2d grid `grid` is filled with water. (Assume 0 represents water and 1 represents land).

```

0 0 0
0 0 0
0 0 0

```

Operation #1: `addLand(0, 0)` turns the water at `grid[0][0]` into a land.

```

1 0 0

```

```
0 0 0   Number of islands = 1
```

```
0 0 0
```

Operation #2: addLand(0, 1) turns the water at grid[0][1] into a land.

```
1 1 0
```

```
0 0 0   Number of islands = 1
```

```
0 0 0
```

Operation #3: addLand(1, 2) turns the water at grid[1][2] into a land.

```
1 1 0
```

```
0 0 1   Number of islands = 2
```

```
0 0 0
```

Operation #4: addLand(2, 1) turns the water at grid[2][1] into a land.

```
1 1 0
```

```
0 0 1   Number of islands = 3
```

```
0 1 0
```

Follow up:

Can you do it in time complexity  $O(k \log mn)$ , where  $k$  is the length of the positions?

*Solution*

05/08/2020 (Union-Find) [Discussion](#):

1. Use a set `islands` to store all the islands being added so far.
2. Use a set `components` to keep track of roots of current connected components (the size of components is the desired number of islands).
3. Whenever a new island is added, check whether the neighboring islands (in all four directions) exist. If there is a such neighbor, remove its root from the components and merge current island with that neighbor.
4. Then insert the root of island into components.
5. Store the size of components.

```
class UnionFind {
private:
    vector<int> id;
    vector<int> sz;

public:
    UnionFind(int n) {
        id.resize(n);
        iota(id.begin(), id.end(), 0);
        sz.resize(n, 1);
    }

    int find(int x) {
```

```

while(x != id[x]) {
    id[x] = id[id[x]];
    x = id[x];
}
return x;
}

bool merge(int x, int y) {
    int i = find(x), j = find(y);
    if (i == j) return false;
    if (sz[i] > sz[j]) {
        sz[i] += sz[j];
        id[j] = i;
    } else {
        sz[j] += sz[i];
        id[i] = j;
    }
    return true;
}
};

class Solution {
public:
    vector<int> numIslands2(int m, int n, vector<vector<int>>& positions) {
        if (m == 0 || n == 0 || positions.empty() || positions[0].empty()) return
{};
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        UnionFind uf(m * n);
        unordered_set<int> islands;
        unordered_set<int> components;
        vector<int> ret;
        for (auto& p : positions) {
            int i = p[0], j = p[1], island = i * n + j;
            islands.insert(island);
            for (int d = 0; d < 4; ++d) {
                int ni = i + dir[d][0], nj = j + dir[d][1], neighbor = ni * n + nj;
                if (ni >= 0 && ni < m && nj >= 0 && nj < n && islands.count(neighbor) >
0) {
                    if (components.count(uf.find(neighbor))) {
                        components.erase(uf.find(neighbor));
                    }
                    uf.merge(island, neighbor);
                }
            }
            components.insert(uf.find(island));
            ret.push_back(components.size());
        }
        return ret;
    }
}

```



```
};
```

## 311. Sparse Matrix Multiplication

### Description

Given two sparse matrices A and B, return the result of AB.

You may assume that A's column number is equal to B's row number.

Example:

Input:

```
A = [  
  [ 1, 0, 0 ],  
  [-1, 0, 3]  
]
```

```
B = [  
  [ 7, 0, 0 ],  
  [ 0, 0, 0 ],  
  [ 0, 0, 1 ]  
]
```

Output:

$$AB = \begin{bmatrix} 1 & 0 & 0 \\ -1 & 0 & 3 \end{bmatrix} \times \begin{bmatrix} 7 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 7 & 0 & 0 \\ -7 & 0 & 3 \end{bmatrix}$$

### Solution

05/20/2020:

```
class Solution {  
public:  
    vector<vector<int>> multiply(vector<vector<int>>& A, vector<vector<int>>& B) {  
        if (A.empty() || A[0].empty()) return {};  
        int nrowA = A.size(), ncolA = A[0].size(), ncolB = B[0].size();  
        vector<vector<int>> ret(nrowA, vector<int>(ncolB, 0));  
        for (int i = 0; i < nrowA; ++i) {  
            for (int k = 0; k < ncolA; ++k) {  
                if (A[i][k] == 0) continue;  
                for (int j = 0; j < ncolB; ++j) {
```

```

        if (B[k][j] == 0) continue;
        ret[i][j] += A[i][k] * B[k][j];
    }
}
return ret;
};

```

```

class Solution {
public:
    vector<vector<int>> multiply(vector<vector<int>>& A, vector<vector<int>>& B) {
        unordered_map<int, unordered_set<int>> matA, matB;
        int nrowA = A.size(), ncolA = A[0].size(), ncolB = B[0].size();
        for (int i = 0; i < nrowA; ++i)
            for (int j = 0; j < ncolA; ++j)
                if (A[i][j] != 0) matA[i].insert(j);
        for (int i = 0; i < ncolA; ++i)
            for (int j = 0; j < ncolB; ++j)
                if (B[i][j] != 0) matB[i].insert(j);
        vector<vector<int>> ret(nrowA, vector<int>(ncolB, 0));
        for (auto& mA : matA) {
            int i = mA.first;
            for (auto& k : mA.second)
                if (matB.count(k) > 0)
                    for (auto& j : matB[k])
                        ret[i][j] += A[i][k] * B[k][j];
        }
        return ret;
    }
};

```

## 314. Binary Tree Vertical Order Traversal

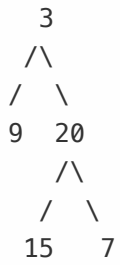
### *Description*

Given a binary tree, return the vertical order traversal of its nodes' values. (ie, from top to bottom, column by column).

If two nodes are in the same row and column, the order should be from left to right.

Examples 1:

Input: [3,9,20,null,null,15,7]



Output:

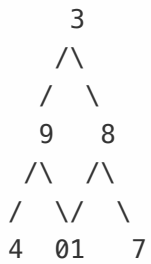
```

[
  [9],
  [3,15],
  [20],
  [7]
]

```

Examples 2:

Input: [3,9,8,4,0,1,7]



Output:

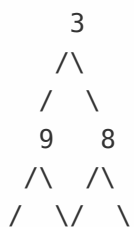
```

[
  [4],
  [9],
  [3,0,1],
  [8],
  [7]
]

```

Examples 3:

Input: [3,9,8,4,0,1,7,null,null,null,2,5] (0's right child is 2 and 1's left child is 5)



```

4  01  7
  ^
 /  \
5    2

```

Output:

```

[
  [4],
  [9,5],
  [3,0,1],
  [8,2],
  [7]
]

```

*Solution*

04/26/2020:

```

/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
class Solution {
public:
    vector<vector<int>> verticalOrder(TreeNode* root) {
        map<int, vector<int>> mp;
        queue<pair<TreeNode*, int>> q;
        q.emplace(root, 0);
        while (!q.empty()) {
            int n = q.size();
            for (int i = 0; i < n; ++i) {
                pair<TreeNode*, int> cur = q.front(); q.pop();
                mp[cur.second].push_back(cur.first->val);
                if (cur.first->left) q.emplace(cur.first->left, cur.second - 1);
                if (cur.first->right) q.emplace(cur.first->right, cur.second + 1);
            }
        }
        vector<vector<int>> ret;
        for (auto& m : mp) ret.push_back(m.second);
        return ret;
    }
};

```

Using unordered\_map:

```
class Solution {
public:
    vector<vector<int>> verticalOrder(TreeNode* root) {
        if (root == nullptr) return {};
        int minCol = 0, maxCol = 0;
        unordered_map<int, vector<int>> mp;
        queue<pair<TreeNode*, int>> q;
        q.emplace(root, 0);
        while (!q.empty()) {
            int n = q.size();
            for (int i = 0; i < n; ++i) {
                pair<TreeNode*, int> cur = q.front(); q.pop();
                mp[cur.second].push_back(cur.first->val);
                minCol = min(minCol, cur.second);
                maxCol = max(maxCol, cur.second);
                if (cur.first->left) q.emplace(cur.first->left, cur.second - 1);
                if (cur.first->right) q.emplace(cur.first->right, cur.second + 1);
            }
        }
        vector<vector<int>> ret;
        for (int i = minCol; i <= maxCol; ++i)
            if (mp.count(i) > 0)
                ret.push_back(mp[i]);
        return ret;
    }
};
```

## 322. Coin Change

---

*Description*

You are given coins of different denominations and a total amount of money amount. Write a function to compute the fewest number of coins that you need to make up that amount. If that amount of money cannot be made up by any combination of the coins, return -1.

Example 1:

Input: coins = [1, 2, 5], amount = 11

Output: 3

Explanation: 11 = 5 + 5 + 1

Example 2:

Input: coins = [2], amount = 3

Output: -1

Note:

You may assume that you have an infinite number of each kind of coin.

*Solution*

05/27/2020:

```
class Solution {
public:
    int coinChange(vector<int>& coins, int amount) {
        vector<int> dp(amount + 1, INT_MAX);
        dp[0] = 0;
        for (int i = 1; i <= amount; ++i)
            for (auto& c : coins)
                if (i - c >= 0 && dp[i - c] != INT_MAX)
                    dp[i] = min(dp[i - c] + 1, dp[i]);
        return dp.back() == INT_MAX ? -1 : dp.back();
    }
};
```

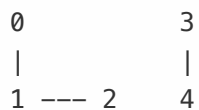
## 323. Number of Connected Components in an Undirected Graph

*Description*

Given  $n$  nodes labeled from 0 to  $n - 1$  and a list of undirected edges (each edge is a pair of nodes), write a function to find the number of connected components in an undirected graph.

Example 1:

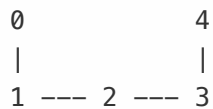
Input: n = 5 and edges = [[0, 1], [1, 2], [3, 4]]



Output: 2

Example 2:

Input: n = 5 and edges = [[0, 1], [1, 2], [2, 3], [3, 4]]



Output: 1

Note:

You can assume that no duplicate edges will appear in edges. Since all edges are undirected, [0, 1] is the same as [1, 0] and thus will not appear together in edges.

*Solution*

05/05/2020 (BFS):

```
class Solution {
public:
    int countComponents(int n, vector<vector<int>>& edges) {
        if (n == 0) return 0;
        unordered_map<int, unordered_set<int>> next;
        for (auto& e : edges) {
            next[e[0]].insert(e[1]);
            next[e[1]].insert(e[0]);
        }
        unordered_set<int> visited;
        int ret = 0;
        for (int i = 0; i < n; ++i) {
            if (visited.count(i) > 0) continue;
            queue<int> q;
            q.push(i);
            visited.insert(i);
            while (!q.empty()) {
                int cur = q.front(); q.pop();
                for (auto& c : next[cur]) {
                    if (visited.count(c) == 0) q.push(c);
                    visited.insert(c);
                }
            }
            ret++;
        }
        return ret;
    }
};
```

```

    }
    ++ret;
}
return ret;
}
};

```

03/05/2020 (Union-Find Set):

```

class UnionFind {
private:
    vector<int> id;
    vector<int> sz;

public:
    UnionFind(int n) {
        id.resize(n);
        iota(id.begin(), id.end(), 0);
        sz.resize(n, 1);
    }

    int find(int x) {
        while (x != id[x]) {
            id[x] = id[id[x]];
            x = id[x];
        }
        return x;
        // if (x == id[x]) return x;
        // return id[x] = find(id[x]);
    }

    bool connected(int x, int y) {
        return find(x) == find(y);
    }

    void merge(int x, int y) {
        int i = find(x), j = find(y);
        if (i == j) return;
        if (sz[i] > sz[j]) {
            id[j] = i;
            sz[i] += sz[j];
        } else {
            id[i] = j;
            sz[j] = sz[i];
        }
    }
};

```



```

class Solution {
public:
    int countComponents(int n, vector<vector<int>>& edges) {
        UnionFind uf(n);
        for (auto& e : edges) {
            uf.merge(e[0], e[1]);
        }

        unordered_set<int> roots;
        for (int i = 0; i < n; ++i) {
            roots.insert(uf.find(i));
        }
        return roots.size();
    }
};

```

## 326. Power of Three

### Description

Given an integer, write a function to determine if it is a power of three.

Example 1:

Input: 27  
Output: true

Example 2:

Input: 0  
Output: false

Example 3:

Input: 9  
Output: true

Example 4:

Input: 45  
Output: false

Follow up:

Could you do it without using any loop / recursion?

### Solution

06/10/2020:

```
class Solution {
public:
    bool isPowerOfThree(int n) {
        return n > 0 && (1162261467 % n == 0);
    }
};
```

```
class Solution {
public:
    bool isPowerOfThree(int n) {
        if (n <= 0) return false;
        for (; n != 1 && n % 3 == 0; n /= 3);
        return n == 1;
    }
};
```

## 328. Odd Even Linked List

### *Description*

Given a singly linked list, group all odd nodes together followed by the even nodes. Please note here we are talking about the node number and not the value in the nodes.

You should try to do it in place. The program should run in  $O(1)$  space complexity and  $O(\text{nodes})$  time complexity.

Example 1:

Input: 1->2->3->4->5->NULL

Output: 1->3->5->2->4->NULL

Example 2:

Input: 2->1->3->5->6->4->7->NULL

Output: 2->3->6->7->1->5->4->NULL

Note:

The relative order inside both the even and odd groups should remain as it was in the input.

The first node is considered odd, the second node even and so on ...

### *Solution*

05/14/2020:

```

/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
public:
    ListNode* oddEvenList(ListNode* head) {
        if (!head || !head->next) return head;
        ListNode *oddHead = head, *evenHead = head->next;
        ListNode *cur = evenHead, *oddCur = oddHead, *evenCur = evenHead;
        while (cur->next) {
            if (cur) {
                oddCur->next = cur->next;
                oddCur = oddCur->next;
                cur = cur->next;
            }
            if (cur) {
                evenCur->next = cur->next;
                evenCur = evenCur->next;
                cur = cur->next;
            }
            if (!cur) break;
        }
        oddCur->next = evenHead;
        return oddHead;
    }
};

```

```

/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
public:
    ListNode* oddEvenList(ListNode* head) {

```

```

if (!head || !head->next) return head;
ListNode *evenHead = head->next, *evenCur = evenHead, *cur = head;
while (cur->next && evenCur->next) {
    cur->next = cur->next->next;
    cur = cur->next;
    evenCur->next = evenCur->next->next;
    evenCur = cur->next;
}
cur->next = evenHead;
return head;
}
};

```

## 344. Reverse String

### Description

Write a function that reverses a string. The input string is given as an array of characters `char[]`.

Do not allocate extra space for another array, you must do this by modifying the input array in-place with  $O(1)$  extra memory.

You may assume all the characters consist of printable ascii characters.

Example 1:

Input: ["h","e","l","l","o"]

Output: ["o","l","l","e","h"]

Example 2:

Input: ["H","a","n","n","a","h"]

Output: ["h","a","n","n","a","H"]

### Solution

05/30/2020:

```

class Solution {
public:
    void reverseString(vector<char>& s) {
        for (int lo = 0, hi = s.size() - 1; lo < hi; ++lo, --hi) swap(s[lo], s[hi]);
    }
};

```

## 345. Reverse Vowels of a String

### Description

Write a function that takes a string as input and reverse only the vowels of a string.

Example 1:

Input: "hello"

Output: "holle"

Example 2:

Input: "leetcode"

Output: "leotcede"

Note:

The vowels does not include the letter "y".

### Solution

06/16/2020:

```

class Solution {
public:
    string reverseVowels(string s) {
        unordered_set<char> vowels{'a', 'e', 'i', 'o', 'u'};
        int n = s.size(), l = 0, r = n - 1;
        while (l < r) {
            while (!isalpha(s[l]) && l < n - 1) ++l;
            while (!isalpha(s[r]) && r > 0) --r;
            while (l < n - 1 && vowels.count(tolower(s[l])) == 0) ++l;
            while (r > 0 && vowels.count(tolower(s[r])) == 0) --r;
            if (l < r)
                if (vowels.count(tolower(s[l])) > 0 && vowels.count(tolower(s[r])) > 0)
                    swap(s[l++], s[r--]);
        }
        return s;
    }
};

```

```
};
```

## 347. Top K Frequent Elements

### Description

Given a non-empty array of integers, return the k most frequent elements.

Example 1:

Input: nums = [1,1,1,2,2,3], k = 2

Output: [1,2]

Example 2:

Input: nums = [1], k = 1

Output: [1]

Note:

You may assume k is always valid,  $1 \leq k \leq$  number of unique elements.

Your algorithm's time complexity must be better than  $O(n \log n)$ , where n is the array's size.

It's guaranteed that the answer is unique, in other words the set of the top k frequent elements is unique.

You can return the answer in any order.

### Solution

05/20/2020:

```
class Solution {
public:
    vector<int> topKFrequent(vector<int>& nums, int k) {
        unordered_map<int, int> cnt;
        for (auto& n : nums) ++cnt[n];
        priority_queue<pair<int, int>, vector<pair<int, int>>, less<pair<int, int>>>
q;
        for (auto m : cnt) q.emplace(m.second, m.first);
        vector<int> ret;
        while (k-- > 0) {
            ret.push_back(q.top().second);
            q.pop();
        }
        return ret;
    }
};
```

```

class Solution {
public:
    vector<int> topKFrequent(vector<int>& nums, int k) {
        unordered_map<int, int> cnt;
        for (auto& n : nums) ++cnt[n];
        vector<pair<int, int>> freq;
        for (auto& m : cnt) freq.emplace_back(m.second, m.first);
        sort(freq.begin(), freq.end(), greater<pair<int, int>>());
        vector<int> ret;
        for (auto& f : freq) {
            if (k-- > 0) {
                ret.push_back(f.second);
            } else {
                break;
            }
        }
        return ret;
    }
};

```

## 348. Design Tic-Tac-Toe

### Description

Design a Tic-tac-toe game that is played between two players on a  $n \times n$  grid.

You may assume the following rules:

A move is guaranteed to be valid and is placed on an empty block.

Once a winning condition is reached, no more moves is allowed.

A player who succeeds in placing  $n$  of their marks in a horizontal, vertical, or diagonal row wins the game.

Example:

Given  $n = 3$ , assume that player 1 is "X" and player 2 is "O" in the board.

```
TicTacToe toe = new TicTacToe(3);
```

```
toe.move(0, 0, 1); -> Returns 0 (no one wins)
```

```
|X| | |
| | | | // Player 1 makes a move at (0, 0).
| | | |
```

```
toe.move(0, 2, 2); -> Returns 0 (no one wins)
```

```
|X| |O|
| | | | // Player 2 makes a move at (0, 2).
```

```

| | | |
toe.move(2, 2, 1); -> Returns 0 (no one wins)
|X| |0|
| | | | // Player 1 makes a move at (2, 2).
| | |X|

toe.move(1, 1, 2); -> Returns 0 (no one wins)
|X| |0|
| |0| | // Player 2 makes a move at (1, 1).
| | |X|

toe.move(2, 0, 1); -> Returns 0 (no one wins)
|X| |0|
| |0| | // Player 1 makes a move at (2, 0).
|X| |X|

toe.move(1, 0, 2); -> Returns 0 (no one wins)
|X| |0|
|0|0| | // Player 2 makes a move at (1, 0).
|X| |X|

toe.move(2, 1, 1); -> Returns 1 (player 1 wins)
|X| |0|
|0|0| | // Player 1 makes a move at (2, 1).
|X|X|X|
Follow up:
Could you do better than O(n2) per move() operation?

```

*Solution*

06/15/2020:

```

class TicTacToe {
    vector<vector<int>> board;
    vector<int> rows;
    vector<int> cols;
    int diag;
    int antiDiag;
    int n;

public:
    /** Initialize your data structure here. */
    TicTacToe(int n) {
        board.assign(n, vector<int>(n, -1));
        rows.assign(n, 0);
        cols.assign(n, 0);
        diag = 0;
    }
};

```



```

    antiDiag = 0;
    this->n = n;
}

/** Player {player} makes a move at ({row}, {col}).
    @param row The row of the board.
    @param col The column of the board.
    @param player The player, can be either 1 or 2.
    @return The current winning condition, can be either:
            0: No one wins.
            1: Player 1 wins.
            2: Player 2 wins. */
int move(int row, int col, int player) {
    if (board[row][col] != -1) return 0;
    board[row][col] = player;
    if (++rows[row] == n) {
        int countRow = 0;
        for (int i = 0; i < n && board[row][i] == player; ++i) ++countRow;
        if (countRow == n) return player;
    }
    if (++cols[col]) {
        int countCol = 0;
        for (int i = 0; i < n && board[i][col] == player; ++i) ++countCol;
        if (countCol == n) return player;
    }
    if (row == col && ++diag == n) {
        int countDiag = 0;
        for (int i = 0; i < n && board[i][i] == player; ++i) ++countDiag;
        if (countDiag == n) return player;
    }
    if (row == n - 1 - col && ++antiDiag == n) {
        int countAntiDiag = 0;
        for (int i = 0; i < n && board[i][n - i - 1] == player; ++i)
++countAntiDiag;
        if (countAntiDiag == n) return player;
    }
    return 0;
}
};

/**
 * Your TicTacToe object will be instantiated and called as such:
 * TicTacToe* obj = new TicTacToe(n);
 * int param_1 = obj->move(row,col,player);
 */

```

## 367. Valid Perfect Square

### Description

Given a positive integer num, write a function which returns True if num is a perfect square else False.

Note: Do not use any built-in library function such as sqrt.

Example 1:

Input: 16  
Output: true

Example 2:

Input: 14  
Output: false

### Solution

05/08/2020:

```
class Solution {
public:
    bool isPerfectSquare(int num) {
        int lo = 0, hi = num;
        while (lo <= hi) {
            int mid = lo + (hi - lo) / 2;
            long long s = (long long)mid * mid;
            if (s == num) {
                return true;
            } else if (s > num) {
                hi = mid - 1;
            } else {
                lo = mid + 1;
            }
        }
        return (long long)hi * hi == num;
    }
};
```

## 368. Largest Divisible Subset

### Description

Given a set of distinct positive integers, find the largest subset such that every pair  $(S_i, S_j)$  of elements in this subset satisfies:

$S_i \% S_j = 0$  or  $S_j \% S_i = 0$ .

If there are multiple solutions, return any subset is fine.

Example 1:

Input: [1,2,3]

Output: [1,2] (of course, [1,3] will also be ok)

Example 2:

Input: [1,2,4,8]

Output: [1,2,4,8]

*Solution*

06/13/2020:

Dynamic Programming:

```
class Solution {
public:
    vector<int> largestDivisibleSubset(vector<int>& nums) {
        if (nums.empty()) return nums;
        sort(nums.begin(), nums.end());
        unordered_map<int, int> dp;
        unordered_map<int, int> pre;
        for (auto& n : nums) {
            dp[n] = 1;
            pre[n] = n;
        }
        int n = nums.size(), max_count = 1, max_num = nums[0];
        for (int i = 0; i < n; ++i) {
            for (int j = 0; j < i; ++j) {
                if (nums[i] % nums[j] == 0) {
                    if (dp[nums[j]] + 1 > dp[nums[i]]) {
                        pre[nums[i]] = nums[j];
                        dp[nums[i]] = dp[nums[j]] + 1;
                        if (dp[nums[i]] > max_count) {
                            max_count = dp[nums[i]];
                            max_num = nums[i];
                        }
                    }
                }
            }
        }
        return {max_num};
    }
};
```

```

vector<int> ret(1, max_num);
while (max_num != pre[max_num]) {
    max_num = pre[max_num];
    ret.push_back(max_num);
}
return ret;
}
};

```

Backtracking (TLE):

```

class Solution {
public:
    vector<int> ret;
    vector<int> largestDivisibleSubset(vector<int>& nums) {
        sort(nums.begin(), nums.end());
        int n = nums.size();
        vector<int> s;
        backtrack(0, n, s, nums);
        return ret;
    }

    void backtrack(int k, int n, vector<int>& s, vector<int>& nums) {
        if (k == n) {
            if (s.size() > ret.size()) ret = s;
            return;
        }
        bool ok = true;
        for (int i = s.size() - 1; i >= 0 && ok; --i)
            ok = nums[k] % s[i] == 0;
        if (ok) {
            s.push_back(nums[k]);
            backtrack(k + 1, n, s, nums);
            s.pop_back();
        }
        backtrack(k + 1, n, s, nums);
    }
};

```

## 374. Guess Number Higher or Lower

### *Description*

We are playing the Guess Game. The game is as follows:

I pick a number from 1 to n. You have to guess which number I picked.

Every time you guess wrong, I'll tell you whether the number is higher or lower.

You call a pre-defined API `guess(int num)` which returns 3 possible results (-1, 1, or 0):

-1 : My number is lower  
1 : My number is higher  
0 : Congrats! You got it!

Example :

Input: n = 10, pick = 6

Output: 6

*Solution*

04/23/2020:

```
/**
 * Forward declaration of guess API.
 * @param num    your guess
 * @return      -1 if num is lower than the guess number
 *              1 if num is higher than the guess number
 *              otherwise return 0
 * int guess(int num);
 */

class Solution {
public:
    int guessNumber(int n) {
        int lo = 1, hi = n;
        while (lo <= hi) {
            int mid = lo + (hi - lo) / 2;
            int ans = guess(mid);
            if (ans == 0) {
                return mid;
            } else if (ans == 1) {
                lo = mid + 1;
            } else {
                hi = mid - 1;
            }
        }
        return lo;
    }
};
```

## 379. Design Phone Directory

### Description

Design a Phone Directory which supports the following operations:

get: Provide a number which is not assigned to anyone.  
check: Check if a number is available or not.  
release: Recycle or release a number.

Example:

```
// Init a phone directory containing a total of 3 numbers: 0, 1, and 2.
PhoneDirectory directory = new PhoneDirectory(3);

// It can return any available phone number. Here we assume it returns 0.
directory.get();

// Assume it returns 1.
directory.get();

// The number 2 is available, so return true.
directory.check(2);

// It returns 2, the only number that is left.
directory.get();

// The number 2 is no longer available, so return false.
directory.check(2);

// Release number 2 back to the pool.
directory.release(2);

// Number 2 is available again, return true.
directory.check(2);
```

Constraints:

```
1 <= maxNumbers <= 10^4
0 <= number < maxNumbers
The total number of call of the methods is between [0 - 20000]
```

### Solution

06/15/2020:

```
class PhoneDirectory {
private:
    vector<bool> numbers;
    int p, maxNumbers;
    stack<int> st;

public:
    /** Initialize your data structure here
        @param maxNumbers - The maximum numbers that can be stored in the phone
        directory. */
    PhoneDirectory(int maxNumbers) {
        numbers.assign(maxNumbers, true);
        this->maxNumbers = maxNumbers;
        p = 0;
    }

    /** Provide a number which is not assigned to anyone.
        @return - Return an available number. Return -1 if none is available. */
    int get() {
        if (!st.empty()) {
            int cur = st.top(); st.pop();
            numbers[cur] = false;
            return cur;
        }
        while (!numbers[p] && p < maxNumbers) p++;
        if (p < maxNumbers) {
            numbers[p] = false;
            return p;
        } else {
            return -1;
        }
    }

    /** Check if a number is available or not. */
    bool check(int number) {
        return numbers[number];
    }

    /** Recycle or release a number. */
    void release(int number) {
        if (!numbers[number]) {
            numbers[number] = true;
            st.push(number);
        }
    }
};
```

```
/**
 * Your PhoneDirectory object will be instantiated and called as such:
 * PhoneDirectory* obj = new PhoneDirectory(maxNumbers);
 * int param_1 = obj->get();
 * bool param_2 = obj->check(number);
 * obj->release(number);
 */
```

```
class PhoneDirectory {
private:
    vector<bool> numbers;
    int p, maxNumbers;

public:
    /** Initialize your data structure here
        @param maxNumbers - The maximum numbers that can be stored in the phone
        directory. */
    PhoneDirectory(int maxNumbers) {
        numbers.assign(maxNumbers, true);
        this->maxNumbers = maxNumbers;
        p = 0;
    }

    /** Provide a number which is not assigned to anyone.
        @return - Return an available number. Return -1 if none is available. */
    int get() {
        while (!numbers[p] && p < maxNumbers) p++;
        if (p < maxNumbers) {
            numbers[p] = false;
            return p;
        } else {
            return -1;
        }
    }

    /** Check if a number is available or not. */
    bool check(int number) {
        return numbers[number];
    }

    /** Recycle or release a number. */
    void release(int number) {
        numbers[number] = true;
        p = min(number, p);
    }
};
```



```
/**
 * Your PhoneDirectory object will be instantiated and called as such:
 * PhoneDirectory* obj = new PhoneDirectory(maxNumbers);
 * int param_1 = obj->get();
 * bool param_2 = obj->check(number);
 * obj->release(number);
 */
```

## 380. Insert Delete GetRandom O(1)

### *Description*

Design a data structure that supports all following operations in average  $O(1)$  time.

`insert(val)`: Inserts an item `val` to the set if not already present.

`remove(val)`: Removes an item `val` from the set if present.

`getRandom`: Returns a random element from current set of elements. Each element must have the same probability of being returned.

Example:

```
// Init an empty set.
```

```
RandomizedSet randomSet = new RandomizedSet();
```

```
// Inserts 1 to the set. Returns true as 1 was inserted successfully.
```

```
randomSet.insert(1);
```

```
// Returns false as 2 does not exist in the set.
```

```
randomSet.remove(2);
```

```
// Inserts 2 to the set, returns true. Set now contains [1,2].
```

```
randomSet.insert(2);
```

```
// getRandom should return either 1 or 2 randomly.
```

```
randomSet.getRandom();
```

```
// Removes 1 from the set, returns true. Set now contains [2].
```

```
randomSet.remove(1);
```

```
// 2 was already in the set, so return false.
```

```
randomSet.insert(2);
```

```
// Since 2 is the only number in the set, getRandom always return 2.
```

```
randomSet.getRandom();
```

06/12/2020:

```
class RandomizedSet {
private:
    unordered_map<int, int> cnt;
    vector<int> nums;
    int n;

public:
    /** Initialize your data structure here. */
    RandomizedSet() {
        nums.clear();
        cnt.clear();
        n = 0;
    }

    /** Inserts a value to the set. Returns true if the set did not already
    contain the specified element. */
    bool insert(int val) {
        if (cnt.count(val) > 0) return false;
        nums.push_back(val);
        cnt[val] = n++;
        return true;
    }

    /** Removes a value from the set. Returns true if the set contained the
    specified element. */
    bool remove(int val) {
        if (cnt.count(val) == 0) return false;
        cnt[nums[n - 1]] = cnt[val];
        swap(nums[cnt[val]], nums[n - 1]);
        cnt.erase(val);
        nums.pop_back();
        --n;
        return true;
    }

    /** Get a random element from the set. */
    int getRandom() {
        if (n == 0) return -1;
        return nums[rand() % n];
    }
};

/**
 * Your RandomizedSet object will be instantiated and called as such:
 * RandomizedSet* obj = new RandomizedSet();
 */
```

```
* bool param_1 = obj->insert(val);
* bool param_2 = obj->remove(val);
* int param_3 = obj->getRandom();
*/
```

```
class RandomizedSet {
private:
    unordered_set<int> cnt;
    vector<int> nums;
    bool ok = false;

public:
    /** Initialize your data structure here. */
    RandomizedSet() {
        nums.clear();
        cnt.clear();
    }

    /** Inserts a value to the set. Returns true if the set did not already
    contain the specified element. */
    bool insert(int val) {
        if (cnt.count(val) > 0) return false;
        cnt.insert(val);
        nums.push_back(val);
        return true;
    }

    /** Removes a value from the set. Returns true if the set contained the
    specified element. */
    bool remove(int val) {
        if (cnt.count(val) == 0) return false;
        cnt.erase(val);
        return true;
    }

    /** Get a random element from the set. */
    int getRandom() {
        if (nums.empty()) return -1;
        int i;
        for (i = rand() % nums.size(); nums.size() != 0 && cnt.count(nums[i]) == 0;
i = nums.empty() ? 0 : rand() % nums.size())
            nums.erase(nums.begin() + i);
        return nums.empty() ? -1 : nums[i];
    }
};

/**
 * Your RandomizedSet object will be instantiated and called as such:
```

```
* RandomizedSet* obj = new RandomizedSet();
* bool param_1 = obj->insert(val);
* bool param_2 = obj->remove(val);
* int param_3 = obj->getRandom();
*/
```

## 384. Shuffle an Array

### *Description*

Shuffle a set of numbers without duplicates.

Example:

```
// Init an array with set 1, 2, and 3.
int[] nums = {1,2,3};
Solution solution = new Solution(nums);

// Shuffle the array [1,2,3] and return its result. Any permutation of [1,2,3]
must equally likely to be returned.
solution.shuffle();

// Resets the array back to its original configuration [1,2,3].
solution.reset();

// Returns the random shuffling of array [1,2,3].
solution.shuffle();
```

### *Solution*

05/25/2020:

```
class Solution {
private:
    vector<int> nums;
    std::function<int()> mt19937_rand;
public:
    Solution(vector<int>& nums) {
        this->nums = nums;
        mt19937_rand = bind(uniform_int_distribution<int>(0, INT_MAX),
mt19937(time(0)));
    }

    /** Resets the array to its original configuration and return it. */
    vector<int> reset() {
```

```

    return nums;
}

/** Returns a random shuffling of the array. */
vector<int> shuffle() {
    vector<int> shuffled(nums);
    int n = nums.size();
    for (int i = 0; i < n; ++i) {
        int j = mt19937_rand() % (i + 1);
        swap(shuffled[i], shuffled[j]);
    }
    return shuffled;
}
};

/**
 * Your Solution object will be instantiated and called as such:
 * Solution* obj = new Solution(nums);
 * vector<int> param_1 = obj->reset();
 * vector<int> param_2 = obj->shuffle();
 */

```

```

class Solution {
private:
    vector<int> nums;
    mt19937 mt19937_rand;
    unsigned seed;
public:
    Solution(vector<int>& nums) {
        this->nums = nums;
        seed = std::chrono::system_clock::now().time_since_epoch().count();
        mt19937_rand.seed(seed);
    }

    /** Resets the array to its original configuration and return it. */
    vector<int> reset() {
        seed = std::chrono::system_clock::now().time_since_epoch().count();
        mt19937_rand.seed(seed);
        return nums;
    }

    /** Returns a random shuffling of the array. */
    vector<int> shuffle() {
        vector<int> shuffled(nums);
        int n = nums.size();
        for (int i = 0; i < n; ++i) {
            int j = mt19937_rand() % (i + 1);
            swap(shuffled[i], shuffled[j]);
        }
    }
};

```

```
    }  
    return shuffled;  
  }  
};
```

## 387. First Unique Character in a String

### *Description*

Given a string, find the first non-repeating character in it and return its index. If it doesn't exist, return -1.

Examples:

s = "leetcode"  
return 0.

s = "loveleetcode",  
return 2.

Note: You may assume the string contains only lowercase letters.

### *Solution*

05/05/2020:

```
class Solution {  
public:  
    int firstUniqChar(string s) {  
        unordered_map<char, int> mp;  
        for (auto& c : s) ++mp[c];  
        int n = s.size();  
        for (int i = 0; i < n; ++i)  
            if (mp[s[i]] == 1)  
                return i;  
        return -1;  
    }  
};
```

## 389. Find the Difference

### *Description*

Given two strings *s* and *t* which consist of only lowercase letters.

String *t* is generated by random shuffling string *s* and then add one more letter at a random position.

Find the letter that was added in *t*.

Example:

Input:

*s* = "abcd"

*t* = "abcde"

Output:

e

Explanation:

'e' is the letter that was added.

*Solution*

05/17/2020:

```
class Solution {
public:
    char findTheDifference(string s, string t) {
        vector<int> cntS(26, 0);
        vector<int> cntT(26, 0);
        for (auto& c : s) {
            ++cntS[c - 'a'];
        }
        for (auto& c : t) {
            ++cntT[c - 'a'];
        }
        for (int i = 0; i < 26; ++i) {
            if (cntS[i] < cntT[i]) {
                return i + 'a';
            }
        }
        return 'a';
    }
};
```

392. Is Subsequence

## Description

Given a string `s` and a string `t`, check if `s` is subsequence of `t`.

You may assume that there is only lower case English letters in both `s` and `t`. `t` is potentially a very long (length  $\sim 500,000$ ) string, and `s` is a short string ( $\leq 100$ ).

A subsequence of a string is a new string which is formed from the original string by deleting some (can be none) of the characters without disturbing the relative positions of the remaining characters. (ie, "ace" is a subsequence of "abcde" while "aec" is not).

Example 1:

`s = "abc", t = "ahbgdc"`

Return true.

Example 2:

`s = "axc", t = "ahbgdc"`

Return false.

Follow up:

If there are lots of incoming `S`, say `S1, S2, ... , Sk` where  $k \geq 1B$ , and you want to check one by one to see if `T` has its subsequence. In this scenario, how would you change your code?

Credits:

Special thanks to @pbrother for adding this problem and creating all test cases.

## Solution

01/14/2020 (Dynamic Programming, Memory Limit Exceeded):

```
class Solution {
public:
    bool isSubsequence(string s, string t) {
        if (s.size() >= t.size()) return s == t;
        return s.back() == t.back() ? isSubsequence(s.substr(0, s.size() - 1),
t.substr(0, t.size() - 1)) :
        isSubsequence(s, t.substr(0, t.size() - 1));
    }
};
```

01/14/2020 (Dynamic Programming (bottom-up), Two pointers):



```

class Solution {
public:
    bool isSubsequence(string s, string t) {
        int ps = 0, pt = 0;
        for (; ps < s.size() && pt < t.size(); ++pt)
            if (s[ps] == t[pt]) ++ps;
        return ps == s.size();
    }
};

```

06/09/2020:

```

class Solution {
public:
    bool isSubsequence(string s, string t) {
        auto sit = s.begin(), tit = t.begin();
        for (; sit != s.end() && tit != t.end(); ++tit)
            if (*sit == *tit) ++sit;
        return sit == s.end();
    }
};

```

## 402. Remove K Digits

### *Description*

Given a non-negative integer num represented as a string, remove k digits from the number so that the new number is the smallest possible.

**Note:**

The length of num is less than 10002 and will be  $\geq k$ .

The given num does not contain any leading zero.

Example 1:

Input: num = "1432219", k = 3

Output: "1219"

Explanation: Remove the three digits 4, 3, and 2 to form the new number 1219 which is the smallest.

Example 2:

Input: num = "10200", k = 1

Output: "200"

Explanation: Remove the leading 1 and the number is 200. Note that the output must not contain leading zeroes.

Example 3:

Input: num = "10", k = 2

Output: "0"

Explanation: Remove all the digits from the number and it is left with nothing which is 0.

*Solution*

05/12/2020:

```
class Solution {
public:
    string removeKdigits(string num, int k) {
        deque<char> q;
        string ret;
        for (int i = 0; i < num.size(); ++i) {
            while (!q.empty() && q.back() > num[i] && k-- > 0) q.pop_back();
            q.push_back(num[i]);
        }
        while (!q.empty() && k-- > 0) q.pop_back();
        while (!q.empty() && q.front() == '0') q.pop_front();
        while (!q.empty()) ret += q.front(), q.pop_front();
        return ret.empty() ? string(1, '0') : ret;
    }
};
```

```
class Solution {
public:
    string removeKdigits(string num, int k) {
        string ret;
        for (int i = 0; i < num.size(); ++i) {
            while (!ret.empty() && ret.back() > num[i] && k-- > 0)
                ret.pop_back();
            if (!ret.empty() || num[i] != '0')
                ret += num[i];
        }
        while (!ret.empty() && k-- > 0)
            ret.pop_back();
        return ret.empty() ? "0" : ret;
    }
};
```

## 404. Sum of Left Leaves

## Description

Find the sum of all left leaves in a given binary tree.

Example:

```
    3
   / \
  9  20
   / \
  15  7
```

There are two left leaves in the binary tree, with values 9 and 15 respectively. Return 24.

## Solution

## Discussion

Recursive:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
class Solution {
public:
    int sumOfLeftLeaves(TreeNode* root, bool isLeft = false) {
        if (root == nullptr) return 0;
        int s = !root->left && !root->right && isLeft ? root->val : 0;
        return s + sumOfLeftLeaves(root->left, true) + sumOfLeftLeaves(root->right,
false);
    }
};
```

Iterative:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
```

```

*     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
* };
*/
class Solution {
public:
    int sumOfLeftLeaves(TreeNode* root) {
        if (root == nullptr) return 0;
        int s = 0;
        stack<pair<TreeNode*, bool>> st;
        st.emplace(root, false);
        while (!st.empty()) {
            pair<TreeNode*, bool> cur = st.top(); st.pop();
            s += cur.first->left == nullptr && cur.first->right == nullptr &&
cur.second ? cur.first->val : 0;
            if (cur.first->left) st.emplace(cur.first->left, true);
            if (cur.first->right) st.emplace(cur.first->right, false);
        }
        return s;
    }
};

```

## 406. Queue Reconstruction by Height

### Description

Suppose you have a random list of people standing in a queue. Each person is described by a pair of integers (h, k), where h is the height of the person and k is the number of people in front of this person who have a height greater than or equal to h. Write an algorithm to reconstruct the queue.

#### Note:

The number of people is less than 1,100.

#### Example

##### Input:

```
[[7,0], [4,4], [7,1], [5,0], [6,1], [5,2]]
```

##### Output:

```
[[5,0], [7,0], [5,2], [6,1], [4,4], [7,1]]
```

### Solution

06/06/2020:

```

class Solution {
public:
    vector<vector<int>> reconstructQueue(vector<vector<int>>& people) {
        if (people.empty() || people[0].empty()) return people;
        sort(people.begin(), people.end(), [](vector<int>& p1, vector<int>& p2) {
            if (p1[0] == p2[0]) return p1[1] < p2[1];
            return p1[0] > p2[0];
        });
        list<vector<int>> ret;
        for (int i = 0; i < (int)people.size(); ++i)
            ret.insert(next(ret.begin()), people[i][1], people[i]);
        return vector<vector<int>>(ret.begin(), ret.end());
    }
};

```

```

class Solution {
public:
    vector<vector<int>> reconstructQueue(vector<vector<int>>& people) {
        if (people.empty() || people[0].empty()) return people;
        sort(people.begin(), people.end(), [](vector<int>& p1, vector<int>& p2) {
            if (p1[0] == p2[0]) return p1[1] < p2[1];
            return p1[0] > p2[0];
        });
        vector<vector<int>> ret;
        for (int i = 0; i < (int)people.size(); ++i)
            ret.insert(ret.begin() + people[i][1], people[i]);
        return ret;
    }
};

```

## 409. Longest Palindrome

### Description

Given a string which consists of lowercase or uppercase letters, find the length of the longest palindromes that can be built with those letters.

This is case sensitive, for example "Aa" is not considered a palindrome here.

Note:

Assume the length of given string will not exceed 1,010.

Example:

Input:  
"abccccdd"

Output:  
7

Explanation:  
One longest palindrome that can be built is "dccaccd", whose length is 7.

*Solution*

05/18/2020:

```
class Solution {
public:
    int longestPalindrome(string s) {
        vector<int> cnt(128, 0);
        for (auto& c : s) {
            ++cnt[c];
        }
        bool odd = false;
        int even = 0;
        for (int i = 0; i < 128; ++i) {
            if (cnt[i] % 2 == 0) {
                even += cnt[i];
            } else {
                odd = true;
                even += max(cnt[i] - 1, 0);
            }
        }
        return odd ? even + 1 : even;
    }
};
```

## 415. Add Strings

*Description*

Given two non-negative integers num1 and num2 represented as string, return the sum of num1 and num2.

Note:

The length of both num1 and num2 is < 5100.

Both num1 and num2 contains only digits 0-9.

Both num1 and num2 does not contain any leading zero.

You must not use any built-in BigInteger library or convert the inputs to integer directly.

*Solution*

06/09/2020:

```
class Solution {
public:
    string addStrings(string num1, string num2) {
        int carry = 0;
        int n1 = num1.size(), n2 = num2.size();
        string ret;
        for (int i1 = n1 - 1, i2 = n2 - 1; i1 >= 0 || i2 >= 0 || carry > 0; --i1, --i2) {
            if (i1 >= 0) carry += num1[i1] - '0';
            if (i2 >= 0) carry += num2[i2] - '0';
            ret.push_back(carry % 10 + '0');
            carry /= 10;
        }
        reverse(ret.begin(), ret.end());
        return ret;
    }
};
```

## 429. N-ary Tree Level Order Traversal

*Description*

Given an n-ary tree, return the level order traversal of its nodes' values.

Nary-Tree input serialization is represented in their level order traversal, each group of children is separated by the null value (See examples).

Example 1:

Input: root = [1,null,3,2,4,null,5,6]

Output: [[1],[3,2,4],[5,6]]

Example 2:

Input: root =

[1,null,2,3,4,5,null,null,6,7,null,8,null,9,10,null,null,11,null,12,null,13,null,null,14]

Output: [[1],[2,3,4,5],[6,7,8,9,10],[11,12,13],[14]]

Constraints:

The height of the n-ary tree is less than or equal to 1000

The total number of nodes is between [0, 10<sup>4</sup>]

*Solution*

04/26/2020:

```
/*
// Definition for a Node.
class Node {
public:
    int val;
    vector<Node*> children;

    Node() {}

    Node(int _val) {
        val = _val;
    }

    Node(int _val, vector<Node*> _children) {
        val = _val;
        children = _children;
    }
};
*/

class Solution {
public:
    vector<vector<int>> levelOrder(Node* root) {
        if (root == nullptr) return {};
    }
};
```



```

queue<Node*> q;
q.push(root);
vector<vector<int>> ret;
ret.push_back({root->val});
while (!q.empty()) {
    int n = q.size();
    vector<int> level;
    for (int i = 0; i < n; ++i) {
        Node* cur = q.front(); q.pop();
        for (auto& c : cur->children) {
            level.push_back(c->val);
            q.push(c);
        }
    }
    if (!level.empty()) ret.push_back(level);
}
return ret;
}
};

```

## 438. Find All Anagrams in a String

### Description

Given a string *s* and a non-empty string *p*, find all the start indices of *p*'s anagrams in *s*.

Strings consists of lowercase English letters only and the length of both strings *s* and *p* will not be larger than 20,100.

The order of output does not matter.

Example 1:

Input:

s: "cbaebabacd" p: "abc"

Output:

[0, 6]

Explanation:

The substring with start index = 0 is "cba", which is an anagram of "abc".

The substring with start index = 6 is "bac", which is an anagram of "abc".

Example 2:

Input:

s: "abab" p: "ab"

Output:

[0, 1, 2]

Explanation:

The substring with start index = 0 is "ab", which is an anagram of "ab".

The substring with start index = 1 is "ba", which is an anagram of "ab".

The substring with start index = 2 is "ab", which is an anagram of "ab".

*Solution*

05/16/2020:

```
class Solution {
public:
    vector<int> findAnagrams(string s, string p) {
        vector<int> sch(26, 0), pch(26, 0), ret;
        for (auto& c : p) ++pch[c - 'a'];
        int m = p.size(), n = s.size();
        if (m > n) return {};
        for (int i = 0; i < m - 1; ++i) ++sch[s[i] - 'a'];
        bool ok = false;
        for (int i = m - 1; i < n; ++i) {
            if (ok) {
                if (s[i] == s[i - m]) {
                    ret.push_back(i - m + 1);
                } else {
                    --sch[s[i - m] - 'a'];
                    ++sch[s[i] - 'a'];
                    ok = false;
                }
            } else {
                if (i >= m) --sch[s[i - m] - 'a'];
                ++sch[s[i] - 'a'];
                ok = true;
                for (int j = 0; j < 26; ++j) {
                    if (pch[j] != sch[j]) {
                        ok = false;
                        break;
                    }
                }
                if (ok) ret.push_back(i - m + 1);
            }
        }
        return ret;
    }
};
```

```

class Solution {
public:
    vector<int> findAnagrams(string s, string p) {
        vector<int> sch(26, 0), pch(26, 0), ret;
        for (auto& c : p) ++pch[c - 'a'];
        int m = p.size(), n = s.size();
        for (int i = 0; i < n; ++i) {
            if (i - m >= 0) --sch[s[i - m] - 'a'];
            ++sch[s[i] - 'a'];
            bool ok = true;
            for (int j = 0; j < 26; ++j) {
                if (pch[j] != sch[j]) {
                    ok = false;
                    break;
                }
            }
            if (ok) ret.push_back(i - m + 1);
        }
        return ret;
    }
};

```

## 447. Number of Boomerangs

### *Description*

Given  $n$  points in the plane that are all pairwise distinct, a "boomerang" is a tuple of points  $(i, j, k)$  such that the distance between  $i$  and  $j$  equals the distance between  $i$  and  $k$  (the order of the tuple matters).

Find the number of boomerangs. You may assume that  $n$  will be at most 500 and coordinates of points are all in the range  $[-10000, 10000]$  (inclusive).

Example:

Input:  
 $[[0,0], [1,0], [2,0]]$

Output:  
 2

Explanation:  
 The two boomerangs are  $[[1,0], [0,0], [2,0]]$  and  $[[1,0], [2,0], [0,0]]$

### *Solution*

05/20/2020:

```
class Solution {
public:
    int numberOfBoomerangs(vector<vector<int>>& points) {
        if (points.empty() || points[0].empty()) return 0;
        int n = points.size(), cnt = 0;
        vector<unordered_map<double, int>> mp(n);
        for (int i = 0; i < n - 1; ++i) {
            for (int j = i + 1; j < n; ++j) {
                int x1 = points[i][0], x2 = points[j][0];
                int y1 = points[i][1], y2 = points[j][1];
                double d = hypot(x1 - x2, y1 - y2);
                ++mp[i][d];
                ++mp[j][d];
            }
        }
        for (int i = 0; i < n; ++i)
            for (auto& m : mp[i])
                cnt += nchoosek(m.second, 2) * 2;
        return cnt;
    }

    int nchoosek(int n, int k) {
        if (n < k) return 0;
        // n! / (k!) (n - k)!;
        // n * (n - 1) * ... * (n - k + 1) / k!
        long long num = 1, dem = 1;
        int kk = min(n - k, k);
        for (int i = 1; i <= kk; ++i) {
            num *= (n - i + 1);
            dem *= i;
        }
        return num / dem;
    }
};
```

## 451. Sort Characters By Frequency

### Description

Given a string, sort it in decreasing order based on the frequency of characters.

Example 1:

Input:  
"tree"

Output:  
"eert"

Explanation:

'e' appears twice while 'r' and 't' both appear once.

So 'e' must appear before both 'r' and 't'. Therefore "eetr" is also a valid answer.

Example 2:

Input:  
"cccaaa"

Output:  
"cccaaa"

Explanation:

Both 'c' and 'a' appear three times, so "aaaccc" is also a valid answer.

Note that "cacaca" is incorrect, as the same characters must be together.

Example 3:

Input:  
"Aabb"

Output:  
"bbAa"

Explanation:

"bbaA" is also a valid answer, but "Aabb" is incorrect.

Note that 'A' and 'a' are treated as two different characters.

*Solution*

05/20/2020:

```

class Solution {
public:
    string frequencySort(string s) {
        unordered_map<char, int> mp;
        for (auto& c : s) ++mp[c];
        vector<pair<int, char>> str;
        for (auto& m : mp) str.emplace_back(m.second, m.first);
        sort(str.begin(), str.end(), greater<pair<int, char>>());
        string ret = "";
        for (auto& s : str) ret += string(s.first, s.second);
        return ret;
    }
};

```

## 468. Validate IP Address

### Description

Write a function to check whether an input string is a valid IPv4 address or IPv6 address or neither.

IPv4 addresses are canonically represented in dot-decimal notation, which consists of four decimal numbers, each ranging from 0 to 255, separated by dots ("."), e.g.,172.16.254.1;

Besides, leading zeros in the IPv4 is invalid. For example, the address 172.16.254.01 is invalid.

IPv6 addresses are represented as eight groups of four hexadecimal digits, each group representing 16 bits. The groups are separated by colons (":"). For example, the address 2001:0db8:85a3:0000:0000:8a2e:0370:7334 is a valid one. Also, we could omit some leading zeros among four hexadecimal digits and some low-case characters in the address to upper-case ones, so 2001:db8:85a3:0:0:8A2E:0370:7334 is also a valid IPv6 address(Omit leading zeros and using upper cases).

However, we don't replace a consecutive group of zero value with a single empty group using two consecutive colons (::) to pursue simplicity. For example, 2001:0db8:85a3::8A2E:0370:7334 is an invalid IPv6 address.

Besides, extra leading zeros in the IPv6 is also invalid. For example, the address 02001:0db8:85a3:0000:0000:8a2e:0370:7334 is invalid.

Note: You may assume there is no extra space or special characters in the input string.

Example 1:

Input: "172.16.254.1"

Output: "IPv4"

Explanation: This is a valid IPv4 address, return "IPv4".

Example 2:

Input: "2001:0db8:85a3:0:0:8A2E:0370:7334"

Output: "IPv6"

Explanation: This is a valid IPv6 address, return "IPv6".

Example 3:

Input: "256.256.256.256"

Output: "Neither"

Explanation: This is neither a IPv4 address nor a IPv6 address.

*Solution*

06/16/2020:

```
class Solution {
public:
    string validIPAddress(string IP) {
        int n = IP.size();
        vector<string> ipv4 = split(IP, '.');
        vector<string> ipv6 = split(IP, ':');
        bool isIpv4 = ipv4.size() == 4 && isdigit(IP.front()) && isdigit(IP.back());
        bool isIpv6 = ipv6.size() == 8 && (isalpha(IP.front()) ||
        isdigit(IP.front())) && (isalpha(IP.back()) || isdigit(IP.back()));
        bool hasLeadingZero = false;
        if (isIpv4) {
            for (auto& s : ipv4) {
                isIpv4 = isIpv4 && s.size() <= 3 && s.size() > 0;
                for (auto& c : s) isIpv4 = isIpv4 && isdigit(c);
                if (!isIpv4) break;
                int n = stoi(s);
                isIpv4 = isIpv4 && (0 <= n && n <= 255 && to_string(n) == s);
            }
        } else if (isIpv6) {
            for (auto& s : ipv6) {
                isIpv6 = isIpv6 && s.size() > 0 && s.size() < 5;
                for (auto& c : s) isIpv6 = isIpv6 && (('0' <= c && c <= '9') || ('a' <=
                c && c <= 'f') || ('A' <= c && c <= 'F'));
                if (!isIpv6) break;
            }
        }
    }
};
```

```

    }
}
return isIpv4 ? "IPv4" : isIpv6 ? "IPv6" : "Neither";
}

vector<string> split(string IP, char delimiter) {
    istringstream iss(IP);
    string s;
    vector<string> ret;
    while (getline(iss, s, delimiter)) ret.push_back(s);
    return ret;
}
};

```

## 470. Implement Rand10() Using Rand7()

### *Description*

Given a function rand7 which generates a uniform random integer in the range 1 to 7, write a function rand10 which generates a uniform random integer in the range 1 to 10.

Do NOT use system's Math.random().

Example 1:

Input: 1

Output: [7]

Example 2:

Input: 2

Output: [8,4]

Example 3:

Input: 3

Output: [8,1,10]

Note:

rand7 is predefined.

Each testcase has one argument: n, the number of times that rand10 is called.



Follow up:

What is the expected value for the number of calls to rand7() function?  
Could you minimize the number of calls to rand7()?

*Solution*

05/27/2020:

```
// The rand7() API is already defined for you.
// int rand7();
// @return a random integer in the range 1 to 7

class Solution {
public:
    int rand10() {
        int s = 41;
        while (s > 40) s = (rand7() - 1) * 7 + rand7();
        return (s - 1) % 10 + 1;
    }
};
```

## 475. Heaters

*Description*

Winter is coming! Your first job during the contest is to design a standard heater with fixed warm radius to warm all the houses.

Now, you are given positions of houses and heaters on a horizontal line, find out minimum radius of heaters so that all houses could be covered by those heaters.

So, your input will be the positions of houses and heaters separately, and your expected output will be the minimum radius standard of heaters.

Note:

Numbers of houses and heaters you are given are non-negative and will not exceed 25000.

Positions of houses and heaters you are given are non-negative and will not exceed  $10^9$ .

As long as a house is in the heaters' warm radius range, it can be warmed.

All the heaters follow your radius standard and the warm radius will be the same.

Example 1:

Input: [1,2,3],[2]

Output: 1

Explanation: The only heater was placed in the position 2, and if we use the radius 1 standard, then all the houses can be warmed.

Example 2:

Input: [1,2,3,4],[1,4]

Output: 1

Explanation: The two heater was placed in the position 1 and 4. We need to use radius 1 standard, then all the houses can be warmed.

*Solution*

04/23/2020:

```
class Solution {
public:
    int findRadius(vector<int>& houses, vector<int>& heaters) {
        if ((int)houses.size() <= 0) return 0;
        int n = heaters.size();
        if (n <= 0) return INT_MAX;
        sort(heaters.begin(), heaters.end());
        for (auto& h : houses) {
            int lo = 0, hi = n - 1;
            while (lo <= hi) {
                int mid = lo + (hi - lo) / 2;
                if (heaters[mid] <= h) {
                    lo = mid + 1;
                } else {
                    hi = mid - 1;
                }
            }
            int left = INT_MAX, right = INT_MAX;
            if (lo - 1 >= 0) left = min(right, h - heaters[lo - 1]);
            if (lo < n) right = min(left, heaters[lo] - h);
            h = min(left, right);
        }
        return *max_element(houses.begin(), houses.end());
    }
};
```

## 482. License Key Formatting

### Description

You are given a license key represented as a string  $S$  which consists only alphanumeric character and dashes. The string is separated into  $N+1$  groups by  $N$  dashes.

Given a number  $K$ , we would want to reformat the strings such that each group contains exactly  $K$  characters, except for the first group which could be shorter than  $K$ , but still must contain at least one character. Furthermore, there must be a dash inserted between two groups and all lowercase letters should be converted to uppercase.

Given a non-empty string  $S$  and a number  $K$ , format the string according to the rules described above.

Example 1:

Input:  $S = "5F3Z-2e-9-w"$ ,  $K = 4$

Output:  $"5F3Z-2E9W"$

Explanation: The string  $S$  has been split into two parts, each part has 4 characters.

Note that the two extra dashes are not needed and can be removed.

Example 2:

Input:  $S = "2-5g-3-J"$ ,  $K = 2$

Output:  $"2-5G-3J"$

Explanation: The string  $S$  has been split into three parts, each part has 2 characters except the first part as it could be shorter as mentioned above.

Note:

The length of string  $S$  will not exceed 12,000, and  $K$  is a positive integer. String  $S$  consists only of alphanumerical characters (a-z and/or A-Z and/or 0-9) and dashes(-).

String  $S$  is non-empty.

### Solution

02/05/2020:

```
class Solution {  
public:
```

```

string licenseKeyFormatting(string S, int K) {
    string ret;
    int cnt = 0;
    for (auto it = S.rbegin(); it != S.rend(); ++it) {
        if (*it == '-') {
            continue;
        } else {
            ret += toupper(*it);
        }
        if (++cnt % K == 0) {
            cnt == 0;
            ret += '-';
        }
    }
    if (ret.back() == '-') {
        ret.pop_back();
    }
    reverse(ret.begin(), ret.end());
    return ret;
}
};

```

## 498. Diagonal Traverse

### *Description*

Given a matrix of  $M \times N$  elements ( $M$  rows,  $N$  columns), return all elements of the matrix in diagonal order as shown in the below image.

Example:

Input:

```

[
  [ 1, 2, 3 ],
  [ 4, 5, 6 ],
  [ 7, 8, 9 ]
]

```

Output: [1,2,4,7,5,3,6,8,9]

Explanation:

Note:

The total number of elements of the given matrix will not exceed 10,000.

*Solution*

05/20/2020:

```
class Solution {
public:
    vector<int> findDiagonalOrder(vector<vector<int>>& matrix) {
        if (matrix.empty() || matrix[0].empty()) return {};
        int m = matrix.size(), n = matrix[0].size();
        vector<int> ret;
        bool ascending = false;
        for (int k = 0; k < m + n; ++k) {
            if (ascending)
                for (int i = max(0, k - n + 1); i <= min(m - 1, k); ++i)
                    ret.push_back(matrix[i][k - i]);
            else
                for (int i = min(m - 1, k); i >= max(0, k - n + 1); --i)
                    ret.push_back(matrix[i][k - i]);
            ascending = !ascending;
        }
        return ret;
    }
};
```

## 510. Inorder Successor in BST II

*Description*

Given a node in a binary search tree, find the in-order successor of that node in the BST.

If that node has no in-order successor, return null.

The successor of a node is the node with the smallest key greater than node.val.

You will have direct access to the node but not to the root of the tree. Each node will have a reference to its parent node. Below is the definition for Node:

```
class Node {
    public int val;
    public Node left;
```

```
    public Node right;
    public Node parent;
}
```

Follow up:

Could you solve it without looking up any of the node's values?

Example 1:

Input: tree = [2,1,3], node = 1

Output: 2

Explanation: 1's in-order successor node is 2. Note that both the node and the return value is of Node type.

Example 2:

Input: tree = [5,3,6,2,4,null,null,1], node = 6

Output: null

Explanation: There is no in-order successor of the current node, so the answer is null.

Example 3:

Input: tree =

[15,6,18,3,7,17,20,2,4,null,13,null,null,null,null,null,null,null,9], node = 15

Output: 17

Example 4:

Input: tree =

[15,6,18,3,7,17,20,2,4,null,13,null,null,null,null,null,null,null,9], node = 13

Output: 15

Example 5:

Input: tree = [0], node = 0

Output: null

Constraints:

$-10^5 \leq \text{Node.val} \leq 10^5$

$1 \leq \text{Number of Nodes} \leq 10^4$

All Nodes will have unique values.

*Solution*

05/24/2020:

```
/*
// Definition for a Node.
class Node {
public:
    int val;
    Node* left;
    Node* right;
    Node* parent;
};
*/

class Solution {
public:
    Node* inorderSuccessor(Node* node) {
        if (node->right != nullptr) return leftmostChild(node->right);
        Node* parent = smallestParentGreaterThanNode(node);
        return parent == nullptr || parent->val < node->val ? nullptr : parent;
    }

    Node* leftmostChild(Node* root) {
        if (root == nullptr) return nullptr;
        if (root->left == nullptr) return root;
        return leftmostChild(root->left);
    }

    Node* smallestParentGreaterThanNode(Node* node) {
        if (node == nullptr || node->parent == nullptr) return nullptr;
        if (node->parent->val > node->val) return node->parent;
        return smallestParentGreaterThanNode(node->parent);
    }
};
```

```
class Solution {
public:
    Node* inorderSuccessor(Node* node) {
        if (node->right != nullptr) return leftmost(node->right);
        Node* parent = node->parent;
        while (true) {
            if (parent != nullptr && parent->val < node->val && parent->parent !=
                nullptr) {
                parent = parent->parent;
            }
        }
    }
};
```

```

        } else {
            break;
        }
    }
    return parent != nullptr && parent->val < node->val ? nullptr : parent;
}

Node* leftmost(Node* root) {
    if (root == nullptr) return nullptr;
    if (root->left == nullptr) return root;
    return leftmost(root->left);
}
};

```

```

class Solution {
public:
    Node* inorderSuccessor(Node* node) {
        Node* root = node;
        while (true) {
            if (root->parent != nullptr)
                root = root->parent;
            else
                break;
        }
        vector<Node*> inorder = inorderTraversal(root);
        auto it = find(inorder.begin(), inorder.end(), node);
        return it != inorder.end() && it + 1 != inorder.end() ? *(it + 1) : nullptr;
    }

    vector<Node*> inorderTraversal(Node* root) {
        if (root == nullptr) return {};
        vector<Node*> leftTraversal = inorderTraversal(root->left);
        vector<Node*> rightTraversal = inorderTraversal(root->right);
        leftTraversal.insert(leftTraversal.end(), root);
        leftTraversal.insert(leftTraversal.end(), rightTraversal.begin(),
rightTraversal.end());
        return leftTraversal;
    }
};

```

## 518. Coin Change 2

*Description*



You are given coins of different denominations and a total amount of money. Write a function to compute the number of combinations that make up that amount. You may assume that you have infinite number of each kind of coin.

Example 1:

Input: amount = 5, coins = [1, 2, 5]

Output: 4

Explanation: there are four ways to make up the amount:

5=5

5=2+2+1

5=2+1+1+1

5=1+1+1+1+1

Example 2:

Input: amount = 3, coins = [2]

Output: 0

Explanation: the amount of 3 cannot be made up just with coins of 2.

Example 3:

Input: amount = 10, coins = [10]

Output: 1

Note:

You can assume that

$0 \leq \text{amount} \leq 5000$

$1 \leq \text{coin} \leq 5000$

the number of coins is less than 500

the answer is guaranteed to fit into signed 32-bit integer

*Solution*

05/27/2020:

```
class Solution {
public:
    int change(int amount, vector<int>& coins) {
        vector<int> dp(amount + 1, 0);
        dp[0] = 1;
        for (auto& c : coins) {
            for (int i = c; i <= amount; ++i) {
                if (i - c >= 0) {
                    dp[i] += dp[i - c];
                }
            }
        }
    }
};
```

```

    }
  }
}
return dp.back();
}
};

```

## 525. Contiguous Array

### Description

Given a binary array, find the maximum length of a contiguous subarray with equal number of 0 and 1.

Example 1:

Input: [0,1]

Output: 2

Explanation: [0, 1] is the longest contiguous subarray with equal number of 0 and 1.

Example 2:

Input: [0,1,0]

Output: 2

Explanation: [0, 1] (or [1, 0]) is a longest contiguous subarray with equal number of 0 and 1.

Note: The length of the given binary array will not exceed 50,000.

### Solution

05/26/2020:

```

class Solution {
public:
    int findMaxLength(vector<int>& nums) {
        if (nums.empty()) return 0;
        unordered_map<int, int> firstOccurrence{ {0, -1} };
        int n = nums.size(), ret = 0, prefix = 0;
        for (int i = 0; i < n; ++i) {
            prefix = prefix + (nums[i] == 0 ? 1 : -1);
            if (firstOccurrence.count(prefix) == 0)
                firstOccurrence[prefix] = i;
            else
                ret = max(ret, i - firstOccurrence[prefix]);
        }
        return ret;
    }
}

```

```
};
```

## 528. Random Pick with Weight

### Description

Given an array `w` of positive integers, where `w[i]` describes the weight of index `i`, write a function `pickIndex` which randomly picks an index in proportion to its weight.

Note:

```
1 <= w.length <= 10000
```

```
1 <= w[i] <= 10^5
```

`pickIndex` will be called at most 10000 times.

Example 1:

Input:

```
["Solution","pickIndex"]
```

```
[[[1]],[]]
```

Output: [null,0]

Example 2:

Input:

```
["Solution","pickIndex","pickIndex","pickIndex","pickIndex","pickIndex"]
```

```
[[[1,3]],[],[],[],[],[]]
```

Output: [null,0,1,1,1,0]

Explanation of Input Syntax:

The input is two lists: the subroutines called and their arguments. `Solution`'s constructor has one argument, the array `w`. `pickIndex` has no arguments. Arguments are always wrapped with a list, even if there aren't any.

### Solution

06/05/2020:

```
class Solution {
private:
    vector<int> weights;
    mt19937 rand;

public:
    Solution(vector<int>& w) {
        unsigned seed = std::chrono::system_clock::now().time_since_epoch().count();
```

```

    rand.seed(seed);
    weights.resize(w.size());
    partial_sum(w.begin(), w.end(), weights.begin(), plus<int>());
}

int pickIndex() {
    long long n = (double) rand() / rand.max() * weights.back();
    return upper_bound(weights.begin(), weights.end(), n) - weights.begin();
}
};

/**
 * Your Solution object will be instantiated and called as such:
 * Solution* obj = new Solution(w);
 * int param_1 = obj->pickIndex();
 */

```

## 535. Encode and Decode TinyURL

### Description

Note: This is a companion problem to the System Design problem: Design TinyURL ([https://leetcode.com/discuss/interview-question/124658/Design-a-URL-Shortener-\(-TinyURL-\)-System/](https://leetcode.com/discuss/interview-question/124658/Design-a-URL-Shortener-(-TinyURL-)-System/)).

TinyURL is a URL shortening service where you enter a URL such as <https://leetcode.com/problems/design-tinyurl> and it returns a short URL such as <http://tinyurl.com/4e9iAk>.

Design the encode and decode methods for the TinyURL service. There is no restriction on how your encode/decode algorithm should work. You just need to ensure that a URL can be encoded to a tiny URL and the tiny URL can be decoded to the original URL.

### Solution

05/20/2020:

```

class Solution {
private:
    unordered_map<string, string> decodeTable;
    unordered_map<string, string> encodeTable;
    string base = "http://tinyurl.com/";
    const int MOD = 1e9 + 7;

    int hash(string s) {

```

```

    long long h = 1;
    for (auto& c : s) h = (h * 131 + c) % MOD;
    return h;
}

public:

    // Encodes a URL to a shortened URL.
    string encode(string longUrl) {
        if (encodeTable.count(longUrl) == 0) {
            string shortUrl = base + to_string(hash(longUrl));
            encodeTable[longUrl] = shortUrl;
            decodeTable[shortUrl] = longUrl;
        }
        return encodeTable[longUrl];
    }

    // Decodes a shortened URL to its original URL.
    string decode(string shortUrl) {
        return decodeTable[shortUrl];
    }
};

// Your Solution object will be instantiated and called as such:
// Solution solution;
// solution.decode(solution.encode(url));

```

## 540. Single Element in a Sorted Array

### Description

You are given a sorted array consisting of only integers where every element appears exactly twice, except for one element which appears exactly once. Find this single element that appears only once.

Follow up: Your solution should run in  $O(\log n)$  time and  $O(1)$  space.

Example 1:

Input: nums = [1,1,2,3,3,4,4,8,8]

Output: 2

Example 2:

Input: nums = [3,3,7,7,10,11,11]

Output: 10

Constraints:

$1 \leq \text{nums.length} \leq 10^5$

$0 \leq \text{nums}[i] \leq 10^5$

*Solution*

05/12/2020:

```
class Solution {
public:
    int singleNonDuplicate(vector<int>& nums) {
        int lo = 0, hi = nums.size() - 1;
        while (lo < hi) {
            int mid = lo + (hi - lo) / 2;
            if (mid % 2 == 0) {
                if (nums[mid] != nums[mid + 1]) {
                    hi = mid;
                } else {
                    lo = mid + 1;
                }
            } else {
                if (nums[mid] != nums[mid - 1]) {
                    hi = mid;
                } else {
                    lo = mid + 1;
                }
            }
        }
        return nums[lo];
    }
};
```

```
class Solution {
public:
    int singleNonDuplicate(vector<int>& nums) {
        int lo = 0, hi = nums.size() - 1;
        while (lo < hi) {
            int mid = lo + (hi - lo) / 2;
            if ((!(mid & 1) && nums[mid] != nums[mid + 1]) || ((mid & 1) && nums[mid]
!= nums[mid - 1])) {
                hi = mid;
            } else {
                lo = mid + 1;
            }
        }
    }
};
```

```
    }  
  }  
  return nums[lo];  
}  
};
```

```
class Solution {  
public:  
    int singleNonDuplicate(vector<int>& nums) {  
        int ret = 0;  
        for (auto& n : nums) {  
            ret ^= n;  
        }  
        return ret;  
    }  
};
```

## 542. 01 Matrix

### *Description*

Given a matrix consists of 0 and 1, find the distance of the nearest 0 for each cell.

The distance between two adjacent cells is 1.

Example 1:

Input:

```
[[0,0,0],  
 [0,1,0],  
 [0,0,0]]
```

Output:

```
[[0,0,0],  
 [0,1,0],  
 [0,0,0]]
```

Example 2:

Input:

```
[[0,0,0],  
 [0,1,0],  
 [1,1,1]]
```

Output:  
[[0,0,0],  
 [0,1,0],  
 [1,2,1]]

Note:

The number of elements of the given matrix will not exceed 10,000.  
There are at least one 0 in the given matrix.  
The cells are adjacent in only four directions: up, down, left and right.

*Solution*

06/09/2020: BFS:

```
class Solution {
public:
    vector<vector<int>> updateMatrix(vector<vector<int>>& matrix) {
        if (matrix.empty() || matrix[0].empty()) return matrix;
        int m = matrix.size(), n = matrix[0].size();
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        queue<pair<int, int >> q;
        for (int i = 0; i < m; ++i)
            for (int j = 0; j < n; ++j)
                if (matrix[i][j] == 0)
                    q.emplace(i, j);
                else
                    matrix[i][j] = INT_MAX;

        while (!q.empty()) {
            pair<int, int> cur = q.front(); q.pop();
            int i = cur.first, j = cur.second;
            for (int d = 0; d < 4; ++d) {
                int ni = i + dir[d][0], nj = j + dir[d][1];
                if (ni < 0 || ni >= m || nj < 0 || nj >= n) continue;
                if (matrix[ni][nj] > matrix[i][j] + 1) {
                    matrix[ni][nj] = matrix[i][j] + 1;
                    q.emplace(ni, nj);
                }
            }
        }
        return matrix;
    }
};
```



```

class Solution {
public:
    vector<vector<int>> updateMatrix(vector<vector<int>>& matrix) {
        if (matrix.empty() || matrix[0].empty()) return matrix;
        int m = matrix.size(), n = matrix[0].size();
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (matrix[i][j] == 0) continue;
                int cur_min = INT_MAX;
                queue<pair<int, int>> q;
                q.emplace(i, j);
                bool found = false;
                unordered_set<string> visited;
                while (!q.empty() && !found) {
                    int sz = q.size();
                    for (int s = 0; s < sz && !found; ++s) {
                        pair<int, int> cur = q.front(); q.pop();
                        string key = to_string(cur.first) + "," + to_string(cur.second);
                        if (visited.count(key) > 0) continue;
                        for (int d = 0; d < 4; ++d) {
                            int ni = cur.first + dir[d][0], nj = cur.second + dir[d][1];
                            string key = to_string(ni) + "," + to_string(nj);
                            if (ni < 0 || ni >= m || nj < 0 || nj >= n) continue;
                            if (matrix[ni][nj] == 0) {
                                found = true;
                                matrix[i][j] = (abs(i - ni) + abs(j - nj));
                                break;
                            } else {
                                q.emplace(ni, nj);
                            }
                        }
                    }
                    visited.insert(key);
                }
            }
        }
        return matrix;
    }
};

```

Dynamic Programming:

```

class Solution {
public:
    vector<vector<int>> updateMatrix(vector<vector<int>>& matrix) {
        if (matrix.empty() || matrix[0].empty()) return matrix;
        int m = matrix.size(), n = matrix[0].size();

```

```

int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
queue<pair<int, int >> q;
for (int i = 0; i < m; ++i)
    for (int j = 0; j < n; ++j)
        if (matrix[i][j] != 0)
            matrix[i][j] = INT_MAX;
for (int i = 0; i < m; ++i) {
    for (int j = 0; j < n; ++j) {
        if (i > 0 && matrix[i - 1][j] != INT_MAX) matrix[i][j] = min(matrix[i]
[j], matrix[i - 1][j] + 1);
        if (j > 0 && matrix[i][j - 1] != INT_MAX) matrix[i][j] = min(matrix[i]
[j], matrix[i][j - 1] + 1);
    }
}
for (int i = m - 1; i >= 0; --i) {
    for (int j = n - 1; j >= 0; --j) {
        if (i < m - 1 && matrix[i + 1][j] != INT_MAX) matrix[i][j] =
min(matrix[i][j], matrix[i + 1][j] + 1);
        if (j < n - 1 && matrix[i][j + 1] != INT_MAX) matrix[i][j] =
min(matrix[i][j], matrix[i][j + 1] + 1);
    }
}
return matrix;
}
};

```

## 551. Student Attendance Record I

### Description

You are given a string representing an attendance record for a student. The record only contains the following three characters:

'A' : Absent.

'L' : Late.

'P' : Present.

A student could be rewarded if his attendance record doesn't contain more than one 'A' (absent) or more than two continuous 'L' (late).

You need to return whether the student could be rewarded according to his attendance record.

Example 1:

Input: "PPALLP"

Output: True

Example 2:

Input: "PPALLL"  
Output: False

*Solution*

02/05/2020:

```
class Solution {
public:
    bool checkRecord(string s) {
        int cntA = 0, cntLL = 0;
        for (auto i = 0; i < s.size(); ++i) {
            if (s[i] == 'A') {
                ++cntA;
            } else if (s[i] == 'L') {
                if (i > 0 && s[i - 1] == 'L') {
                    ++cntLL;
                } else {
                    cntLL = 1;
                }
            }
        }
        if (cntA > 1 || cntLL > 2) {
            return false;
        }
        return true;
    }
};
```

## 567. Permutation in String

*Description*

Given two strings s1 and s2, write a function to return true if s2 contains the permutation of s1. In other words, one of the first string's permutations is the substring of the second string.

Example 1:

Input: s1 = "ab" s2 = "eidbaooo"

Output: True

Explanation: s2 contains one permutation of s1 ("ba").

Example 2:

Input:s1= "ab" s2 = "eidboao"

Output: False

Note:

The input strings only contain lower case letters.  
The length of both given strings is in range [1, 10,000].

*Solution*

05/18/2020:

```
class Solution {
public:
    bool checkInclusion(string s, string t) {
        vector<int> cntS(26, 0), cntT(26, 0);
        int n = s.size(), m = t.size();
        for (auto& c : s) ++cntS[c - 'a'];
        for (int i = 0; i < m; ++i) {
            ++cntT[t[i] - 'a'];
            bool ok = true;
            if (i - n >= 0) --cntT[t[i - n] - 'a'];
            for (int j = 0; j < 26; ++j)
                if (cntS[j] != cntT[j]) ok = false;
            if (ok) return true;
        }
        return false;
    }
};
```

## 609. Find Duplicate File in System

*Description*

Given a list of directory info including directory path, and all the files with contents in this directory, you need to find out all the groups of duplicate files in the file system in terms of their paths.

A group of duplicate files consists of at least two files that have exactly the same content.

A single directory info string in the input list has the following format:

```
"root/d1/d2/.../dm f1.txt(f1_content) f2.txt(f2_content) ... fn.txt(fn_content)"
```

It means there are  $n$  files (`f1.txt`, `f2.txt` ... `fn.txt` with content `f1_content`, `f2_content` ... `fn_content`, respectively) in directory `root/d1/d2/.../dm`. Note that  $n \geq 1$  and  $m \geq 0$ . If  $m = 0$ , it means the directory is just the root directory.

The output is a list of group of duplicate file paths. For each group, it contains all the file paths of the files that have the same content. A file path is a string that has the following format:

```
"directory_path/file_name.txt"
```

Example 1:

Input:

```
["root/a 1.txt(abcd) 2.txt(efgh)", "root/c 3.txt(abcd)", "root/c/d 4.txt(efgh)",  
"root 4.txt(efgh)"]
```

Output:

```
[["root/a/2.txt","root/c/d/4.txt","root/4.txt"],["root/a/1.txt","root/c/3.txt"]]
```

Note:

No order is required for the final output.

You may assume the directory name, file name and file content only has letters and digits, and the length of file content is in the range of  $[1,50]$ .

The number of files given is in the range of  $[1,20000]$ .

You may assume no files or directories share the same name in the same directory.

You may assume each given directory info represents a unique directory.

Directory path and file info are separated by a single blank space.

Follow-up beyond contest:

Imagine you are given a real file system, how will you search files? DFS or BFS?

If the file content is very large (GB level), how will you modify your solution?

If you can only read the file by 1kb each time, how will you modify your solution?

What is the time complexity of your modified solution? What is the most time-consuming part and memory consuming part of it? How to optimize?

How to make sure the duplicated files you find are not false positive?

*Solution*

05/20/2020:

```

class Solution {
public:
    vector<vector<string>> findDuplicate(vector<string>& paths) {
        unordered_map<string, vector<string>> files;
        for (auto& p : paths) {
            istringstream iss(p);
            string base, file_content;
            iss >> base;
            while (iss >> file_content) {
                int i = 0, n = file_content.size();
                for (; file_content[i] != '(' && i < n; ++i);
                string filename = file_content.substr(0, i);
                string content = file_content.substr(i, n - i);
                files[content].push_back(base + "/" + filename);
            }
        }
        vector<vector<string>> ret;
        for (auto& f : files)
            if (f.second.size() > 1)
                ret.push_back(f.second);
        return ret;
    }
};

```

## 616. Add Bold Tag in String

### *Description*

Given a string *s* and a list of strings *dict*, you need to add a closed pair of bold tag `<b>` and `</b>` to wrap the substrings in *s* that exist in *dict*. If two such substrings overlap, you need to wrap them together by only one pair of closed bold tag. Also, if two substrings wrapped by bold tags are consecutive, you need to combine them.

Example 1:

Input:

*s* = "abcxyz123"

*dict* = ["abc", "123"]

Output:

"<b>abc</b>xyz<b>123</b>"

Example 2:

Input:

```
s = "aaabbcc"
dict = ["aaa","aab","bc"]
Output:
"<b>aaabbcc</b>c"
```

Constraints:

The given dict won't contain duplicates, and its length won't exceed 100.

All the strings in input have length in range [1, 1000].

Note: This question is the same as 758: <https://leetcode.com/problems/bold-words-in-string/>

*Solution*

06/09/2020:

```
class Solution {
public:
    string addBoldTag(string s, vector<string>& dict) {
        unordered_set<string> dicts(dict.begin(), dict.end());
        vector<int> lengths;
        for (auto& d : dict) lengths.push_back(d.size());
        sort(lengths.rbegin(), lengths.rend());
        int n = s.size();
        vector<bool> isBold(n, false);
        for (int i = 0; i < n; ++i) {
            bool wrap = false;
            int wrapLength = 0;
            for (auto& len : lengths) {
                if (i + len <= n && dicts.count(s.substr(i, len)) > 0) {
                    wrap = true;
                    wrapLength = len;
                    break;
                }
            }
            if (wrap) {
                for (int j = i; j < i + wrapLength; ++j)
                    isBold[j] = true;
            }
        }
        string ret;
        int last = 0;
        for (int i = 0; i < n; i++) {
            if (isBold[i]) {
                while (i + 1 < n && isBold[i + 1]) ++i;
                ret += "<b>";
                ret += s.substr(last, i - last + 1);
            }
            last = i + 1;
        }
        ret += s.substr(last, n - last);
        return ret;
    }
};
```

```

        ret += "</b>";
        last = ++i;
    } else {
        ret += s[i];
        last = ++i;
    }
}
return ret;
}
};

```

## 658. Find K Closest Elements

### Description

Given a sorted array, two integers  $k$  and  $x$ , find the  $k$  closest elements to  $x$  in the array. The result should also be sorted in ascending order. If there is a tie, the smaller elements are always preferred.

Example 1:

Input: [1,2,3,4,5],  $k=4$ ,  $x=3$

Output: [1,2,3,4]

Example 2:

Input: [1,2,3,4,5],  $k=4$ ,  $x=-1$

Output: [1,2,3,4]

Note:

The value  $k$  is positive and will always be smaller than the length of the sorted array.

Length of the given array is positive and will not exceed 104

Absolute value of elements in the array and  $x$  will not exceed 104

### Solution

04/28/2020:

```

class Solution {
public:
    vector<int> findClosestElements(vector<int>& arr, int k, int x) {
        if (arr.empty()) return {};
        int n = arr.size();
        int lo = 0, hi = n - 1;
        int l, r;
    }
};

```



```

if (x < arr.front()) {
    l = 0;
    r = 0;
} else if (x > arr.back()) {
    l = n - 1;
    r = n - 1;
} else {
    while (lo <= hi) {
        int mid = lo + (hi - lo) / 2;
        if (arr[mid] == x) {
            l = mid;
            r = mid;
            break;
        } else if (arr[mid] > x) {
            hi = mid - 1;
            r = hi;
        } else {
            lo = mid + 1;
            l = lo;
        }
    }
    if (abs(arr[l] - x) < abs(arr[r] - x)) {
        r = l;
    } else {
        l = r;
    }
}
deque<int> q;
if (k-- > 0) {
    q.push_front(arr[l]);
    l--;
    r++;
}
while (k > 0) {
    if (l < 0) {
        q.push_back(arr[r++]);
    } else if (r > n - 1) {
        q.push_front(arr[l--]);
    } else if (abs(arr[l] - x) <= abs(arr[r] - x)) {
        q.push_front(arr[l--]);
    } else if (abs(arr[l] - x) > abs(arr[r] - x)) {
        q.push_back(arr[r++]);
    }
    --k;
}
return vector<int>(q.begin(), q.end());
}
};

```

# 669. Trim a Binary Search Tree

## Description

Given a binary search tree and the lowest and highest boundaries as L and R, trim the tree so that all its elements lies in [L, R] ( $R \geq L$ ). You might need to change the root of the tree, so the result should return the new root of the trimmed binary search tree.

Example 1:

Input:

```
  1
 /  \
0    2
```

L = 1

R = 2

Output:

```
  1
   \
    2
```

Example 2:

Input:

```
  3
 /  \
0    4
   \
    2
   /
  1
```

L = 1

R = 3

Output:

```
  3
 /
 2
 /
 1
```

## Solution

05/10/2020:

```
/**
```

```

* Definition for a binary tree node.
* struct TreeNode {
*     int val;
*     TreeNode *left;
*     TreeNode *right;
*     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
* };
*/
class Solution {
public:
    TreeNode* trimBST(TreeNode* root, int L, int R) {
        if (root == nullptr) return root;
        if (root->val < L) return trimBST(root->right, L, R);
        if (root->val > R) return trimBST(root->left, L, R);
        root->left = trimBST(root->left, L, R);
        root->right = trimBST(root->right, L, R);
        return root;
    }
};

```

## 684. Redundant Connection

### Description

In this problem, a tree is an undirected graph that is connected and has no cycles.

The given input is a graph that started as a tree with  $N$  nodes (with distinct values  $1, 2, \dots, N$ ), with one additional edge added. The added edge has two different vertices chosen from  $1$  to  $N$ , and was not an edge that already existed.

The resulting graph is given as a 2D-array of edges. Each element of edges is a pair  $[u, v]$  with  $u < v$ , that represents an undirected edge connecting nodes  $u$  and  $v$ .

Return an edge that can be removed so that the resulting graph is a tree of  $N$  nodes. If there are multiple answers, return the answer that occurs last in the given 2D-array. The answer edge  $[u, v]$  should be in the same format, with  $u < v$ .

Example 1:

Input:  $[[1,2], [1,3], [2,3]]$

Output:  $[2,3]$

Explanation: The given undirected graph will be like this:

```

  1
 / \
2 - 3

```

Example 2:

Input: [[1,2], [2,3], [3,4], [1,4], [1,5]]

Output: [1,4]

Explanation: The given undirected graph will be like this:

```
5 - 1 - 2
    |   |
    4 - 3
```

Note:

The size of the input 2D-array will be between 3 and 1000.

Every integer represented in the 2D-array will be between 1 and N, where N is the size of the input array.

*Solution*

06/10/2020:

```
class UnionFind {
private:
    vector<int> id;
    vector<int> sz;

public:
    UnionFind(int n) {
        id.resize(n);
        iota(id.begin(), id.end(), 0);
        sz.resize(n, 1);
    }

    int find(int x) {
        if (x == id[x]) return x;
        return id[x] = find(id[x]);
    }

    bool connected(int x, int y) {
        return find(x) == find(y);
    }

    bool merge(int x, int y) {
        int i = find(x), j = find(y);
        if (i == j) return false;
        if (sz[i] > sz[j]) {
            sz[i] += sz[j];
            id[j] = i;
        } else {
            sz[j] += sz[i];
            id[i] = j;
        }
        return true;
    }
};
```

```

    }
};

class Solution {
public:
    vector<int> findRedundantConnection(vector<vector<int>>& edges) {
        int n = edges.size();
        UnionFind uf(n + 1);
        vector<int> ret;
        for (auto& e : edges)
            if (uf.connected(e[0], e[1]))
                ret = e;
            else
                uf.merge(e[0], e[1]);
        return ret;
    }
};

```

## 679. 24 Game

### *Description*

You have 4 cards each containing a number from 1 to 9. You need to judge whether they could be operated through  $*$ ,  $/$ ,  $+$ ,  $-$ ,  $($ ,  $)$  to get the value of 24.

Example 1:

Input: [4, 1, 8, 7]

Output: True

Explanation:  $(8-4) * (7-1) = 24$

Example 2:

Input: [1, 2, 1, 2]

Output: False

Note:

The division operator  $/$  represents real division, not integer division. For example,  $4 / (1 - 2/3) = 12$ .

Every operation done is between two numbers. In particular, we cannot use  $-$  as a unary operator. For example, with [1, 1, 1, 1] as input, the expression  $-1 - 1 - 1 - 1$  is not allowed.

You cannot concatenate numbers together. For example, if the input is [1, 2, 1, 2], we cannot write this as  $12 + 12$ .

### *Solution*

06/10/2020:

```

class Solution {
public:
    bool judgePoint24(vector<int>& nums) {
        unordered_set<string> seen;
        string operators = "+-*/";
        sort(nums.begin(), nums.end());
        do {
            int a = nums[0], b = nums[1], c = nums[2], d = nums[3];
            string key = to_string(a) + "," + to_string(b) + "," + to_string(c) + ","
+ to_string(d);
            if (seen.count(key) > 0) continue;
            seen.insert(key);
            for (int i = 0; i < 4; ++i) {
                for (int j = 0; j < 4; ++j) {
                    for (int k = 0; k < 4; ++k) {
                        vector<double> results;
                        results.push_back(op(op(op(a, b, operators[i]), c, operators[j]), d,
operators[k]));
                        results.push_back(op(op(a, b, operators[i]), op(c, d, operators[k]),
operators[j]));
                        results.push_back(op(op(a, op(b, c, operators[j]), operators[i]), d,
operators[k]));
                        results.push_back(op(a, op(op(b, c, operators[j]), d, operators[k]),
operators[i]));
                        results.push_back(op(a, op(b, op(c, d, operators[k]),
operators[j]), operators[i]));
                        for (auto& r : results)
                            if (fabs(r - 24) < 1e-6)
                                return true;
                    }
                }
            }
        } while (next_permutation(nums.begin(), nums.end()));
        return false;
    }

    double op(double a, double b, char op) {
        switch(op) {
            case '+': return a + b;
            case '-': return a - b;
            case '*': return a * b;
            case '/': return b != 0 ? a / b : -1e6;
        }
        return 0;
    }
};

```

```

class Solution {
public:
    bool judgePoint24(vector<int>& nums) {
        bool res = false;
        double eps = 0.001;
        vector<double> arr(nums.begin(), nums.end());
        dfs(arr, eps, res);
        return res;
    }

    void dfs(vector<double>& nums, double eps, bool& res) {
        if(res) return;
        if(nums.size()==1){
            if(abs(nums[0]-24)<eps) res=true;
            return;
        }
        for(int i=0;i<nums.size();i++){
            for(int j=0;j<i;j++){
                double p=nums[i], q=nums[j];
                vector<double>t{p+q, p-q, q-p, p*q};
                if (p > eps) t.push_back(q / p);
                if (q > eps) t.push_back(p / q);
                nums.erase(nums.begin() + i);
                nums.erase(nums.begin() + j);
                for (double d : t) {
                    nums.push_back(d);
                    dfs(nums, eps, res);
                    nums.pop_back();
                }
                nums.insert(nums.begin() + j, q);
                nums.insert(nums.begin() + i, p);
            }
        }
    }
};

```

## 693. Binary Number with Alternating Bits

### Description

Given a positive integer, check whether it has alternating bits: namely, if two adjacent bits will always have different values.

Example 1:

Input: 5  
Output: True  
Explanation:  
The binary representation of 5 is: 101  
Example 2:

Input: 7  
Output: False  
Explanation:  
The binary representation of 7 is: 111.  
Example 3:

Input: 11  
Output: False  
Explanation:  
The binary representation of 11 is: 1011.  
Example 4:

Input: 10  
Output: True  
Explanation:  
The binary representation of 10 is: 1010.

*Solution*

05/10/2020:

```
class Solution {
public:
    bool hasAlternatingBits(int n) {
        int cnt = n & 1 ? 1 : 0;
        while (n != 0) {
            n >>= 1;
            if (n & 1) {
                ++cnt;
            } else {
                --cnt;
            }
            if (cnt < 0 || cnt > 1) return false;
        }
        return true;
    }
};
```

## 694. Number of Distinct Islands

---



## Description

Given a non-empty 2D array grid of 0's and 1's, an island is a group of 1's (representing land) connected 4-directionally (horizontal or vertical.) You may assume all four edges of the grid are surrounded by water.

Count the number of distinct islands. An island is considered to be the same as another if and only if one island can be translated (and not rotated or reflected) to equal the other.

Example 1:

```
11000
11000
00011
00011
```

Given the above grid map, return 1.

Example 2:

```
11011
10000
00001
11011
```

Given the above grid map, return 3.

Notice that:

```
11
1
and
1
11
```

are considered different island shapes, because we do not consider reflection / rotation.

Note: The length of each dimension in the given grid does not exceed 50.

## Solution

05/08/2020 [Discussion](#):

1. Use a UnionFind set to find all the islands.
2. Eliminate the islands that are the "same" by hashing: translate all the islands to the top-left corner (subtract offset from the id for each cell). Then hash the sorted sequence of the cell. Note: the offset is obtained from the minimum of the row index  $min_i$  and column index  $min_j$ :  $offset = min_i * n + min_j$ .

```
class UnionFind {
private:
    vector<int> id;
```

```

vector<int> sz;

public:
UnionFind(int n) {
    id.resize(n);
    iota(id.begin(), id.end(), 0);
    sz.resize(n, 1);
}

int find(int x) {
    while (x != id[x]) {
        id[x] = id[id[x]];
        x = id[x];
    }
    return x;
}

bool merge(int x, int y) {
    int i = find(x), j = find(y);
    if (i == j) return false;
    if (sz[i] > sz[j]) {
        sz[i] += sz[j];
        id[j] = i;
    } else {
        sz[j] += sz[i];
        id[i] = j;
    }
    return true;
}
};

class Solution {
public:
    int numDistinctIslands(vector<vector<int>>& grid) {
        if (grid.empty() || grid[0].empty()) return 0;
        int m = grid.size(), n = grid[0].size();

        UnionFind uf(m * n);
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (grid[i][j] == 0) continue;
                int island = i * n + j;
                for (int d = 0; d < 4; ++d) {
                    int ni = i + dir[d][0], nj = j + dir[d][1], neighbor = ni * n + nj;
                    if (ni >= 0 && ni < m && nj >= 0 && nj < n && grid[ni][nj] == 1) {
                        bool merged = uf.merge(island, neighbor);
                    }
                }
            }
        }
    }
};

```

```

    }
}

unordered_map<int, vector<int>> components;
for (int i = 0; i < m; ++i) {
    for (int j = 0; j < n; ++j) {
        if (grid[i][j] == 0) continue;
        int island = i * n + j;
        components[uf.find(island)].push_back(island);
    }
}

unordered_set<int> islands;
for (auto& c : components) {
    sort(c.second.begin(), c.second.end());
    int mini = c.second.front() / n, minj = n;
    for (auto& i : c.second) minj = min(minj, i % n);
    int offset = mini * n + minj;
    long long hash = 1;
    const int MOD = 1e9 + 7;
    for (auto& i : c.second) hash = (hash * 31 + i - offset) % MOD;
    islands.insert(hash);
}
return islands.size();
}
};

```

```

class UnionFind {
private:
    vector<int> id;
    vector<int> sz;

public:
    UnionFind(int n) {
        id.resize(n);
        iota(id.begin(), id.end(), 0);
        sz.resize(n, 1);
    }

    int find(int x) {
        if (x == id[x]) return x;
        return id[x] = find(id[x]);
    }

    bool merge(int x, int y) {
        int i = find(x), j = find(y);
        if (i == j) return false;
        if (sz[i] > sz[j]) {

```

```

        sz[i] += sz[j];
        id[j] = i;
    } else {
        sz[j] += sz[i];
        id[i] = j;
    }
    return true;
}
};

class Solution {
private:
    int m, n;
public:
    int numDistinctIslands(vector<vector<int>>& grid) {
        if (grid.empty() || grid[0].empty()) return 0;
        m = grid.size(), n = grid[0].size();
        UnionFind uf(m * n);
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (grid[i][j] == 0) continue;
                int island = i * n + j;
                for (int d = 0; d < 4; ++d) {
                    int ni = i + dir[d][0], nj = j + dir[d][1], neighbor = ni * n + nj;
                    if (ni >= 0 && ni < m && nj >= 0 && nj < n && grid[ni][nj] == 1) {
                        uf.merge(island, neighbor);
                    }
                }
            }
        }

        unordered_map<int, vector<pair<int, int>>> components;
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (grid[i][j] == 0) continue;
                components[uf.find(i * n + j)].emplace_back(i, j);
            }
        }

        unordered_set<int> seen;
        int cnt = 0;
        for (auto& c : components) {
            vector<int> hashes = hash(c.second);
            bool exists = false;
            for (auto& h : hashes) {
                if (seen.count(h)) {
                    exists = true;
                    break;
                }
            }
        }
    }
};

```

```

    }
}
for (auto& h : hashes) seen.insert(h);
if (!exists) ++cnt;
}
return cnt;
}

vector<int> hash(vector<pair<int, int>> nums) {
    const int MOD = 1e9 + 7;
    int ops[4][2] = { {1, 1}, {-1, 1}, {-1, -1}, {1, -1} };
    int N = 1;
    vector<long long> hashes(N, 1);
    for (int d = 0; d < N; ++d) {
        int offset_x = INT_MAX, offset_y = INT_MAX;
        vector<pair<int, int>> t_nums(nums);
        for (auto& p : t_nums) {
            int& x = p.first;
            int& y = p.second;
            if (d >= 4) swap(x, y);
            x *= ops[d % 4][0];
            y *= ops[d % 4][1];
            offset_x = min(offset_x, x);
            offset_y = min(offset_y, y);
        }
        sort(t_nums.begin(), t_nums.end());
        for (auto& p : t_nums) {
            int& x = p.first;
            int& y = p.second;
            hashes[d] = (hashes[d] * 31 + (x - offset_x) * n + (y - offset_y)) %
MOD;
        }
    }
    return vector<int>(hashes.begin(), hashes.end());
}
};

```

## 700. Search in a Binary Search Tree

### Description

Given the root node of a binary search tree (BST) and a value. You need to find the node in the BST that the node's value equals the given value. Return the subtree rooted with that node. If such node doesn't exist, you should return NULL.

For example,

Given the tree:

```
    4
   / \
  2   7
 / \
1  3
```

And the value to search: 2

You should return this subtree:

```
    2
   / \
  1   3
```

In the example above, if we want to search the value 5, since there is no node with value 5, we should return NULL.

Note that an empty tree is represented by NULL, therefore you would see the expected output (serialized tree format) as [], notnull.

*Solution*

06/15/2020:

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
 * };
 */
class Solution {
public:
    TreeNode* searchBST(TreeNode* root, int val) {
        if (!root) return nullptr;
        if (root->val == val) return root;
        if (root->val > val) return searchBST(root->left, val);
        if (root->val < val) return searchBST(root->right, val);
        return nullptr;
    }
};
```

## 702. Search in a Sorted Array of Unknown Size

### Description

Given an integer array sorted in ascending order, write a function to search target in nums. If target exists, then return its index, otherwise return -1. However, the array size is unknown to you. You may only access the array using an `ArrayReader` interface, where `ArrayReader.get(k)` returns the element of the array at index `k` (0-indexed).

You may assume all integers in the array are less than 10000, and if you access the array out of bounds, `ArrayReader.get` will return 2147483647.

Example 1:

Input: array = [-1,0,3,5,9,12], target = 9

Output: 4

Explanation: 9 exists in nums and its index is 4

Example 2:

Input: array = [-1,0,3,5,9,12], target = 2

Output: -1

Explanation: 2 does not exist in nums so return -1

Note:

You may assume that all elements in the array are unique.

The value of each element in the array will be in the range [-9999, 9999].

### Solution

04/26/2020:

```
/**
 * // This is the ArrayReader's API interface.
 * // You should not implement it, or speculate about its implementation
 * class ArrayReader {
 *   public:
 *     int get(int index);
 * };
 */

class Solution {
public:
```

```

int search(const ArrayReader& reader, int target) {
    int lo = 0, hi = INT_MAX;
    while (lo <= hi) {
        int mid = lo + (hi - lo) / 2;
        int val = reader.get(mid);
        if (val == target) {
            return mid;
        } else if (val == INT_MAX || val > target) {
            hi = mid - 1;
        } else {
            lo = mid + 1;
        }
    }
    return -1;
}
};

```

## 705. Design HashSet

### *Description*

Design a HashSet without using any built-in hash table libraries.

To be specific, your design should include these functions:

add(value): Insert a value into the HashSet.

contains(value) : Return whether the value exists in the HashSet or not.

remove(value): Remove a value in the HashSet. If the value does not exist in the HashSet, do nothing.

Example:

```

MyHashSet hashSet = new MyHashSet();
hashSet.add(1);
hashSet.add(2);
hashSet.contains(1);    // returns true
hashSet.contains(3);    // returns false (not found)
hashSet.add(2);
hashSet.contains(2);    // returns true
hashSet.remove(2);
hashSet.contains(2);    // returns false (already removed)

```

Note:

All values will be in the range of [0, 1000000].

The number of operations will be in the range of [1, 10000].



Please do not use the built-in HashSet library.

*Solution*

05/10/2020:

```
struct Node {
    int key;
    Node* next;
    Node (int k) : key(k), next(nullptr) {}
};

class MyHashSet {
private:
    vector<Node*> hashset;
    int N;

public:
    /** Initialize your data structure here. */
    MyHashSet() {
        N = 100003;
        hashset.resize(N, nullptr);
    }

    void add(int key) {
        Node* cur = hashset[key % N];
        if (cur == nullptr) {
            hashset[key % N] = new Node(key);
        } else {
            while (cur->key != key && cur->next != nullptr) {
                cur = cur->next;
            }
            if (cur->key != key) {
                cur->next = new Node(key);
            }
        }
    }

    void remove(int key) {
        Node* cur = hashset[key % N];
        if (cur == nullptr) return;
        if (cur->key == key) {
            hashset[key % N] = cur->next;
        } else {
            while (cur->next != nullptr && cur->next->key != key) {
                cur = cur->next;
            }
            if (cur->next != nullptr && cur->next->key == key) {
```

```

        cur->next = cur->next->next;
    }
}

/** Returns true if this set contains the specified element */
bool contains(int key) {
    Node* cur = hashset[key % N];
    if (cur == nullptr) return false;
    if (cur->key == key) {
        return true;
    } else {
        while (cur->key != key && cur->next != nullptr) {
            cur = cur->next;
        }
        return cur->key == key;
    }
}
};

/**
 * Your MyHashSet object will be instantiated and called as such:
 * MyHashSet* obj = new MyHashSet();
 * obj->add(key);
 * obj->remove(key);
 * bool param_3 = obj->contains(key);
 */

```

## 706. Design HashMap

### *Description*

Design a HashMap without using any built-in hash table libraries.

To be specific, your design should include these functions:

put(key, value) : Insert a (key, value) pair into the HashMap. If the value already exists in the HashMap, update the value.

get(key): Returns the value to which the specified key is mapped, or -1 if this map contains no mapping for the key.

remove(key) : Remove the mapping for the value key if this map contains the mapping for the key.

Example:

```
MyHashMap hashMap = new MyHashMap();
```

```

hashMap.put(1, 1);
hashMap.put(2, 2);
hashMap.get(1);           // returns 1
hashMap.get(3);           // returns -1 (not found)
hashMap.put(2, 1);        // update the existing value
hashMap.get(2);           // returns 1
hashMap.remove(2);        // remove the mapping for 2
hashMap.get(2);           // returns -1 (not found)

```

Note:

All keys and values will be in the range of [0, 1000000].  
The number of operations will be in the range of [1, 10000].  
Please do not use the built-in HashMap library.

*Solution*

05/10/2020:

```

struct Node {
    int key, val;
    Node* next;
    Node (int k, int v) : key(k), val(v), next(nullptr) {}
};

class MyHashMap {
private:
    vector<Node*> hashmap;
    int N, n;

public:
    /** Initialize your data structure here. */
    MyHashMap() {
        N = 100003;
        n = 0;
        hashmap.resize(N, nullptr);
    }

    /** value will always be non-negative. */
    void put(int key, int value) {
        Node* cur = hashmap[key % N];
        if (cur == nullptr) {
            hashmap[key % N] = new Node(key, value);
        } else {
            while (cur->key != key && cur->next != nullptr) {
                cur = cur->next;
            }
            if (cur->key == key) {

```

```

        cur->val = value;
    } else if (cur->next == nullptr) {
        cur->next = new Node(key, value);
    }
}
}
}

```

/\*\* Returns the value to which the specified key is mapped, or -1 if this map contains no mapping for the key \*/

```

int get(int key) {
    Node* cur = hashmap[key % N];
    if (cur == nullptr) {
        return -1;
    } else if (cur->key == key) {
        return cur->val;
    } else {
        while (cur->key != key && cur->next != nullptr) {
            cur = cur->next;
        }
        return cur->key == key ? cur->val : -1;
    }
}
}

```

/\*\* Removes the mapping of the specified value key if this map contains a mapping for the key \*/

```

void remove(int key) {
    Node* cur = hashmap[key % N];
    if (cur == nullptr) return;
    if (cur->key == key) {
        hashmap[key % N] = cur->next;
    } else {
        while (cur->next != nullptr && cur->next->key != key) {
            cur = cur->next;
        }
        if (cur->next != nullptr && cur->next->key == key) {
            cur->next = cur->next->next;
        }
    }
}
}
};

```

/\*\*

\* Your MyHashMap object will be instantiated and called as such:

\* MyHashMap\* obj = new MyHashMap();

\* obj->put(key,value);

\* int param\_2 = obj->get(key);

\* obj->remove(key);

\*/

# 711. Number of Distinct Islands II

## Description

Given a non-empty 2D array grid of 0's and 1's, an island is a group of 1's (representing land) connected 4-directionally (horizontal or vertical.) You may assume all four edges of the grid are surrounded by water.

Count the number of distinct islands. An island is considered to be the same as another if they have the same shape, or have the same shape after rotation (90, 180, or 270 degrees only) or reflection (left/right direction or up/down direction).

Example 1:

```
11000
10000
00001
00011
```

Given the above grid map, return 1.

Notice that:

```
11
1
and
1
11
```

are considered same island shapes. Because if we make a 180 degrees clockwise rotation on the first island, then two islands will have the same shapes.

Example 2:

```
11100
10001
01001
01110
```

Given the above grid map, return 2.

Here are the two distinct islands:

```
111
1
and
1
1
```

Notice that:

```
111
1
and
```

1

111

are considered same island shapes. Because if we flip the first array in the up/down direction, then they have the same shapes.

Note: The length of each dimension in the given grid does not exceed 50.

*Solution*

05/08/2020 [Discussion](#):

1. Use a UnionFind set to find all the islands.
2. Eliminate the islands that are the "same" by hashing: translate all the transformed islands to the top-left corner (subtract `offset_x` from the `x` and subtract `offset_y` from `y` for each cell). Then hash the sorted sequence of the cell. Note: the `offset_x` and `offset_y` are obtained from the minimum of the row index and column index.

```
class UnionFind {
private:
    vector<int> id;
    vector<int> sz;

public:
    UnionFind(int n) {
        id.resize(n);
        iota(id.begin(), id.end(), 0);
        sz.resize(n, 1);
    }

    int find(int x) {
        if (x == id[x]) return x;
        return id[x] = find(id[x]);
    }

    bool merge(int x, int y) {
        int i = find(x), j = find(y);
        if (i == j) return false;
        if (sz[i] > sz[j]) {
            sz[i] += sz[j];
            id[j] = i;
        } else {
            sz[j] += sz[i];
            id[i] = j;
        }
        return true;
    }
};

class Solution {
```

```

private:
    int m, n;
public:
    int numDistinctIslands2(vector<vector<int>>& grid) {
        if (grid.empty() || grid[0].empty()) return 0;
        m = grid.size(), n = grid[0].size();
        UnionFind uf(m * n);
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (grid[i][j] == 0) continue;
                int island = i * n + j;
                for (int d = 0; d < 4; ++d) {
                    int ni = i + dir[d][0], nj = j + dir[d][1], neighbor = ni * n + nj;
                    if (ni >= 0 && ni < m && nj >= 0 && nj < n && grid[ni][nj] == 1) {
                        uf.merge(island, neighbor);
                    }
                }
            }
        }

        unordered_map<int, vector<pair<int, int>>> components;
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (grid[i][j] == 0) continue;
                components[uf.find(i * n + j)].emplace_back(i, j);
            }
        }

        unordered_set<int> seen;
        int cnt = 0;
        for (auto& c : components) {
            vector<int> hashes = hash(c.second);
            bool exists = false;
            for (auto& h : hashes) {
                if (seen.count(h)) {
                    exists = true;
                    break;
                }
            }
            for (auto& h : hashes) seen.insert(h);
            if (!exists) ++cnt;
        }
        return cnt;
    }

    vector<int> hash(vector<pair<int, int>> nums) {
        const int MOD = 1e9 + 7;
        int ops[4][2] = { {1, 1}, {-1, 1}, {-1, -1}, {1, -1} };
    }

```

```

vector<long long> hashes(8, 1);
for (int d = 0; d < 8; ++d) {
    int offset_x = INT_MAX, offset_y = INT_MAX;
    vector<pair<int, int>> t_nums(nums);
    for (auto& p : t_nums) {
        int& x = p.first;
        int& y = p.second;
        if (d >= 4) swap(x, y);
        x *= ops[d % 4][0];
        y *= ops[d % 4][1];
        offset_x = min(offset_x, x);
        offset_y = min(offset_y, y);
    }
    sort(t_nums.begin(), t_nums.end());
    for (auto& p : t_nums) {
        int& x = p.first;
        int& y = p.second;
        hashes[d] = (hashes[d] * 31 + (x - offset_x) * n + (y - offset_y)) %
MOD;
    }
}
return vector<int>(hashes.begin(), hashes.end());
};

```

## 717. 1-bit and 2-bit Characters

### Description

We have two special characters. The first character can be represented by one bit 0. The second character can be represented by two bits (10 or 11).

Now given a string represented by several bits. Return whether the last character must be a one-bit character or not. The given string will always end with a zero.

Example 1:

Input:

bits = [1, 0, 0]

Output: True

Explanation:

The only way to decode it is two-bit character and one-bit character. So the last character is one-bit character.

Example 2:

Input:

bits = [1, 1, 1, 0]



Output: False

Explanation:

The only way to decode it is two-bit character and two-bit character. So the last character is NOT one-bit character.

Note:

```
1 <= len(bits) <= 1000.  
bits[i] is always 0 or 1.
```

*Solution*

06/09/2020:

```
class Solution {  
public:  
    bool isOneBitCharacter(vector<int>& bits) {  
        int i = 0, n = bits.size();  
        for (; i < n - 1; ++i) {  
            if (bits[i] == 1) {  
                ++i;  
            }  
        }  
        return i == n - 1;  
    }  
};
```

## 720. Longest Word in Dictionary

*Description*

Given a list of strings words representing an English Dictionary, find the longest word in words that can be built one character at a time by other words in words. If there is more than one possible answer, return the longest word with the smallest lexicographical order.

If there is no answer, return the empty string.

Example 1:

Input:

```
words = ["w","wo","wor","worl", "world"]
```

Output: "world"

Explanation:

The word "world" can be built one character at a time by "w", "wo", "wor", and "worl".

Example 2:

Input:

```
words = ["a", "banana", "app", "appl", "ap", "apply", "apple"]
```

Output: "apple"

Explanation:

Both "apply" and "apple" can be built from other words in the dictionary.

However, "apple" is lexicographically smaller than "apply".

Note:

All the strings in the input will only contain lowercase letters.

The length of words will be in the range [1, 1000].

The length of words[i] will be in the range [1, 30].

*Solution*

05/18/2020:

```
class Solution {
public:
    string longestWord(vector<string>& words) {
        unordered_set<string> s(words.begin(), words.end());
        sort(words.begin(), words.end(), [](const string& s1, const string& s2) {
            if (s1.size() == s2.size()) {
                return s1 < s2;
            }
            return s1.size() > s2.size();
        });
        for (auto& w : words) {
            bool ok = true;
            for (int i = 1; i < (int)w.size(); ++i) {
                if (s.count(w.substr(0, i)) == 0) {
                    ok = false;
                    break;
                }
            }
            if (ok) return w;
        }
        return "";
    }
};
```

## 721. Accounts Merge

*Description*

Given a list accounts, each element accounts[i] is a list of strings, where the first element accounts[i][0] is a name, and the rest of the elements are emails representing emails of the account.

Now, we would like to merge these accounts. Two accounts definitely belong to the same person if there is some email that is common to both accounts. Note that even if two accounts have the same name, they may belong to different people as people could have the same name. A person can have any number of accounts initially, but all of their accounts definitely have the same name.

After merging the accounts, return the accounts in the following format: the first element of each account is the name, and the rest of the elements are emails in sorted order. The accounts themselves can be returned in any order.

Example 1:

Input:

```
accounts = [{"John", "johnsmith@mail.com", "john00@mail.com"}, {"John", "johnnybravo@mail.com"}, {"John", "johnsmith@mail.com", "john_newyork@mail.com"}, {"Mary", "mary@mail.com"}]
```

```
Output: [{"John", "john00@mail.com", "john_newyork@mail.com", "johnsmith@mail.com"}, {"John", "johnnybravo@mail.com"}, {"Mary", "mary@mail.com"}]
```

Explanation:

The first and third John's are the same person as they have the common email "johnsmith@mail.com".

The second John and Mary are different people as none of their email addresses are used by other accounts.

We could return these lists in any order, for example the answer [{"Mary", "mary@mail.com"}, {"John", "johnnybravo@mail.com"}, {"John", "john00@mail.com", "john\_newyork@mail.com", "johnsmith@mail.com"}] would still be accepted.

Note:

The length of accounts will be in the range [1, 1000].

The length of accounts[i] will be in the range [1, 10].

The length of accounts[i][j] will be in the range [1, 30].

*Solution*

06/10/2020:

Define UnionFind:

```
class UnionFind {
private:
    vector<int> id;
    vector<int> sz;
```

```

public:
    UnionFind(int n) {
        id.resize(n);
        iota(id.begin(), id.end(), 0);
        sz.resize(n, 1);
    }

    int find(int x) {
        if (x == id[x]) return x;
        return id[x] = find(id[x]);
    }

    bool connected(int x, int y) {
        return find(x) == find(y);
    }

    bool merge(int x, int y) {
        int i = find(x), j = find(y);
        if (i == j) return false;
        if (sz[i] > sz[j]) {
            id[j] = i;
            sz[i] += sz[j];
        } else {
            id[i] = j;
            sz[j] += sz[i];
        }
        return true;
    }
};

```

```

class Solution {
public:
    vector<vector<string>> accountsMerge(vector<vector<string>>& accounts) {
        int id = 0;
        UnionFind uf(10000);
        unordered_map<string, string> email_user;
        unordered_map<string, int> email_id;
        unordered_map<int, string> id_email;
        for (auto& emails : accounts) {
            for (int i = 1; i < (int)emails.size(); ++i) {
                email_user[emails[i]] = emails[0];
                if (email_id.count(emails[i]) == 0) {
                    email_id[emails[i]] = id++;
                    id_email[id - 1] = emails[i];
                }
                if (i > 1) uf.merge(email_id[emails[i - 1]], email_id[emails[i]]);
            }
        }
    }
};

```

```

unordered_map<int, vector<int>> connected_components;
for (int i = 0; i < id; ++i) connected_components[uf.find(i)].push_back(i);

vector<vector<string>> ret;
for (auto& c : connected_components) {
    vector<string> account;
    account.push_back(email_user[id_email[c.first]]);
    for (auto& i : c.second) account.push_back(id_email[i]);
    sort(account.begin() + 1, account.end());
    ret.push_back(account);
}
return ret;
}
};

```

```

class Solution {
public:
    vector<vector<string>> accountsMerge(vector<vector<string>>& accounts) {
        int n = accounts.size();
        UnionFind uf(n);
        for (int i = 0; i < n; ++i) {
            unordered_set<string> accounts_set(accounts[i].begin() + 1,
accounts[i].end());
            for (int j = i + 1; j < n; ++j)
                if (accounts[i][0] == accounts[j][0])
                    for (int k = 1; !uf.connected(i, j) && k < (int)accounts[j].size();
++k)
                        if (accounts_set.count(accounts[j][k]) > 0)
                            uf.merge(i, j);
        }

        unordered_map<int, vector<int>> connected_components;
        for (int i = 0; i < n; ++i)
            if (connected_components.count(uf.find(i)) == 0)
                connected_components[uf.find(i)] = {i};
            else
                connected_components[uf.find(i)].push_back(i);

        vector<vector<string>> ret;
        for (auto& c : connected_components) {
            unordered_set<string> acct;
            for (auto i : c.second)
                for (int j = 1; j < (int)accounts[i].size(); ++j)
                    acct.insert(accounts[i][j]);
            vector<string> sorted_acct(acct.begin(), acct.end());
            sort(sorted_acct.rbegin(), sorted_acct.rend());
            sorted_acct.push_back(accounts[c.first][0]);
        }
    }
};

```

```
        reverse(sorted_acct.begin(), sorted_acct.end());
        ret.push_back(sorted_acct);
    }
    return ret;
}
};
```

## 733. Flood Fill

### *Description*

An image is represented by a 2-D array of integers, each integer representing the pixel value of the image (from 0 to 65535).

Given a coordinate (sr, sc) representing the starting pixel (row and column) of the flood fill, and a pixel value newColor, "flood fill" the image.

To perform a "flood fill", consider the starting pixel, plus any pixels connected 4-directionally to the starting pixel of the same color as the starting pixel, plus any pixels connected 4-directionally to those pixels (also with the same color as the starting pixel), and so on. Replace the color of all of the aforementioned pixels with the newColor.

At the end, return the modified image.

Example 1:

Input:

```
image = [[1,1,1],[1,1,0],[1,0,1]]
```

```
sr = 1, sc = 1, newColor = 2
```

```
Output: [[2,2,2],[2,2,0],[2,0,1]]
```

Explanation:

From the center of the image (with position (sr, sc) = (1, 1)), all pixels connected

by a path of the same color as the starting pixel are colored with the new color.

Note the bottom corner is not colored 2, because it is not 4-directionally connected to the starting pixel.

Note:

The length of image and image[0] will be in the range [1, 50].

The given starting pixel will satisfy  $0 \leq sr < \text{image.length}$  and  $0 \leq sc < \text{image}[0].\text{length}$ .

The value of each color in image[i][j] and newColor will be an integer in [0, 65535].

05/11/2020:

```

class Solution {
public:
    vector<vector<int>> floodFill(vector<vector<int>>& image, int sr, int sc, int
newColor) {
        if (image.empty() || image[0].empty()) return image;
        if (image[sr][sc] == newColor) return image;
        int m = image.size(), n = image[0].size();
        int oldColor = image[sr][sc];
        image[sr][sc] = newColor;
        queue<pair<int, int>> q;
        q.emplace(sr, sc);
        int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
        while (!q.empty()) {
            int sz = q.size();
            for (int s = 0; s < sz; ++s) {
                pair<int, int> cur = q.front(); q.pop();
                int i = cur.first, j = cur.second;
                for (int d = 0; d < 4; ++d) {
                    int ni = i + dir[d][0], nj = j + dir[d][1];
                    if (ni >= 0 && ni < m && nj >= 0 && nj < n && image[ni][nj] ==
oldColor) {
                        image[ni][nj] = newColor;
                        q.emplace(ni, nj);
                    }
                }
            }
        }
        return image;
    }
};

```

## 734. Sentence Similarity

### Description

Given two sentences `words1`, `words2` (each represented as an array of strings), and a list of similar word pairs `pairs`, determine if two sentences are similar.

For example, "great acting skills" and "fine drama talent" are similar, if the similar word pairs are `pairs = [{"great", "fine"}, {"acting", "drama"}, {"skills", "talent"}]`.

Note that the similarity relation is not transitive. For example, if "great" and "fine" are similar, and "fine" and "good" are similar, "great" and "good" are not necessarily similar.

However, similarity is symmetric. For example, "great" and "fine" being similar is the same as "fine" and "great" being similar.

Also, a word is always similar with itself. For example, the sentences words1 = ["great"], words2 = ["great"], pairs = [] are similar, even though there are no specified similar word pairs.

Finally, sentences can only be similar if they have the same number of words. So a sentence like words1 = ["great"] can never be similar to words2 = ["doubleplus", "good"].

Note:

The length of words1 and words2 will not exceed 1000.

The length of pairs will not exceed 2000.

The length of each pairs[i] will be 2.

The length of each words[i] and pairs[i][j] will be in the range [1, 20].

*Solution*

06/10/2020:

```
class Solution {
public:
    bool areSentencesSimilar(vector<string>& words1, vector<string>& words2,
vector<vector<string>>& pairs) {
        if (words1.size() != words2.size()) return false;
        bool isSimilar = true;
        unordered_map<string, unordered_set<string>> similarWords;
        for (auto& p : pairs) {
            similarWords[p[0]].insert(p[1]);
            similarWords[p[1]].insert(p[0]);
        }
        for (int i = 0; i < words1.size(); ++i)
            if (words1[i] != words2[i])
                if (similarWords[words1[i]].count(words2[i]) == 0 &&
similarWords[words2[i]].count(words1[i]) == 0)
                    return false;
        return true;
    }
};
```

05/20/2020:



```

class Solution {
public:
    bool areSentencesSimilar(vector<string>& words1, vector<string>& words2,
vector<vector<string>>& pairs) {
        if (words1.empty() && words2.empty()) return true;
        if (words1.size() != words2.size()) return false;
        unordered_map<string, unordered_set<string>> similarWords;
        for (auto& p : pairs) {
            if (p.empty()) continue;
            similarWords[p[0]].insert(p[1]);
            similarWords[p[1]].insert(p[0]);
        }
        int n = words1.size();
        for (int i = 0; i < n; ++i) {
            bool isSimilar = false;
            if (words1[i] == words2[i]) {
                isSimilar = true;
            } else {
                if (similarWords[words1[i]].count(words2[i]) > 0)
                    isSimilar = true;
            }
            if (isSimilar == false) return false;
        }
        return true;
    }
};

```

## 737. Sentence Similarity II

### *Description*

Given two sentences words1, words2 (each represented as an array of strings), and a list of similar word pairs pairs, determine if two sentences are similar.

For example, words1 = ["great", "acting", "skills"] and words2 = ["fine", "drama", "talent"] are similar, if the similar word pairs are pairs = [{"great", "good"}, {"fine", "good"}, {"acting", "drama"}, {"skills", "talent"}].

Note that the similarity relation is transitive. For example, if "great" and "good" are similar, and "fine" and "good" are similar, then "great" and "fine" are similar.

Similarity is also symmetric. For example, "great" and "fine" being similar is the same as "fine" and "great" being similar.

Also, a word is always similar with itself. For example, the sentences words1 = ["great"], words2 = ["great"], pairs = [] are similar, even though there are no specified similar word pairs.

Finally, sentences can only be similar if they have the same number of words. So a sentence like words1 = ["great"] can never be similar to words2 = ["doubleplus", "good"].

Note:

The length of words1 and words2 will not exceed 1000.

The length of pairs will not exceed 2000.

The length of each pairs[i] will be 2.

The length of each words[i] and pairs[i][j] will be in the range [1, 20].

*Solution*

06/10/2020:

```
class UnionFind {
private:
    vector<int> id;
    vector<int> sz;

public:
    UnionFind(int n) {
        id.resize(n);
        iota(id.begin(), id.end(), 0);
        sz.resize(n, 1);
    }

    int find(int x) {
        if (x == id[x]) return x;
        return id[x] = find(id[x]);
    }

    bool connected(int x, int y) {
        return find(x) == find(y);
    }

    bool merge(int x, int y) {
        int i = find(x), j = find(y);
        if (i == j) return false;
        if (sz[i] > sz[j]) {
            sz[i] += sz[j];
            id[j] = i;
        } else {
            sz[j] += sz[i];
            id[i] = j;
        }
    }
};
```

```

        id[i] = j;
    }
    return true;
}
};

class Solution {
public:
    bool areSentencesSimilarTwo(vector<string>& words1, vector<string>& words2,
vector<vector<string>>& pairs) {
        if (words1.size() != words2.size()) return false;
        UnionFind uf(6000);
        unordered_map<string, int> word_id;
        int id = 0;
        for (auto& w : words1) assign_id(w, id, word_id);
        for (auto& w : words2) assign_id(w, id, word_id);
        for (auto& p : pairs) {
            assign_id(p[0], id, word_id);
            assign_id(p[1], id, word_id);
            uf.merge(word_id[p[0]], word_id[p[1]]);
        }
        for (int i = 0; i < (int)words1.size(); ++i)
            if (words1[i] != words2[i] && !uf.connected(word_id[words1[i]],
word_id[words2[i]]))
                return false;
        return true;
    }

    void assign_id(const string& word, int& id, unordered_map<string, int>&
word_id) {
        if (word_id.count(word) == 0) word_id[word] = id++;
    }
};

```

## 748. Shortest Completing Word

### Description

Find the minimum length word from a given dictionary words, which has all the letters from the string licensePlate. Such a word is said to complete the given string licensePlate

Here, for letters we ignore case. For example, "P" on the licensePlate still matches "p" on the word.

It is guaranteed an answer exists. If there are multiple answers, return the one that occurs first in the array.

The license plate might have the same letter occurring multiple times. For example, given a licensePlate of "PP", the word "pair" does not complete the licensePlate, but the word "supper" does.

Example 1:

Input: licensePlate = "1s3 PSt", words = ["step", "steps", "stripe", "stepple"]

Output: "steps"

Explanation: The smallest length word that contains the letters "S", "P", "S", and "T".

Note that the answer is not "step", because the letter "s" must occur in the word twice.

Also note that we ignored case for the purposes of comparing whether a letter exists in the word.

Example 2:

Input: licensePlate = "1s3 456", words = ["looks", "pest", "stew", "show"]

Output: "pest"

Explanation: There are 3 smallest length words that contains the letters "s". We return the one that occurred first.

Note:

licensePlate will be a string with length in range [1, 7].

licensePlate will contain digits, spaces, or letters (uppercase or lowercase).

words will have a length in the range [10, 1000].

Every words[i] will consist of lowercase letters, and have length in range [1, 15].

*Solution*

05/17/2020:

```
class Solution {
public:
    string shortestCompletingWord(string licensePlate, vector<string>& words) {
        vector<int> trimmedLicensePlate(26, 0);
        for (auto& l : licensePlate) {
            if (isalpha(l)) {
                ++trimmedLicensePlate[tolower(l) - 'a'];
            }
        }
        vector<pair<string, int>> validWords;
        for (int k = 0; k < (int)words.size(); ++k) {
            string w = words[k];
            vector<int> cnt(26, 0);
            for (auto& c : w) {
                ++cnt[c - 'a'];
            }
        }
    }
};
```

```

bool isComplete = true;
for (int i = 0; i < 26; ++i) {
    if (cnt[i] < trimmedLicensePlate[i]) {
        isComplete = false;
        break;
    }
}
if (isComplete) validWords.emplace_back(w, k);
}
sort(validWords.begin(), validWords.end(), [](pair<string, int>& p1,
pair<string, int>& p2) {
    if (p1.first.size() == p2.first.size()) {
        return p1.second < p2.second;
    }
    return p1.first.size() < p2.first.size();
});
return validWords.front().first;
}
};

```

## 758. Bold Words in String

### Description

Given a set of keywords `words` and a string `S`, make all appearances of all keywords in `S` bold. Any letters between `<b>` and `</b>` tags become bold.

The returned string should use the least number of tags possible, and of course the tags should form a valid combination.

For example, given that `words = ["ab", "bc"]` and `S = "aabcd"`, we should return `"a<b>abc</b>d"`. Note that returning `"a<b>a<b>b</b>c</b>d"` would use more tags, so it is incorrect.

Constraints:

`words` has length in range `[0, 50]`.

`words[i]` has length in range `[1, 10]`.

`S` has length in range `[0, 500]`.

All characters in `words[i]` and `S` are lowercase letters.

Note: This question is the same as 616: <https://leetcode.com/problems/add-bold-tag-in-string/>

### Solution

06/09/2020:

```
class Solution {
public:
    string boldWords(vector<string>& dict, string s) {
        unordered_set<string> dicts(dict.begin(), dict.end());
        vector<int> lengths;
        for (auto& d : dict) lengths.push_back(d.size());
        sort(lengths.rbegin(), lengths.rend());
        int n = s.size();
        vector<bool> isBold(n, false);
        for (int i = 0; i < n; ++i) {
            bool wrap = false;
            int wrapLength = 0;
            for (auto& len : lengths) {
                if (i + len <= n && dicts.count(s.substr(i, len)) > 0) {
                    wrap = true;
                    wrapLength = len;
                    break;
                }
            }
            if (wrap) {
                for (int j = i; j < i + wrapLength; ++j)
                    isBold[j] = true;
            }
        }
        string ret;
        int last = 0;
        for (int i = 0; i < n;) {
            if (isBold[i]) {
                while (i + 1 < n && isBold[i + 1]) ++i;
                ret += "<b>";
                ret += s.substr(last, i - last + 1);
                ret += "</b>";
                last = ++i;
            } else {
                ret += s[i];
                last = ++i;
            }
        }
        return ret;
    }
};
```

## 773. Sliding Puzzle

### Description

On a 2x3 board, there are 5 tiles represented by the integers 1 through 5, and an empty square represented by 0.

A move consists of choosing 0 and a 4-directionally adjacent number and swapping it.

The state of the board is solved if and only if the board is `[[1,2,3],[4,5,0]]`.

Given a puzzle board, return the least number of moves required so that the state of the board is solved. If it is impossible for the state of the board to be solved, return -1.

Examples:

Input: board = `[[1,2,3],[4,0,5]]`

Output: 1

Explanation: Swap the 0 and the 5 in one move.

Input: board = `[[1,2,3],[5,4,0]]`

Output: -1

Explanation: No number of moves will make the board solved.

Input: board = `[[4,1,2],[5,0,3]]`

Output: 5

Explanation: 5 is the smallest number of moves that solves the board.

An example path:

After move 0: `[[4,1,2],[5,0,3]]`

After move 1: `[[4,1,2],[0,5,3]]`

After move 2: `[[0,1,2],[4,5,3]]`

After move 3: `[[1,0,2],[4,5,3]]`

After move 4: `[[1,2,0],[4,5,3]]`

After move 5: `[[1,2,3],[4,5,0]]`

Input: board = `[[3,2,4],[1,5,0]]`

Output: 14

Note:

board will be a 2 x 3 array as described above.

board[i][j] will be a permutation of [0, 1, 2, 3, 4, 5].

### Solution

05/08/2020 [Discussion](#):

This is the C++ version solution for one of the assignments of the course [Algorithms \(Part 1\)](#):

```

class Board {
public:
    vector<vector<int>> board;
    pair<int, int> zeroTile;
    int rows, cols, manhattanDist, moves, priority;

    Board(const vector<vector<int>>& board) {
        this->board = board;
        rows = board.size();
        cols = board[0].size();
        zeroTile = findZeroTile();
        manhattanDist = manhattan();
        moves = 0;
    }

    bool operator<(const Board& that) const { return this->priority <
that.priority; }

    bool operator>(const Board& that) const { return this->priority >
that.priority; }

    bool operator!=(const Board& that) const {
        for (int i = 0; i < rows; ++i)
            for (int j = 0; j < cols; ++j)
                if (that.board[i][j] != board[i][j])
                    return true;
        return false;
    }

    pair<int, int> findZeroTile() {
        for (int i = 0; i < rows; ++i)
            for (int j = 0; j < cols; ++j)
                if (board[i][j] == 0)
                    return {i, j};
        return {-1, -1};
    }

    int manhattan() {
        int d = 0;
        for (int i = 0; i < rows; ++i)
            for (int j = 0; j < cols; ++j)
                if (board[i][j] != 0)
                    d += abs((board[i][j] - 1) / cols - i) + abs((board[i][j] - 1) % cols
- j);
        return d;
    }

    void updateMoves(int moves) {

```



```

    this->moves = moves;
    priority = manhattanDist + moves;
}

vector<Board> neighbor() {
    int dir[4][2] = { {-1, 0}, {1, 0}, {0, -1}, {0, 1} };
    int i = zeroTile.first, j = zeroTile.second;
    vector<Board> neighborBoards;
    for (int d = 0; d < 4; ++d) {
        int ni = dir[d][0] + i, nj = dir[d][1] + j;
        vector<vector<int>> newBoard(board);
        if (ni >= 0 && ni < rows && nj >= 0 && nj < cols) {
            swap(newBoard[i][j], newBoard[ni][nj]);
            neighborBoards.push_back(Board(newBoard));
            neighborBoards.back().updateMoves(moves + 1);
        }
    }
    return neighborBoards;
}

Board twin () {
    for (int i = 0; i < rows; ++i) {
        for (int j = 0; j < cols - 1; ++j) {
            if (board[i][j] != 0 && board[i][j + 1] != 0) {
                vector<vector<int>> twinBoard(board);
                swap(twinBoard[i][j], twinBoard[i][j + 1]);
                return Board(twinBoard);
            }
        }
    }
    return Board(board);
}

};

class Solution {
public:
    int slidingPuzzle(vector<vector<int>>& board) {
        priority_queue<pair<Board, Board>, vector<pair<Board, Board>>,
greater<pair<Board, Board>>> pq, pqTwin;
        Board initialBoard(board), twinBoard = initialBoard.twin();
        pq.emplace(initialBoard, initialBoard);
        pqTwin.emplace(twinBoard, twinBoard);
        while (!pq.empty() && pq.top().first.manhattanDist > 0 && !pqTwin.empty() &&
pqTwin.top().first.manhattanDist > 0) {
            pair<Board, Board> cur = pq.top(); pq.pop();
            for(auto& p : cur.first.neighbor())
                if (p != cur.second) pq.emplace(p, cur.first);
            cur = pqTwin.top(); pqTwin.pop();
            for(auto& p : cur.first.neighbor())

```

```

        if (p != cur.second) pqTwin.emplace(p, cur.first);
    }
    return pqTwin.top().first.manhattanDist == 0 ? -1 : pq.top().first.moves;
}
};

```

## 787. Cheapest Flights Within K Stops

### Description

There are  $n$  cities connected by  $m$  flights. Each flight starts from city  $u$  and arrives at  $v$  with a price  $w$ .

Now given all the cities and flights, together with starting city  $src$  and the destination  $dst$ , your task is to find the cheapest price from  $src$  to  $dst$  with up to  $k$  stops. If there is no such route, output  $-1$ .

Example 1:

Input:

$n = 3$ ,  $edges = [[0,1,100],[1,2,100],[0,2,500]]$

$src = 0$ ,  $dst = 2$ ,  $k = 1$

Output: 200

Explanation:

The graph looks like this:

The cheapest price from city 0 to city 2 with at most 1 stop costs 200, as marked red in the picture.

Example 2:

Input:

$n = 3$ ,  $edges = [[0,1,100],[1,2,100],[0,2,500]]$

$src = 0$ ,  $dst = 2$ ,  $k = 0$

Output: 500

Explanation:

The graph looks like this:

The cheapest price from city 0 to city 2 with at most 0 stop costs 500, as marked blue in the picture.

Constraints:

The number of nodes  $n$  will be in range  $[1, 100]$ , with nodes labeled from 0 to  $n - 1$ .

The size of flights will be in range  $[0, n * (n - 1) / 2]$ .

The format of each flight will be (src, dst, price).  
The price of each flight will be in the range [1, 10000].  
k is in the range of [0, n - 1].  
There will not be any duplicated flights or self cycles.

*Solution*

06/14/2020:

```
typedef tuple<int, int, int> ti;
class Solution {
public:
    int findCheapestPrice(int n, vector<vector<int>>& flights, int src, int dst,
int K) {
        vector<vector<pair<int, int>>> adj(n);
        for(const auto& f : flights) adj[f[0]].emplace_back(f[1], f[2]);
        priority_queue<ti, vector<ti>, greater<ti>> pq;
        pq.emplace(0, src, K + 1);
        while(!pq.empty()) {
            auto [cost, u, stops] = pq.top(); pq.pop();
            if(u == dst) return cost;
            if(!stops) continue;
            for(auto& to : adj[u]) {
                auto [v, w] = to;
                pq.emplace(cost + w, v, stops - 1);
            }
        }
        return -1;
    }
};
```

## 819. Most Common Word

*Description*

Given a paragraph and a list of banned words, return the most frequent word that is not in the list of banned words. It is guaranteed there is at least one word that isn't banned, and that the answer is unique.

Words in the list of banned words are given in lowercase, and free of punctuation. Words in the paragraph are not case sensitive. The answer is in lowercase.

Example:

Input:

```
paragraph = "Bob hit a ball, the hit BALL flew far after it was hit."
```

```
banned = ["hit"]
```

Output: "ball"

Explanation:

"hit" occurs 3 times, but it is a banned word.

"ball" occurs twice (and no other word does), so it is the most frequent non-banned word in the paragraph.

Note that words in the paragraph are not case sensitive,

that punctuation is ignored (even if adjacent to words, such as "ball,"),

and that "hit" isn't the answer even though it occurs more because it is banned.

Note:

```
1 <= paragraph.length <= 1000.
```

```
0 <= banned.length <= 100.
```

```
1 <= banned[i].length <= 10.
```

The answer is unique, and written in lowercase (even if its occurrences in paragraph may have uppercase symbols, and even if it is a proper noun.)

paragraph only consists of letters, spaces, or the punctuation symbols !?',,;. There are no hyphens or hyphenated words.

There are no hyphens or hyphenated words.

Words only consist of letters, never apostrophes or other punctuation symbols.

*Solution*

05/10/2020:

```
class Solution {
public:
    string mostCommonWord(string paragraph, vector<string>& banned) {
        for (auto& c : paragraph) c = isalpha(c) ? tolower(c) : ' ';
        unordered_set<string> b(banned.begin(), banned.end());
        unordered_map<string, int> mp;
        string ret, s;
        int freq = 0;
        istringstream iss(paragraph);
        while (iss >> s) {
            if (b.count(s) == 0 && ++mp[s] > freq) {
                freq = mp[s];
                ret = s;
            }
        }
        return ret;
    }
};
```

## 844. Backspace String Compare

### Description

Given two strings S and T, return if they are equal when both are typed into empty text editors. # means a backspace character.

Example 1:

Input: S = "ab#c", T = "ad#c"

Output: true

Explanation: Both S and T become "ac".

Example 2:

Input: S = "ab##", T = "c#d#"

Output: true

Explanation: Both S and T become "".

Example 3:

Input: S = "a##c", T = "#a#c"

Output: true

Explanation: Both S and T become "c".

Example 4:

Input: S = "a#c", T = "b"

Output: false

Explanation: S becomes "c" while T becomes "b".

Note:

1 <= S.length <= 200

1 <= T.length <= 200

S and T only contain lowercase letters and '#' characters.

Follow up:

Can you solve it in O(N) time and O(1) space?

### Solution

02/05/2020:

```
class Solution {
public:
    bool backspaceCompare(string S, string T) {
        string s, t;
        for (auto& c : S) {
            if (c == '#') {
                if (!s.empty()) {
```

```

        s.pop_back();
    }
} else {
    s.push_back(c);
}
}
for (auto& c : T) {
    if (c == '#') {
        if (!t.empty()) {
            t.pop_back();
        }
    } else {
        t.push_back(c);
    }
}
return s == t;
}
};

```

## 868. Binary Gap

### *Description*

Given a positive integer  $N$ , find and return the longest distance between two consecutive 1's in the binary representation of  $N$ .

If there aren't two consecutive 1's, return 0.

Example 1:

Input: 22

Output: 2

Explanation:

22 in binary is 0b10110.

In the binary representation of 22, there are three ones, and two consecutive pairs of 1's.

The first consecutive pair of 1's have distance 2.

The second consecutive pair of 1's have distance 1.

The answer is the largest of these two distances, which is 2.

Example 2:

Input: 5

Output: 2

Explanation:

5 in binary is 0b101.

Example 3:

Input: 6

Output: 1

Explanation:

6 in binary is 0b110.

Example 4:

Input: 8

Output: 0

Explanation:

8 in binary is 0b1000.

There aren't any consecutive pairs of 1's in the binary representation of 8, so we return 0.

Note:

$1 \leq N \leq 10^9$

*Solution*

05/10/2020:

```
class Solution {
public:
    int binaryGap(int N) {
        if (N == 0) return 0;
        int ret = INT_MIN;
        while ((N & 1) == 0) {
            N >>= 1;
        }
        int cnt = 0;
        while (N != 0) {
            N >>= 1;
            if ((N & 1) == 1) {
                ret = max(ret, cnt);
                cnt = 0;
            } else {
                ++cnt;
            }
        }
        return ret == INT_MIN ? 0 : ret + 1;
    }
};
```

## 917. Reverse Only Letters

### Description

Given a string *S*, return the "reversed" string where all characters that are not a letter stay in the same place, and all letters reverse their positions.

Example 1:

Input: "ab-cd"

Output: "dc-ba"

Example 2:

Input: "a-bC-dEf-ghIj"

Output: "j-Ih-gfE-dCba"

Example 3:

Input: "Test1ng-Leet=code-Q!"

Output: "Qedo1ct-eeLg=ntse-T!"

Note:

*S*.length <= 100

33 <= *S*[*i*].ASCIIcode <= 122

*S* doesn't contain \ or "

### Solution

05/11/2020:

```
class Solution {
public:
    string reverseOnlyLetters(string S) {
        int n = S.size(), l = 0, r = n - 1;
        while (l < r) {
            if (isalpha(S[l]) && isalpha(S[r])) {
                swap(S[l++], S[r--]);
            } else if (!isalpha(S[l])) {
                ++l;
            } else { // if (!isalpha(S[r]))
                --r;
            }
        }
        return S;
    }
};
```



```
}  
};
```

```
class Solution {  
public:  
    string reverseOnlyLetters(string S) {  
        int n = S.size(), l = 0, r = n - 1;  
        string ret(S);  
        while (l < n && r >= 0) {  
            for (; l < n && !isalpha(ret[l]); ++l);  
            for (; r >= 0 && !isalpha(S[r]); --r);  
            if (l < n && r >= 0) ret[l++] = S[r--];  
        }  
        return ret;  
    }  
};
```

## 973. K Closest Points to Origin

### Description

We have a list of points on the plane. Find the K closest points to the origin (0, 0).

(Here, the distance between two points on a plane is the Euclidean distance.)

You may return the answer in any order. The answer is guaranteed to be unique (except for the order that it is in.)

Example 1:

Input: points = [[1,3],[-2,2]], K = 1

Output: [[-2,2]]

Explanation:

The distance between (1, 3) and the origin is  $\sqrt{10}$ .

The distance between (-2, 2) and the origin is  $\sqrt{8}$ .

Since  $\sqrt{8} < \sqrt{10}$ , (-2, 2) is closer to the origin.

We only want the closest K = 1 points from the origin, so the answer is just [[-2,2]].

Example 2:

Input: points = [[3,3],[5,-1],[-2,4]], K = 2

Output: [[3,3],[-2,4]]

(The answer `[[-2,4],[3,3]]` would also be accepted.)

Note:

```
1 <= K <= points.length <= 10000
-10000 < points[i][0] < 10000
-10000 < points[i][1] < 10000
```

*Solution*

05/28/2020:

```
class Solution {
public:
    vector<vector<int>> kClosest(vector<vector<int>>& points, int K) {
        priority_queue<pair<double, vector<int>>, vector<pair<double, vector<int>>>,
greater<pair<double, vector<int>>>> q;
        for (auto& p : points) q.push({hypot(p[0], p[1]), p});
        vector<vector<int>> ret;
        while (K-- > 0) {
            pair<double, vector<int>> cur = q.top(); q.pop();
            ret.push_back(cur.second);
        }
        return ret;
    }
};
```

## 986. Interval List Intersections

*Description*

Given two lists of closed intervals, each list of intervals is pairwise disjoint and in sorted order.

Return the intersection of these two interval lists.

(Formally, a closed interval  $[a, b]$  (with  $a \leq b$ ) denotes the set of real numbers  $x$  with  $a \leq x \leq b$ . The intersection of two closed intervals is a set of real numbers that is either empty, or can be represented as a closed interval. For example, the intersection of  $[1, 3]$  and  $[2, 4]$  is  $[2, 3]$ .)

Example 1:

Input: A = [[0,2],[5,10],[13,23],[24,25]], B = [[1,5],[8,12],[15,24],[25,26]]

Output: [[1,2],[5,5],[8,10],[15,23],[24,24],[25,25]]

Reminder: The inputs and the desired output are lists of Interval objects, and not arrays or lists.

Note:

$0 \leq A.length < 1000$

$0 \leq B.length < 1000$

$0 \leq A[i].start, A[i].end, B[i].start, B[i].end < 10^9$

NOTE: input types have been changed on April 15, 2019. Please reset to default code definition to get new method signature.

*Solution*

05/23/2020:

```
class Solution {
public:
    vector<vector<int>> intervalIntersection(vector<vector<int>>& A,
vector<vector<int>>& B) {
        auto ia = A.begin();
        auto ib = B.begin();
        vector<vector<int>> ret;
        for (; ia != A.end() && ib != B.end(); ) {
            if ((*ia)[1] <= (*ib)[1]) {
                if ((*ib)[0] <= (*ia)[1]) {
                    ret.push_back({max((*ia)[0], (*ib)[0]), (*ia)[1]});
                }
                ++ia;
            } else {
                if ((*ib)[1] >= (*ia)[0]) {
                    ret.push_back({max((*ia)[0], (*ib)[0]), (*ib)[1]});
                }
                ++ib;
            }
        }
        return ret;
    }
};
```

## 988. Smallest String Starting From Leaf

### *Description*

Given the root of a binary tree, each node has a value from 0 to 25 representing the letters 'a' to 'z': a value of 0 represents 'a', a value of 1 represents 'b', and so on.

Find the lexicographically smallest string that starts at a leaf of this tree and ends at the root.

(As a reminder, any shorter prefix of a string is lexicographically smaller: for example, "ab" is lexicographically smaller than "aba". A leaf of a node is a node that has no children.)

Example 1:

Input: [0,1,2,3,4,3,4]

Output: "dba"

Example 2:

Input: [25,1,3,1,3,0,2]

Output: "adz"

Example 3:

Input: [2,2,1,null,1,0,null,0]

Output: "abc"

Note:

The number of nodes in the given tree will be between 1 and 8500.  
Each node in the tree will have a value between 0 and 25.

### *Solution*

06/15/2020:

```
/**
```

```

* Definition for a binary tree node.
* struct TreeNode {
*     int val;
*     TreeNode *left;
*     TreeNode *right;
*     TreeNode() : val(0), left(nullptr), right(nullptr) {}
*     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
*     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),
right(right) {}
* };
*/
class Solution {
public:
    string smallestFromLeaf(TreeNode* root) {
        if (!root) return "";
        string ret(8500, 'z'), s;
        backtrack(s, ret, root);
        return ret;
    }

    void backtrack(string& s, string& ret, TreeNode* root) {
        if (!root) return;
        s.push_back(root->val + 'a');
        if (!root->left && !root->right) ret = min(string(s.rbegin(), s.rend()),
ret);
        backtrack(s, ret, root->left);
        backtrack(s, ret, root->right);
        s.pop_back();
    }
};

```

## 997. Find the Town Judge

### Description

In a town, there are  $N$  people labelled from 1 to  $N$ . There is a rumor that one of these people is secretly the town judge.

If the town judge exists, then:

The town judge trusts nobody.

Everybody (except for the town judge) trusts the town judge.

There is exactly one person that satisfies properties 1 and 2.

You are given `trust`, an array of pairs `trust[i] = [a, b]` representing that the person labelled `a` trusts the person labelled `b`.

If the town judge exists and can be identified, return the label of the town judge. Otherwise, return -1.

Example 1:

Input: N = 2, trust = [[1,2]]

Output: 2

Example 2:

Input: N = 3, trust = [[1,3],[2,3]]

Output: 3

Example 3:

Input: N = 3, trust = [[1,3],[2,3],[3,1]]

Output: -1

Example 4:

Input: N = 3, trust = [[1,2],[2,3]]

Output: -1

Example 5:

Input: N = 4, trust = [[1,3],[1,4],[2,3],[2,4],[4,3]]

Output: 3

Note:

1 <= N <= 1000

trust.length <= 10000

trust[i] are all different

trust[i][0] != trust[i][1]

1 <= trust[i][0], trust[i][1] <= N

*Solution*

05/10/2020:

```
class Solution {
public:
    int findJudge(int N, vector<vector<int>>& trust) {
        if (N == 1 && trust.empty()) return 1;
        if (trust.empty() || trust[0].empty()) return -1;
        unordered_map<int, vector<int>> r;
        unordered_set<int> truster;
        for (auto& t : trust) {
            r[t[1]].push_back(t[0]);
        }
    }
};
```

```

        truster.insert(t[0]);
    }
    int judge = -1;
    int cnt = 0;
    for (auto& s : r) {
        if ((int)s.second.size() == N - 1 && truster.count(s.first) == 0) {
            ++cnt;
            judge = s.first;
        }
    }
    return cnt == 1 ? judge : -1;
}
};

```

## 1029. Two City Scheduling

### *Description*

There are  $2N$  people a company is planning to interview. The cost of flying the  $i$ -th person to city A is  $\text{costs}[i][0]$ , and the cost of flying the  $i$ -th person to city B is  $\text{costs}[i][1]$ .

Return the minimum cost to fly every person to a city such that exactly  $N$  people arrive in each city.

Example 1:

Input:  $[[10,20],[30,200],[400,50],[30,20]]$

Output: 110

Explanation:

The first person goes to city A for a cost of 10.

The second person goes to city A for a cost of 30.

The third person goes to city B for a cost of 50.

The fourth person goes to city B for a cost of 20.

The total minimum cost is  $10 + 30 + 50 + 20 = 110$  to have half the people interviewing in each city.

Note:

$1 \leq \text{costs.length} \leq 100$

It is guaranteed that  $\text{costs.length}$  is even.

$1 \leq \text{costs}[i][0], \text{costs}[i][1] \leq 1000$

Solution

06/03/2020:

Backtracking (TLE):

```
class Solution {
public:
    int ret, numOfCityA;
    int twoCitySchedCost(vector<vector<int>>& costs) {
        int n = costs.size(), totalCost = 0;
        ret = INT_MAX, numOfCityA = 0;
        backtrack(0, n, numOfCityA, totalCost, costs);
        return ret;
    }

    void backtrack(int k, int n, int numOfCityA, int totalCost,
vector<vector<int>>& costs) {
        if (k == n) {
            if (numOfCityA == n / 2) ret = min(ret, totalCost);
            return;
        }
        if (numOfCityA > n / 2) return;
        backtrack(k + 1, n, numOfCityA + 1, totalCost + costs[k][0], costs);
        backtrack(k + 1, n, numOfCityA, totalCost + costs[k][1], costs);
    }
};
```

Dynamic Programming:

```
class Solution {
public:
    int twoCitySchedCost(vector<vector<int>>& costs) {
        int n = costs.size() / 2;
        vector<vector<int>> dp(2 * n + 1, vector<int>(n + 1, INT_MAX));
        dp[0][0] = 0;
        // dp[i][j]: the total cost among costs[0..i - 1]: pick j to city A.
        // dp[i][j] = min(dp[i - 1][j - 1] + costs[i - 1][0], dp[i - 1][j] + costs[i - 1][1])
        for (int i = 1; i <= 2 * n; ++i) {
            for (int j = 0; j <= i && j <= n; ++j) {
                if (i - j > n) continue;
                if (dp[i - 1][j] != INT_MAX) {
                    dp[i][j] = min(dp[i][j], dp[i - 1][j] + costs[i - 1][1]);
                }
                if (j > 0 && dp[i - 1][j - 1] != INT_MAX) {
                    dp[i][j] = min(dp[i][j], dp[i - 1][j - 1] + costs[i - 1][0]);
                }
            }
        }
    }
};
```



```

    }
  }
}
return dp[2 * n][n];
}
};

```

Greedy:

```

class Solution {
public:
    int twoCitySchedCost(vector<vector<int>>& costs) {
        sort(costs.begin(), costs.end(), [](vector<int>& nums1, vector<int>& nums2)
        {
            return nums1[0] - nums1[1] < nums2[0] - nums2[1];
        });
        int n = costs.size(), ret = 0;
        for (int i = 0; i < n; ++i) ret += i < n / 2 ? costs[i][0] : costs[i][1];
        return ret;
    }
};

```

## 1057. Campus Bikes

### Description

On a campus represented as a 2D grid, there are  $N$  workers and  $M$  bikes, with  $N \leq M$ . Each worker and bike is a 2D coordinate on this grid.

Our goal is to assign a bike to each worker. Among the available bikes and workers, we choose the (worker, bike) pair with the shortest Manhattan distance between each other, and assign the bike to that worker. (If there are multiple (worker, bike) pairs with the same shortest Manhattan distance, we choose the pair with the smallest worker index; if there are multiple ways to do that, we choose the pair with the smallest bike index). We repeat this process until there are no available workers.

The Manhattan distance between two points  $p1$  and  $p2$  is  $\text{Manhattan}(p1, p2) = |p1.x - p2.x| + |p1.y - p2.y|$ .

Return a vector `ans` of length  $N$ , where `ans[i]` is the index (0-indexed) of the bike that the  $i$ -th worker is assigned to.

Example 1:

Input: workers = [[0,0],[2,1]], bikes = [[1,2],[3,3]]

Output: [1,0]

Explanation:

Worker 1 grabs Bike 0 as they are closest (without ties), and Worker 0 is assigned Bike 1. So the output is [1, 0].

Example 2:

Input: workers = [[0,0],[1,1],[2,0]], bikes = [[1,0],[2,2],[2,1]]

Output: [0,2,1]

Explanation:

Worker 0 grabs Bike 0 at first. Worker 1 and Worker 2 share the same distance to Bike 2, thus Worker 1 is assigned to Bike 2, and Worker 2 will take Bike 1. So the output is [0,2,1].

Note:

$0 \leq \text{workers}[i][j], \text{bikes}[i][j] < 1000$

All worker and bike locations are distinct.

$1 \leq \text{workers.length} \leq \text{bikes.length} \leq 1000$

*Solution*

06/11/2020: Using bucket-sort:

```
class Solution {
public:
    vector<int> assignBikes(vector<vector<int>>& workers, vector<vector<int>>&
bikes) {
        vector<vector<pair<int, int>>> bucket(2001);
        int m = workers.size(), n = bikes.size();
        for (int i = 0; i < m; ++i)
            for (int j = 0; j < n; ++j)
                bucket[abs(workers[i][0] - bikes[j][0]) + abs(workers[i][1] - bikes[j]
[1])].emplace_back(i, j);
        vector<bool> chosen_workers(m, false), chosen_bikes(n, false);
        vector<int> ret(m, -1);
        for (int i = 0; i < 2001; ++i) {
            for (int j = 0; j < (int)bucket[i].size(); ++j) {
                int w = bucket[i][j].first, b = bucket[i][j].second;
                if (!chosen_workers[w] && !chosen_bikes[b]) {
                    ret[w] = b;
                }
            }
        }
    }
};
```

```

        chosen_workers[w] = chosen_bikes[b] = true;
    }
}
return ret;
}
};

```

Using sort:

```

class Solution {
public:
    vector<int> assignBikes(vector<vector<int>>& workers, vector<vector<int>>&
bikes) {
        int m = workers.size(), n = bikes.size();
        vector<pair<int, pair<int, int>>> dist;
        for (int i = 0; i < m; ++i)
            for (int j = 0; j < n; ++j)
                dist.push_back({abs(workers[i][0] - bikes[j][0]) + abs(workers[i][1] -
bikes[j][1]), {i, j}});
        sort(dist.begin(), dist.end());
        vector<bool> chosen_workers(m, false), chosen_bikes(n, false);
        vector<int> ret(m, -1);
        for (auto& d : dist) {
            int i = d.second.first, j = d.second.second;
            if (!chosen_workers[i] && !chosen_bikes[j]) {
                ret[i] = j;
                chosen_workers[i] = chosen_bikes[j] = true;
            }
        }
        return ret;
    }
};

```

Using priority\_queue (much slower than sort):

```

class Solution {
public:
    vector<int> assignBikes(vector<vector<int>>& workers, vector<vector<int>>&
bikes) {
        int m = workers.size(), n = bikes.size();
        priority_queue<pair<int, pair<int, int>>, vector<pair<int, pair<int, int>>>,
greater<pair<int, pair<int, int>>>> pq;
        for (int i = 0; i < m; ++i)
            for (int j = 0; j < n; ++j)
                pq.push({abs(workers[i][0] - bikes[j][0]) + abs(workers[i][1] - bikes[j]
[1]), {i, j}});

```

```

vector<bool> chosen_workers(m, false), chosen_bikes(n, false);
vector<int> ret(m, -1);
while (!pq.empty()) {
    pair<int, pair<int, int>> p = pq.top(); pq.pop();
    int i = p.second.first, j = p.second.second;
    if (!chosen_workers[i] && !chosen_bikes[j]) {
        ret[i] = j;
        chosen_workers[i] = chosen_bikes[j] = true;
    }
}
return ret;
};

```

## 1065. Index Pairs of a String

### Description

Given a text string and words (a list of strings), return all index pairs  $[i, j]$  so that the substring  $\text{text}[i] \dots \text{text}[j]$  is in the list of words.

Example 1:

Input:  $\text{text} = \text{"thestoryoffleetcodeandme"}$ ,  $\text{words} = [\text{"story"}, \text{"fleet"}, \text{"leetcode"}]$   
Output:  $[[3,7], [9,13], [10,17]]$

Example 2:

Input:  $\text{text} = \text{"ababa"}$ ,  $\text{words} = [\text{"aba"}, \text{"ab"}]$   
Output:  $[[0,1], [0,2], [2,3], [2,4]]$

Explanation:

Notice that matches can overlap, see "aba" is found in  $[0,2]$  and  $[2,4]$ .

Note:

All strings contains only lowercase English letters.

It's guaranteed that all strings in words are different.

$1 \leq \text{text.length} \leq 100$

$1 \leq \text{words.length} \leq 20$

$1 \leq \text{words}[i].\text{length} \leq 50$

Return the pairs  $[i,j]$  in sorted order (i.e. sort them by their first coordinate in case of ties sort them by their second coordinate).

Solution

05/10/2020:

```
class Solution {
public:
    vector<vector<int>> indexPairs(string text, vector<string>& words) {
        vector<vector<int>> ret;
        int n = text.size();
        for (auto& w : words) {
            int m = w.size();
            if (m > n) continue;
            for (int i = 0; i < n - m + 1; ++i) {
                if (text.substr(i, m) == w) {
                    ret.push_back({i, i + m - 1});
                }
            }
        }
        sort(ret.begin(), ret.end());
        return ret;
    }
};
```

## 1066. Campus Bikes II

*Description*

On a campus represented as a 2D grid, there are  $N$  workers and  $M$  bikes, with  $N \leq M$ . Each worker and bike is a 2D coordinate on this grid.

We assign one unique bike to each worker so that the sum of the Manhattan distances between each worker and their assigned bike is minimized.

The Manhattan distance between two points  $p1$  and  $p2$  is  $\text{Manhattan}(p1, p2) = |p1.x - p2.x| + |p1.y - p2.y|$ .

Return the minimum possible sum of Manhattan distances between each worker and their assigned bike.

Example 1:

Input: workers = [[0,0],[2,1]], bikes = [[1,2],[3,3]]

Output: 6

Explanation:

We assign bike 0 to worker 0, bike 1 to worker 1. The Manhattan distance of both assignments is 3, so the output is 6.

Example 2:

Input: workers = [[0,0],[1,1],[2,0]], bikes = [[1,0],[2,2],[2,1]]

Output: 4

Explanation:

We first assign bike 0 to worker 0, then assign bike 1 to worker 1 or worker 2, bike 2 to worker 2 or worker 1. Both assignments lead to sum of the Manhattan distances as 4.

Note:

$0 \leq \text{workers}[i][0], \text{workers}[i][1], \text{bikes}[i][0], \text{bikes}[i][1] < 1000$

All worker and bike locations are distinct.

$1 \leq \text{workers.length} \leq \text{bikes.length} \leq 10$

*Solution*

06/11/2020:

Using Dynamic Programming:

```
class Solution {
public:
    int assignBikes(vector<vector<int>>& workers, vector<vector<int>>& bikes) {
        function dist = [](vector<int>& worker, vector<int>& bike) {
            return abs(worker[0] - bike[0]) + abs(worker[1] - bike[1]);
        };
        int m = workers.size(), n = bikes.size();
        vector<vector<int>> dp(m + 1, vector<int>(1 << n, INT_MAX - 4000));
        dp[0][0] = 0;
        int minDist = INT_MAX;
        for (int i = 1; i <= m; ++i) {
            for (int s = 1; s < (1 << n); ++s) {
                for (int j = 0; j < n; ++j) {
                    if ((s & (1 << j)) != 0) {
                        int prev = s ^ (1 << j);
                        dp[i][s] = min(dp[i][s], dp[i - 1][prev] + dist(workers[i - 1],
bikes[j]));
                    }
                    minDist = i == m ? min(minDist, dp[i][s]) : minDist;
                }
            }
        }
    }
};
```

```

    }
}
return minDist;
}
};

```

Using backtracking (TLE):

```

class Solution {
public:
    int minDist = INT_MAX;
    vector<bool> chosen_bikes;
    int assignBikes(vector<vector<int>>& workers, vector<vector<int>>& bikes) {
        int m = workers.size(), n = bikes.size(), dist = 0;
        chosen_bikes.assign(1000, false);
        backtrack(0, dist, workers, bikes);
        return minDist;
    }
    void backtrack(int i, int dist, vector<vector<int>>& workers,
vector<vector<int>>& bikes) {
        if (i == workers.size()) {
            minDist = min(dist, minDist);
            return;
        }
        if (dist > minDist) return; // pruning
        for (int j = 0; j < (int)bikes.size(); ++j) {
            if (chosen_bikes[j] == false) {
                chosen_bikes[j] = true;
                backtrack(i + 1, dist + abs(workers[i][0] - bikes[j][0]) +
abs(workers[i][1] - bikes[j][1]), workers, bikes);
                chosen_bikes[j] = false;
            }
        }
    }
};

```

## 1143. Longest Common Subsequence

### *Description*

Given two strings text1 and text2, return the length of their longest common subsequence.

A subsequence of a string is a new string generated from the original string with some characters (can be none) deleted without changing the relative order of the remaining characters. (eg, "ace" is a subsequence of "abcde" while "aec" is not). A common subsequence of two strings is a subsequence that is common to both strings.

If there is no common subsequence, return 0.

Example 1:

Input: text1 = "abcde", text2 = "ace"

Output: 3

Explanation: The longest common subsequence is "ace" and its length is 3.

Example 2:

Input: text1 = "abc", text2 = "abc"

Output: 3

Explanation: The longest common subsequence is "abc" and its length is 3.

Example 3:

Input: text1 = "abc", text2 = "def"

Output: 0

Explanation: There is no such common subsequence, so the result is 0.

Constraints:

$1 \leq \text{text1.length} \leq 1000$

$1 \leq \text{text2.length} \leq 1000$

The input strings consist of lowercase English characters only.

*Solution*

[Discussion:](#)



```

class Solution {
public:
    int longestCommonSubsequence(string text1, string text2) {
        int m = text1.size(), n = text2.size();
        vector<vector<int>> dp(m + 1, vector<int>(n + 1, 0));
        for (int i = 0; i < m; ++i)
            for (int j = 0; j < n; ++j)
                dp[i + 1][j + 1] = text1[i] == text2[j] ? dp[i][j] + 1 : max(dp[i][j + 1], dp[i + 1][j]);
        return dp.back().back();
    }
};

```

```

class Solution {
public:
    int longestCommonSubsequence(string text1, string text2) {
        int m = text1.size(), n = text2.size();
        vector<int> dp(n + 1, 0);
        for (int i = 0; i < m; ++i) {
            vector<int> tmp(dp);
            for (int j = 0; j < n; ++j) {
                dp[j + 1] = text1[i] == text2[j] ? tmp[j] + 1 : max(dp[j + 1], dp[j]);
            }
        }
        return dp.back();
    }
};

```

## 1170. Compare Strings by Frequency of the Smallest Character

### *Description*

Let's define a function  $f(s)$  over a non-empty string  $s$ , which calculates the frequency of the smallest character in  $s$ . For example, if  $s = "dcce"$  then  $f(s) = 2$  because the smallest character is "c" and its frequency is 2.

Now, given string arrays `queries` and `words`, return an integer array `answer`, where each `answer[i]` is the number of words such that  $f(queries[i]) < f(W)$ , where  $W$  is a word in `words`.

Example 1:

Input: queries = ["cbd"], words = ["zaaaz"]

Output: [1]

Explanation: On the first query we have  $f(\text{"cbd"}) = 1$ ,  $f(\text{"zaaaz"}) = 3$  so  $f(\text{"cbd"}) < f(\text{"zaaaz"})$ .

Example 2:

Input: queries = ["bbb","cc"], words = ["a","aa","aaa","aaaa"]

Output: [1,2]

Explanation: On the first query only  $f(\text{"bbb"}) < f(\text{"aaaa"})$ . On the second query both  $f(\text{"aaa"})$  and  $f(\text{"aaaa"})$  are both  $> f(\text{"cc"})$ .

Constraints:

$1 \leq \text{queries.length} \leq 2000$

$1 \leq \text{words.length} \leq 2000$

$1 \leq \text{queries}[i].\text{length}, \text{words}[i].\text{length} \leq 10$

$\text{queries}[i][j], \text{words}[i][j]$  are English lowercase letters.

*Solution*

02/05/2020:

```
class Solution {
public:
    vector<int> numSmallerByFrequency(vector<string>& queries, vector<string>&
words) {
        int n = queries.size(), N = 11;
        vector<int> ret(n, 0), cnt(N, 0);
        for (auto& w : words) ++cnt[f(w) - 1];
        for (int i = N - 2; i > -1; --i) cnt[i] += cnt[i + 1];
        for (int i = 0; i < n; ++i) ret[i] = cnt[f(queries[i])];
        return ret;
    }

    int f(string& s) {
        char ch = *min_element(s.begin(), s.end());
        return count(s.begin(), s.end(), ch);
    }
};
```

1277. Count Square Submatrices with All Ones

## Description

Given a  $m \times n$  matrix of ones and zeros, return how many square submatrices have all ones.

Example 1:

Input: matrix =

```
[
  [0,1,1,1],
  [1,1,1,1],
  [0,1,1,1]
]
```

Output: 15

Explanation:

There are 10 squares of side 1.

There are 4 squares of side 2.

There is 1 square of side 3.

Total number of squares =  $10 + 4 + 1 = 15$ .

Example 2:

Input: matrix =

```
[
  [1,0,1],
  [1,1,0],
  [1,1,0]
]
```

Output: 7

Explanation:

There are 6 squares of side 1.

There is 1 square of side 2.

Total number of squares =  $6 + 1 = 7$ .

Constraints:

$1 \leq \text{arr.length} \leq 300$

$1 \leq \text{arr}[0].\text{length} \leq 300$

$0 \leq \text{arr}[i][j] \leq 1$

## Solution

01/14/2020 (Dynamic Programming):

```
class Solution {
public:
```

```

int countSquares(vector<vector<int>>& matrix) {
    int m = matrix.size(), n = matrix[0].size(), ret = 0;;
    vector<vector<int>> dp(m, vector<int>(n, 0));
    for (int i = 0; i < m; ++i) {
        for (int j = 0; j < n; ++j) {
            if (i == 0 || j == 0)
                dp[i][j] = matrix[i][j];
            else
                dp[i][j] = matrix[i][j] == 0 ? 0 : min(min(dp[i - 1][j - 1], dp[i - 1][j]), dp[i][j - 1]) + 1;
            ret += dp[i][j];
        }
    }
    return ret;
}
};

```

```

class Solution {
public:
    int countSquares(vector<vector<int>>& matrix) {
        if (matrix.empty() || matrix[0].empty()) return 0;
        int m = matrix.size(), n = matrix[0].size();
        int ret = 0;
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                if (i >= 1 && j >= 1 && matrix[i][j] == 1)
                    matrix[i][j] = min(min(matrix[i - 1][j], matrix[i][j - 1]), matrix[i - 1][j - 1]) + 1;
                ret += matrix[i][j];
            }
        }
        return ret;
    }
};

```

## 1287. Element Appearing More Than 25% In Sorted Array

### Description

Given an integer array sorted in non-decreasing order, there is exactly one integer in the array that occurs more than 25% of the time.

Return that integer.

Example 1:

Input: arr = [1,2,2,6,6,6,6,7,10]

Output: 6

Constraints:

$1 \leq \text{arr.length} \leq 10^4$

$0 \leq \text{arr}[i] \leq 10^5$

*Solution*

05/10/2020:

```
class Solution {
public:
    int findSpecialInteger(vector<int>& arr) {
        int n = arr.size();
        int m = (n + 1) / 4;
        for (int i = 0; i < n - m; ++i) {
            if (arr[i] == arr[i + m]) {
                return arr[i];
            }
        }
        return -1;
    }
};
```

## 1314. Matrix Block Sum

*Description*

Given a  $m \times n$  matrix mat and an integer K, return a matrix answer where each answer[i][j] is the sum of all elements mat[r][c] for  $i - K \leq r \leq i + K$ ,  $j - K \leq c \leq j + K$ , and (r, c) is a valid position in the matrix.

Example 1:

Input: mat = [[1,2,3],[4,5,6],[7,8,9]], K = 1

Output: [[12,21,16],[27,45,33],[24,39,28]]

Example 2:

Input: mat = [[1,2,3],[4,5,6],[7,8,9]], K = 2  
Output: [[45,45,45],[45,45,45],[45,45,45]]

Constraints:

```
m == mat.length
n == mat[i].length
1 <= m, n, K <= 100
1 <= mat[i][j] <= 100
```

*Solution*

01/14/2020 (Definition):

```
class Solution {
public:
    vector<vector<int>> matrixBlockSum(vector<vector<int>>& mat, int K) {
        int m = mat.size(), n = mat[0].size();
        vector<vector<int>> ret(m, vector<int>(n, 0));
        for (int i = 0; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                for (int k = -K; k <= K; ++k) {
                    for (int l = -K; l <= K; ++l) {
                        if (i + k >= 0 && i + k < m && j + l >= 0 && j + l < n) {
                            ret[i][j] += mat[i + k][j + l];
                        }
                    }
                }
            }
        }
        return ret;
    }
};
```

01/14/2020 (Dynamic Programming):

```
class Solution {
public:
    vector<vector<int>> matrixBlockSum(vector<vector<int>>& mat, int K) {
        int m = mat.size(), n = mat[0].size();
        vector<vector<int>> dp(m, vector<int>(n, 0));
        vector<vector<int>> ret(m, vector<int>(n, 0));
        dp[0][0] = mat[0][0];
        for (int i = 1; i < m; ++i) dp[i][0] += dp[i - 1][0] + mat[i][0];
        for (int j = 1; j < n; ++j) dp[0][j] += dp[0][j - 1] + mat[0][j];
        for (int i = 1; i < m; ++i)
```

```

    for (int j = 1; j < n; ++j)
        dp[i][j] = mat[i][j] + dp[i][j - 1] + dp[i - 1][j] - dp[i - 1][j - 1];
    for (int i = 0; i < m; ++i) {
        for (int j = 0; j < n; ++j) {
            int r1 = max(i - K, 0);
            int c1 = max(j - K, 0);
            int r2 = min(i + K, m - 1);
            int c2 = min(j + K, n - 1);
            ret[i][j] = dp[r2][c2];
            if (r1 > 0) ret[i][j] -= dp[r1 - 1][c2];
            if (c1 > 0) ret[i][j] -= dp[r2][c1 - 1];
            if (r1 > 0 && c1 > 0) ret[i][j] += dp[r1 - 1][c1 - 1];
        }
    }
    return ret;
}
};

```

## 1331. Rank Transform of an Array

### Description

Given an array of integers `arr`, replace each element with its rank.

The rank represents how large the element is. The rank has the following rules:

Rank is an integer starting from 1.

The larger the element, the larger the rank. If two elements are equal, their rank must be the same.

Rank should be as small as possible.

Example 1:

Input: `arr = [40,10,20,30]`

Output: `[4,1,2,3]`

Explanation: 40 is the largest element. 10 is the smallest. 20 is the second smallest. 30 is the third smallest.

Example 2:

Input: `arr = [100,100,100]`

Output: `[1,1,1]`

Explanation: Same elements share the same rank.

Example 3:

Input: `arr = [37,12,28,9,100,56,80,5,12]`

Output: [5,3,4,2,8,6,7,1,3]

Constraints:

$0 \leq \text{arr.length} \leq 105$

$-109 \leq \text{arr}[i] \leq 109$

*Solution*

05/10/2020:

```
class Solution {
public:
    vector<int> arrayRankTransform(vector<int>& arr) {
        if (arr.empty()) return {};
        set<int> keys(arr.begin(), arr.end());
        unordered_map<int, int> mp;
        int i = 1;
        for (auto& k : keys) {
            mp[k] = i++;
        }
        vector<int> ret;
        for (auto& a : arr) {
            ret.push_back(mp[a]);
        }
        return ret;
    }
};
```

## 1422. Maximum Score After Splitting a String

*Description*

Given a string *s* of zeros and ones, return the maximum score after splitting the string into two non-empty substrings (i.e. left substring and right substring).

The score after splitting a string is the number of zeros in the left substring plus the number of ones in the right substring.

Example 1:

Input: *s* = "011101"



Output: 5

Explanation:

All possible ways of splitting  $s$  into two non-empty substrings are:

left = "0" and right = "11101", score = 1 + 4 = 5

left = "01" and right = "1101", score = 1 + 3 = 4

left = "011" and right = "101", score = 1 + 2 = 3

left = "0111" and right = "01", score = 1 + 1 = 2

left = "01110" and right = "1", score = 2 + 1 = 3

Example 2:

Input:  $s = "00111"$

Output: 5

Explanation: When left = "00" and right = "111", we get the maximum score = 2 + 3 = 5

Example 3:

Input:  $s = "1111"$

Output: 3

Constraints:

$2 \leq s.length \leq 500$

The string  $s$  consists of characters '0' and '1' only.

*Solution*

04/25/2020:

```
class Solution {
public:
    int maxScore(string s) {
        int zeros = 0, ones = 0;
        for (auto& c : s) {
            if (c == '0') {
                ++zeros;
            } else {
                ++ones;
            }
        }
        int n = s.size();
        int cur_zeros = 0;
        int ret = 0;
        for (int i = 0; i < n - 1; ++i) {
            if (s[i] == '0') {
                ++cur_zeros;
            }
            ret = max(ret, cur_zeros + (n - (i + 1) - (zeros - cur_zeros)));
        }
    }
};
```

```
    }  
    return ret;  
  }  
};
```

## 1424. Diagonal Traverse II

### Description

Given a list of lists of integers, `nums`, return all elements of `nums` in diagonal order as shown in the below images.

Example 1:

Input: `nums = [[1,2,3],[4,5,6],[7,8,9]]`

Output: `[1,4,2,7,5,3,8,6,9]`

Example 2:

Input: `nums = [[1,2,3,4,5],[6,7],[8],[9,10,11],[12,13,14,15,16]]`

Output: `[1,6,2,8,7,3,9,4,12,10,5,13,11,14,15,16]`

Example 3:

Input: `nums = [[1,2,3],[4],[5,6,7],[8],[9,10,11]]`

Output: `[1,4,2,5,3,8,6,9,7,10,11]`

Example 4:

Input: `nums = [[1,2,3,4,5,6]]`

Output: `[1,2,3,4,5,6]`

Constraints:

`1 <= nums.length <= 105`

`1 <= nums[i].length <= 105`

`1 <= nums[i][j] <= 109`

There at most `105` elements in `nums`.

### Solution

04/25/2020:

---

```

class Solution {
public:
    vector<int> findDiagonalOrder(vector<vector<int>>& nums) {
        int m = nums.size();
        if (m <= 0) return {};
        map<int, vector<int>> mp;
        for (int i = 0; i < m; ++i) {
            int n = nums[i].size();
            for (int j = 0; j < n; ++j) {
                mp[i + j].push_back(nums[i][j]);
            }
        }
        vector<int> ret;
        for (auto& m : mp) {
            ret.insert(ret.end(), m.second.rbegin(), m.second.rend());
        }
        return ret;
    }
};

```

## 1438. Longest Continuous Subarray With Absolute Diff Less Than or Equal to Limit

### Description

Given an array of integers `nums` and an integer `limit`, return the size of the longest continuous subarray such that the absolute difference between any two elements is less than or equal to `limit`.

In case there is no subarray satisfying the given condition return `0`.

Example 1:

Input: `nums = [8,2,4,7]`, `limit = 4`

Output: `2`

Explanation: All subarrays are:

`[8]` with maximum absolute diff  $|8-8| = 0 \leq 4$ .

`[8,2]` with maximum absolute diff  $|8-2| = 6 > 4$ .

`[8,2,4]` with maximum absolute diff  $|8-2| = 6 > 4$ .

`[8,2,4,7]` with maximum absolute diff  $|8-2| = 6 > 4$ .

`[2]` with maximum absolute diff  $|2-2| = 0 \leq 4$ .

`[2,4]` with maximum absolute diff  $|2-4| = 2 \leq 4$ .

[2,4,7] with maximum absolute diff  $|2-7| = 5 > 4$ .  
[4] with maximum absolute diff  $|4-4| = 0 \leq 4$ .  
[4,7] with maximum absolute diff  $|4-7| = 3 \leq 4$ .  
[7] with maximum absolute diff  $|7-7| = 0 \leq 4$ .  
Therefore, the size of the longest subarray is 2.

Example 2:

Input: nums = [10,1,2,4,7,2], limit = 5

Output: 4

Explanation: The subarray [2,4,7,2] is the longest since the maximum absolute diff is  $|2-7| = 5 \leq 5$ .

Example 3:

Input: nums = [4,2,2,2,4,4,2,2], limit = 0

Output: 3

Constraints:

$1 \leq \text{nums.length} \leq 10^5$

$1 \leq \text{nums}[i] \leq 10^9$

$0 \leq \text{limit} \leq 10^9$

*Solution*

05/03/2020:

```
class Solution {
public:
    int longestSubarray(vector<int>& nums, int limit, int start = 0, int stop =
INT_MAX) {
        if (stop == INT_MAX) stop = nums.size() - 1;
        if (start > stop) return 0;
        auto min_it = min_element(nums.begin() + start, nums.begin() + stop + 1);
        auto max_it = max_element(nums.begin() + start, nums.begin() + stop + 1);
        if (*max_it - *min_it <= limit) return stop - start + 1;
        auto lo = min_it, hi = min_it;
        for (; lo - nums.begin() > start && *(lo - 1) - *min_it <= limit; --lo);
        for (; hi - nums.begin() < stop && *(hi + 1) - *min_it <= limit; ++hi);
        int left = int(min_it - nums.begin() - 1) - start + 1 >= hi - lo + 1 ?
longestSubarray(nums, limit, start, int(min_it - nums.begin() - 1)) : 0;
        int right = stop - int(min_it - nums.begin() - 1) + 1 >= hi - lo + 1 ?
longestSubarray(nums, limit, int(min_it - nums.begin() + 1), stop) : 0;
        return max(int(hi - lo + 1), max(left, right));
    }
};
```

# 1441. Build an Array With Stack Operations

## Description

Given an array `target` and an integer `n`. In each iteration, you will read a number from `list = {1,2,3..., n}`.

Build the target array using the following operations:

Push: Read a new element from the beginning list, and push it in the array.

Pop: delete the last element of the array.

If the target array is already built, stop reading more elements.

You are guaranteed that the target array is strictly increasing, only containing numbers between 1 to `n` inclusive.

Return the operations to build the target array.

You are guaranteed that the answer is unique.

Example 1:

Input: `target = [1,3]`, `n = 3`

Output: `["Push","Push","Pop","Push"]`

Explanation:

Read number 1 and automatically push in the array -> `[1]`

Read number 2 and automatically push in the array then Pop it -> `[1]`

Read number 3 and automatically push in the array -> `[1,3]`

Example 2:

Input: `target = [1,2,3]`, `n = 3`

Output: `["Push","Push","Push"]`

Example 3:

Input: `target = [1,2]`, `n = 4`

Output: `["Push","Push"]`

Explanation: You only need to read the first 2 numbers and stop.

Example 4:

Input: `target = [2,3,4]`, `n = 4`

Output: `["Push","Pop","Push","Push","Push"]`

Constraints:

`1 <= target.length <= 100`

`1 <= target[i] <= 100`

```
1 <= n <= 100
target is strictly increasing.
```

*Solution*

05/10/2020:

```
class Solution {
public:
    vector<string> buildArray(vector<int>& target, int n) {
        int m = *max_element(target.begin(), target.end());
        unordered_set<int> s(target.begin(), target.end());
        vector<string> ret;
        for (int i = 1; i <= m; ++i) {
            ret.push_back("Push");
            if (s.count(i) == 0) {
                ret.push_back("Pop");
            }
        }
        return ret;
    }
};
```

## 1447. Simplified Fractions

*Description*

Given an integer  $n$ , return a list of all simplified fractions between 0 and 1 (exclusive) such that the denominator is less-than-or-equal-to  $n$ . The fractions can be in any order.

Example 1:

Input:  $n = 2$

Output: ["1/2"]

Explanation: "1/2" is the only unique fraction with a denominator less-than-or-equal-to 2.

Example 2:

Input:  $n = 3$

Output: ["1/2", "1/3", "2/3"]

Example 3:

Input: n = 4

Output: ["1/2","1/3","1/4","2/3","3/4"]

Explanation: "2/4" is not a simplified fraction because it can be simplified to "1/2".

Example 4:

Input: n = 1

Output: []

Constraints:

$1 \leq n \leq 100$

*Solution*

05/16/2020:

```
class Solution {
public:
    vector<string> simplifiedFractions(int n) {
        unordered_set<string> ret;
        for (int i = 1; i < n; ++i) {
            for (int j = i + 1; j <= n; ++j) {
                int gcd = __gcd(i, j);
                int numerator = i / gcd;
                int denominator = j / gcd;
                ret.insert(to_string(numerator) + "/" + to_string(denominator));
            }
        }
        return vector<string>(ret.begin(), ret.end());
    }
};
```

## 1452. People Whose List of Favorite Companies Is Not a Subset of Another List

*Description*

Given the array favoriteCompanies where favoriteCompanies[i] is the list of favorites companies for the ith person (indexed from 0).

Return the indices of people whose list of favorite companies is not a subset of any other list of favorites companies. You must return the indices in increasing order.

Example 1:

Input: favoriteCompanies = [ ["leetcode", "google", "facebook"],  
 ["google", "microsoft"], ["google", "facebook"], ["google"], ["amazon"] ]

Output: [0,1,4]

Explanation:

Person with index=2 has favoriteCompanies[2]=["google","facebook"] which is a subset of favoriteCompanies[0]=["leetcode","google","facebook"] corresponding to the person with index 0.

Person with index=3 has favoriteCompanies[3]=["google"] which is a subset of favoriteCompanies[0]=["leetcode","google","facebook"] and favoriteCompanies[1]= ["google","microsoft"].

Other lists of favorite companies are not a subset of another list, therefore, the answer is [0,1,4].

Example 2:

Input: favoriteCompanies = [ ["leetcode", "google", "facebook"],  
 ["leetcode", "amazon"], ["facebook", "google"] ]

Output: [0,1]

Explanation: In this case favoriteCompanies[2]=["facebook","google"] is a subset of favoriteCompanies[0]=["leetcode","google","facebook"], therefore, the answer is [0,1].

Example 3:

Input: favoriteCompanies = [ ["leetcode"], ["google"], ["facebook"], ["amazon"] ]

Output: [0,1,2,3]

Constraints:

1 <= favoriteCompanies.length <= 100

1 <= favoriteCompanies[i].length <= 500

1 <= favoriteCompanies[i][j].length <= 20

All strings in favoriteCompanies[i] are distinct.

All lists of favorite companies are distinct, that is, If we sort alphabetically each list then favoriteCompanies[i] != favoriteCompanies[j].

All strings consist of lowercase English letters only.

*Solution*

05/16/2020:



```

class Solution {
public:
    vector<int> peopleIndexes(vector<vector<string>>& favoriteCompanies) {
        vector<pair<unordered_set<string>, int>> companies;
        int n = favoriteCompanies.size();
        for (int i = 0; i < n; ++i) {
            companies.emplace_back(unordered_set<string>(favoriteCompanies[i].begin(),
favoriteCompanies[i].end()), i);
        }
        sort(companies.begin(), companies.end(), [](pair<unordered_set<string>,
int>& p1, pair<unordered_set<string>, int>& p2) {
            return p1.first.size() > p2.first.size();
        });
        vector<int> ret;
        ret.push_back(companies[0].second);
        for (int i = 1; i < n; ++i) {
            int cnt = 0;
            for (int j = 0; j < i; ++j) {
                bool isSubset = true;
                for (auto& s : companies[i].first) {
                    if (companies[j].first.count(s) == 0) {
                        isSubset = false;
                        break;
                    }
                }
                if (!isSubset) {
                    ++cnt;
                }
            }
            if (cnt == i) ret.push_back(companies[i].second);
        }
        sort(ret.begin(), ret.end());
        return ret;
    }
};

```

## 1461. Check If a String Contains All Binary Codes of Size K

### *Description*

Given a binary string  $s$  and an integer  $k$ .

Return True if all binary codes of length  $k$  is a substring of  $s$ . Otherwise, return False.

Example 1:

Input: s = "00110110", k = 2

Output: true

Explanation: The binary codes of length 2 are "00", "01", "10" and "11". They can be all found as substrings at indices 0, 1, 3 and 2 respectively.

Example 2:

Input: s = "00110", k = 2

Output: true

Example 3:

Input: s = "0110", k = 1

Output: true

Explanation: The binary codes of length 1 are "0" and "1", it is clear that both exist as a substring.

Example 4:

Input: s = "0110", k = 2

Output: false

Explanation: The binary code "00" is of length 2 and doesn't exist in the array.

Example 5:

Input: s = "0000000001011100", k = 4

Output: false

Constraints:

$1 \leq s.length \leq 5 * 10^5$

s consists of 0's and 1's only.

$1 \leq k \leq 20$

*Solution*

05/30/2020:

```
class Solution {
public:
    bool hasAllCodes(string s, int k) {
        bool ret = true;
        int n = 1 << k;
        int sz = s.size() - k;
        unordered_set<string> required;
        for (int i = 0; i <= sz; ++i) {
            string sub = s.substr(i, k);
```

```
        required.insert(sub);
    }
    return required.size() == n;
}
};
```

## 1463. Cherry Pickup II

### *Description*

Given a rows x cols matrix grid representing a field of cherries. Each cell in grid represents the number of cherries that you can collect.

You have two robots that can collect cherries for you, Robot #1 is located at the top-left corner (0,0) , and Robot #2 is located at the top-right corner (0, cols-1) of the grid.

Return the maximum number of cherries collection using both robots by following the rules below:

From a cell (i,j), robots can move to cell (i+1, j-1) , (i+1, j) or (i+1, j+1). When any robot is passing through a cell, It picks it up all cherries, and the cell becomes an empty cell (0).

When both robots stay on the same cell, only one of them takes the cherries.

Both robots cannot move outside of the grid at any moment.

Both robots should reach the bottom row in the grid.

Example 1:

Input: grid = [[3,1,1],[2,5,1],[1,5,5],[2,1,1]]

Output: 24

Explanation: Path of robot #1 and #2 are described in color green and blue respectively.

Cherries taken by Robot #1, (3 + 2 + 5 + 2) = 12.

Cherries taken by Robot #2, (1 + 5 + 5 + 1) = 12.

Total of cherries: 12 + 12 = 24.

Example 2:

Input: grid = [[1,0,0,0,0,0,1],[2,0,0,0,0,3,0],[2,0,9,0,0,0,0],[0,3,0,5,4,0,0],[1,0,2,3,0,0,6]]

Output: 28

Explanation: Path of robot #1 and #2 are described in color green and blue respectively.

Cherries taken by Robot #1,  $(1 + 9 + 5 + 2) = 17$ .

Cherries taken by Robot #2,  $(1 + 3 + 4 + 3) = 11$ .

Total of cherries:  $17 + 11 = 28$ .

Example 3:

Input: grid = `[[1,0,0,3],[0,0,0,3],[0,0,3,3],[9,0,3,3]]`

Output: 22

Example 4:

Input: grid = `[[1,1],[1,1]]`

Output: 4

Constraints:

```
rows == grid.length
cols == grid[i].length
2 <= rows, cols <= 70
0 <= grid[i][j] <= 100
```

*Solution*

06/07/2020:

```
int dp[71][71][71];

class Solution {
public:
    int cherryPickup(vector<vector<int>>& grid) {
        fill_n(dp[0][0], 71 * 71 * 71, -1);
        int m = grid.size(), n = grid[0].size();
        dp[0][0][n - 1] = grid[0][0] + grid[0][n - 1];
        for (int i = 1; i < m; ++i) {
            for (int j = 0; j < n; ++j) {
                for (int k = 0; k < n; ++k) {
                    int cur = -1;
                    for (int u = -1; u <= 1; ++u) {
                        int nj = j + u;
                        for (int v = -1; v <= 1; ++v) {
                            int nk = k + v;
                            if (nj >= 0 && nj < n && nk >= 0 && nk < n && dp[i - 1][nj][nk] !=
-1)
                                cur = max(cur, dp[i - 1][nj][nk] + grid[i][j] + grid[i][k]);
                        }
                    }
                    if (j == k) cur -= grid[i][j];
                }
            }
        }
    }
};
```

```

        dp[i][j][k] = max(dp[i][j][k], cur);
    }
}
}
int ret = -1;
for (int j = 0; j < n; ++j) ret = max(ret, *max_element(dp[m - 1][j], dp[m -
1][j] + n));
return ret;
}
};

```

## 1471. The k Strongest Values in an Array

### Description

Given an array of integers `arr` and an integer `k`.

A value `arr[i]` is said to be stronger than a value `arr[j]` if  $|arr[i] - m| > |arr[j] - m|$  where `m` is the median of the array.

If  $|arr[i] - m| == |arr[j] - m|$ , then `arr[i]` is said to be stronger than `arr[j]` if `arr[i] > arr[j]`.

Return a list of the strongest `k` values in the array. return the answer in any arbitrary order.

Median is the middle value in an ordered integer list. More formally, if the length of the list is `n`, the median is the element in position  $((n - 1) / 2)$  in the sorted list (0-indexed).

For `arr = [6, -3, 7, 2, 11]`, `n = 5` and the median is obtained by sorting the array `arr = [-3, 2, 6, 7, 11]` and the median is `arr[m]` where  $m = ((5 - 1) / 2) = 2$ . The median is 6.

For `arr = [-7, 22, 17, 3]`, `n = 4` and the median is obtained by sorting the array `arr = [-7, 3, 17, 22]` and the median is `arr[m]` where  $m = ((4 - 1) / 2) = 1$ . The median is 3.

Example 1:

Input: `arr = [1,2,3,4,5]`, `k = 2`

Output: `[5,1]`

Explanation: Median is 3, the elements of the array sorted by the strongest are `[5,1,4,2,3]`. The strongest 2 elements are `[5, 1]`. `[1, 5]` is also accepted answer.

Please note that although  $|5 - 3| == |1 - 3|$  but 5 is stronger than 1 because  $5 > 1$ .

Example 2:

Input: arr = [1,1,3,5,5], k = 2

Output: [5,5]

Explanation: Median is 3, the elements of the array sorted by the strongest are [5,5,1,1,3]. The strongest 2 elements are [5, 5].

Example 3:

Input: arr = [6,7,11,7,6,8], k = 5

Output: [11,8,6,6,7]

Explanation: Median is 7, the elements of the array sorted by the strongest are [11,8,6,6,7,7].

Any permutation of [11,8,6,6,7] is accepted.

Example 4:

Input: arr = [6,-3,7,2,11], k = 3

Output: [-3,11,2]

Example 5:

Input: arr = [-7,22,17,3], k = 2

Output: [22,17]

Constraints:

$1 \leq \text{arr.length} \leq 10^5$

$-10^5 \leq \text{arr}[i] \leq 10^5$

$1 \leq k \leq \text{arr.length}$

*Solution*

06/06/2020:

```
int m;
class Solution {
public:
    vector<int> getStrongest(vector<int>& arr, int k) {
        int n = arr.size();
        sort(arr.begin(), arr.end());
        m = arr[(n - 1) / 2];
        sort(arr.begin(), arr.end(), [](int a, int b) {
            if (abs(a - m) == abs(b - m)) return a > b;
            return abs(a - m) > abs(b - m);
        });
        vector<int> ret;
        for (int i = 0; i < k; ++i)
            ret.push_back(arr[i]);
        return ret;
    }
};
```

```
}  
};
```

## Check If a String Is a Valid Sequence from Root to Leaves Path in a Binary Tree

### *Description*

Given a binary tree where each path going from the root to any leaf form a valid sequence, check if a given string is a valid sequence in such binary tree.

We get the given string from the concatenation of an array of integers arr and the concatenation of all values of the nodes along a path results in a sequence in the given binary tree.

Example 1:

Input: root = [0,1,0,0,1,0,null,null,1,0,0], arr = [0,1,0,1]

Output: true

Explanation:

The path 0 -> 1 -> 0 -> 1 is a valid sequence (green color in the figure).

Other valid sequences are:

0 -> 1 -> 1 -> 0

0 -> 0 -> 0

Example 2:

Input: root = [0,1,0,0,1,0,null,null,1,0,0], arr = [0,0,1]

Output: false

Explanation: The path 0 -> 0 -> 1 does not exist, therefore it is not even a sequence.

Example 3:

Input: root = [0,1,0,0,1,0,null,null,1,0,0], arr = [0,1,1]

Output: false

Explanation: The path 0 -> 1 -> 1 is a sequence, but it is not a valid sequence.

Constraints:

$1 \leq \text{arr.length} \leq 5000$

$0 \leq \text{arr}[i] \leq 9$

Each node's value is between  $[0 - 9]$ .

*Solution*

[Discussion](#)

- Using `unordered_map` is straightforward: convert numbers visited to a string and store seen strings in an `unordered_map`. Then just check whether the target string generated by the `arr` is in the `unordered_map` or not.
- Time complexity:  $O(n)$
- Space complexity:  $O(n^2)$

```
class Solution {
private:
    unordered_set<string> s;
    string target;
public:
    bool isValidSequence(TreeNode* root, vector<int>& arr) {
        for (auto& n : arr) target += to_string(n);
        dfs(root);
        return s.count(target);
    }

    void dfs(TreeNode* root, string str="") {
        if (root == nullptr) return;
        if (root->left == nullptr && root->right == nullptr) {
            str += to_string(root->val);
            if (!str.empty()) s.insert(str);
            return;
        }
        dfs(root->left, str + to_string(root->val));
        dfs(root->right, str + to_string(root->val));
    }
};
```

- Another way is more straightforward, we check the elements in the `arr` while traversing the nodes in the tree:
- Time complexity:  $O(n)$
- Space complexity:  $O(n)$



```

class Solution {
public:
    bool isValidSequence(TreeNode* root, vector<int>& arr, int i = 0) {
        if (root == nullptr) return false;
        if (i >= (int)arr.size()) return false;
        if (root->left == nullptr && root->right == nullptr && i == (int)arr.size()
- 1) return arr.back() == root->val;
        return root->val == arr[i] && (isValidSequence(root->left, arr, i + 1) ||
isValidSequence(root->right, arr, i + 1));
    }
};

```

## 1477. Find Two Non-overlapping Sub-arrays Each With Target Sum

### Description

Given an array of integers `arr` and an integer `target`.

You have to find two non-overlapping sub-arrays of `arr` each with sum equal `target`. There can be multiple answers so you have to find an answer where the sum of the lengths of the two sub-arrays is minimum.

Return the minimum sum of the lengths of the two required sub-arrays, or return `-1` if you cannot find such two sub-arrays.

Example 1:

Input: `arr = [3,2,2,4,3]`, `target = 3`

Output: 2

Explanation: Only two sub-arrays have sum = 3 (`[3]` and `[3]`). The sum of their lengths is 2.

Example 2:

Input: `arr = [7,3,4,7]`, `target = 7`

Output: 2

Explanation: Although we have three non-overlapping sub-arrays of sum = 7 (`[7]`, `[3,4]` and `[7]`), but we will choose the first and third sub-arrays as the sum of their lengths is 2.

Example 3:

Input: `arr = [4,3,2,6,2,3,4]`, `target = 6`

Output: -1

Explanation: We have only one sub-array of sum = 6.

Example 4:

Input: arr = [5,5,4,4,5], target = 3

Output: -1

Explanation: We cannot find a sub-array of sum = 3.

Example 5:

Input: arr = [3,1,1,1,5,1,2,1], target = 3

Output: 3

Explanation: Note that sub-arrays [1,2] and [2,1] cannot be an answer because they overlap.

Constraints:

1 <= arr.length <= 10<sup>5</sup>

1 <= arr[i] <= 1000

1 <= target <= 10<sup>8</sup>

*Solution*

06/13/2020:

```
class Solution {
public:
    int minSumOfLengths(vector<int>& arr, int target) {
        queue<pair<int, int>> q;
        int ret = INT_MAX, preMin = INT_MAX;
        int lo = 0, hi = 0, cur = 0;
        while (hi < arr.size()) {
            cur += arr[hi];
            ++hi;
            while (cur > target && lo < hi) {
                cur -= arr[lo];
                ++lo;
            }
            while (!q.empty() && q.front().second <= lo) {
                preMin = min(preMin, q.front().second - q.front().first);
                q.pop();
            }
            if (cur == target) {
                if (preMin != INT_MAX) {
                    ret = min(ret, preMin + hi - lo);
                }
                q.emplace(lo, hi);
            }
        }
    }
};
```

```

    }
    return ret == INT_MAX ? -1 : ret;
}
};

```

```

void solve(vector<int>& v, int t, vector<int>& ret) {
    int n = v.size();
    ret.resize(n + 1);
    fill(ret.begin(), ret.end(), 1e9);
    int cur = 0;
    int lhs = 0;
    for (int i = 1; i <= n; ++i) {
        cur += v[i - 1];
        while (cur > t) cur -= v[lhs++];
        if (cur == t) ret[i] = i - lhs;
    }
}

class Solution {
public:
    int minSumOfLengths(vector<int>& arr, int target) {
        vector<int> lhs, rhs;
        solve(arr, target, lhs);
        reverse(arr.begin(), arr.end());
        solve(arr, target, rhs);
        for (int i = 1; i < (int)lhs.size(); ++i) lhs[i] = min(lhs[i], lhs[i - 1]);
        for (int i = 1; i < (int)rhs.size(); ++i) rhs[i] = min(rhs[i], rhs[i - 1]);
        int ret = 1e9;
        int n = arr.size();
        for (int i = 1; i < n; ++i) ret = min(ret, lhs[i] + rhs[n - i]);
        if (ret >= 1e9) ret = -1;
        return ret;
    }
};

```

## 1480. Running Sum of 1d Array

### Description

Given an array `nums`. We define a running sum of an array as `runningSum[i] = sum(nums[0]...nums[i])`.

Return the running sum of `nums`.

Example 1:

Input: nums = [1,2,3,4]

Output: [1,3,6,10]

Explanation: Running sum is obtained as follows: [1, 1+2, 1+2+3, 1+2+3+4].

Example 2:

Input: nums = [1,1,1,1,1]

Output: [1,2,3,4,5]

Explanation: Running sum is obtained as follows: [1, 1+1, 1+1+1, 1+1+1+1, 1+1+1+1+1].

Example 3:

Input: nums = [3,1,2,10,1]

Output: [3,4,6,16,17]

Constraints:

$1 \leq \text{nums.length} \leq 1000$

$-10^6 \leq \text{nums}[i] \leq 10^6$

*Solution*

06/13/2020:

```
class Solution {
public:
    vector<int> runningSum(vector<int>& nums) {
        partial_sum(nums.begin(), nums.end(), nums.begin());
        return nums;
    }
};
```

## 1482. Minimum Number of Days to Make m Bouquets

*Description*

Given an integer array bloomDay, an integer m and an integer k.

We need to make m bouquets. To make a bouquet, you need to use k adjacent flowers from the garden.

The garden consists of  $n$  flowers, the  $i$ th flower will bloom in the  $\text{bloomDay}[i]$  and then can be used in exactly one bouquet.

Return the minimum number of days you need to wait to be able to make  $m$  bouquets from the garden. If it is impossible to make  $m$  bouquets return  $-1$ .

Example 1:

Input:  $\text{bloomDay} = [1,10,3,10,2]$ ,  $m = 3$ ,  $k = 1$

Output: 3

Explanation: Let's see what happened in the first three days.  $x$  means flower bloomed and  $_$  means flower didn't bloom in the garden.

We need 3 bouquets each should contain 1 flower.

After day 1:  $[x, _, _, _, _]$  // we can only make one bouquet.

After day 2:  $[x, _, _, _, x]$  // we can only make two bouquets.

After day 3:  $[x, _, x, _, x]$  // we can make 3 bouquets. The answer is 3.

Example 2:

Input:  $\text{bloomDay} = [1,10,3,10,2]$ ,  $m = 3$ ,  $k = 2$

Output:  $-1$

Explanation: We need 3 bouquets each has 2 flowers, that means we need 6 flowers. We only have 5 flowers so it is impossible to get the needed bouquets and we return  $-1$ .

Example 3:

Input:  $\text{bloomDay} = [7,7,7,7,12,7,7]$ ,  $m = 2$ ,  $k = 3$

Output: 12

Explanation: We need 2 bouquets each should have 3 flowers.

Here's the garden after the 7 and 12 days:

After day 7:  $[x, x, x, x, _, x, x]$

We can make one bouquet of the first three flowers that bloomed. We cannot make another bouquet from the last three flowers that bloomed because they are not adjacent.

After day 12:  $[x, x, x, x, x, x, x]$

It is obvious that we can make two bouquets in different ways.

Example 4:

Input:  $\text{bloomDay} = [1000000000,1000000000]$ ,  $m = 1$ ,  $k = 1$

Output: 1000000000

Explanation: You need to wait 1000000000 days to have a flower ready for a bouquet.

Example 5:

Input:  $\text{bloomDay} = [1,10,2,9,3,8,4,7,5,6]$ ,  $m = 4$ ,  $k = 2$

Output: 9

Constraints:

```
bloomDay.length == n
1 <= n <= 10^5
1 <= bloomDay[i] <= 10^9
1 <= m <= 10^6
1 <= k <= n
```

*Solution*

06/13/2020:

```
class Solution {
public:
    int minDays(vector<int>& bloomDay, int m, int k) {
        if (bloomDay.size() < (long long)m * k) return -1;
        int lo = 0, hi = 1e9;
        while (lo < hi) {
            int mid = lo + (hi - lo) / 2;
            int amount = 0, candidate = 0;
            for (auto& b : bloomDay) {
                if (b <= mid)
                    ++candidate;
                else
                    candidate = 0;
                if (candidate == k) {
                    candidate = 0;
                    ++amount;
                }
            }
            if (amount >= m)
                hi = mid;
            else
                lo = mid + 1;
        }
        return lo;
    }
};
```

## 1483. Kth Ancestor of a Tree Node

*Description*

You are given a tree with  $n$  nodes numbered from  $0$  to  $n-1$  in the form of a parent array where  $\text{parent}[i]$  is the parent of node  $i$ . The root of the tree is node  $0$ .

Implement the function `getKthAncestor(int node, int k)` to return the  $k$ -th ancestor of the given node. If there is no such ancestor, return `-1`.

The  $k$ -th ancestor of a tree node is the  $k$ -th node in the path from that node to the root.

Example:

Input:

```
["TreeAncestor","getKthAncestor","getKthAncestor","getKthAncestor"]
[[7,[-1,0,0,1,1,2,2]],[3,1],[5,2],[6,3]]
```

Output:

```
[null,1,0,-1]
```

Explanation:

```
TreeAncestor treeAncestor = new TreeAncestor(7, [-1, 0, 0, 1, 1, 2, 2]);
```

```
treeAncestor.getKthAncestor(3, 1); // returns 1 which is the parent of 3
```

```
treeAncestor.getKthAncestor(5, 2); // returns 0 which is the grandparent of 5
```

```
treeAncestor.getKthAncestor(6, 3); // returns -1 because there is no such ancestor
```

Constraints:

```
1 <= k <= n <= 5*10^4
```

```
parent[0] == -1 indicating that 0 is the root node.
```

```
0 <= parent[i] < n for all 0 < i < n
```

```
0 <= node < n
```

```
There will be at most 5*10^4 queries.
```

*Solution*

06/13/2020:

```
const int MAXD = 16;
int dp[50005][MAXD];

class TreeAncestor {
public:
    TreeAncestor(int n, vector<int>& parent) {
        for (int i = 0; i < n; ++i) dp[i][0] = parent[i];
```

```

for (int d = 1; d < MAXD; ++d) {
    for (int i = 0; i < n; ++i) {
        if (dp[i][d - 1] == -1)
            dp[i][d] = -1;
        else
            dp[i][d] = dp[dp[i][d - 1]][d - 1];
    }
}

int getKthAncestor(int node, int k) {
    for (int d = MAXD - 1; d >= 0; --d) {
        if ((1 << d) <= k) {
            k -= 1 << d;
            node = dp[node][d];
            if (node == -1) return node;
        }
    }
    return node;
}
};

/**
 * Your TreeAncestor object will be instantiated and called as such:
 * TreeAncestor* obj = new TreeAncestor(n, parent);
 * int param_1 = obj->getKthAncestor(node,k);
 */

```

## 1488. Avoid Flood in The City

### Description

Your country has an infinite number of lakes. Initially, all the lakes are empty, but when it rains over the  $n$ th lake, the  $n$ th lake becomes full of water. If it rains over a lake which is full of water, there will be a flood. Your goal is to avoid the flood in any lake.

Given an integer array rains where:

rains[i] > 0 means there will be rains over the rains[i] lake.

rains[i] == 0 means there are no rains this day and you can choose one lake this day and dry it.

Return an array ans where:

ans.length == rains.length

ans[i] == -1 if rains[i] > 0.



ans[i] is the lake you choose to dry in the ith day if rains[i] == 0.  
If there are multiple valid answers return any of them. If it is impossible to avoid flood return an empty array.

Notice that if you chose to dry a full lake, it becomes empty, but if you chose to dry an empty lake, nothing changes. (see example 4)

Example 1:

Input: rains = [1,2,3,4]

Output: [-1,-1,-1,-1]

Explanation: After the first day full lakes are [1]

After the second day full lakes are [1,2]

After the third day full lakes are [1,2,3]

After the fourth day full lakes are [1,2,3,4]

There's no day to dry any lake and there is no flood in any lake.

Example 2:

Input: rains = [1,2,0,0,2,1]

Output: [-1,-1,2,1,-1,-1]

Explanation: After the first day full lakes are [1]

After the second day full lakes are [1,2]

After the third day, we dry lake 2. Full lakes are [1]

After the fourth day, we dry lake 1. There is no full lakes.

After the fifth day, full lakes are [2].

After the sixth day, full lakes are [1,2].

It is easy that this scenario is flood-free. [-1,-1,1,2,-1,-1] is another acceptable scenario.

Example 3:

Input: rains = [1,2,0,1,2]

Output: []

Explanation: After the second day, full lakes are [1,2]. We have to dry one lake in the third day.

After that, it will rain over lakes [1,2]. It's easy to prove that no matter which lake you choose to dry in the 3rd day, the other one will flood.

Example 4:

Input: rains = [69,0,0,0,69]

Output: [-1,69,1,1,-1]

Explanation: Any solution on one of the forms [-1,69,x,y,-1], [-1,x,69,y,-1] or [-1,x,y,69,-1] is acceptable where  $1 \leq x,y \leq 10^9$

Example 5:

Input: rains = [10,20,20]

Output: []

Explanation: It will rain over lake 20 two consecutive days. There is no chance to dry any lake.

Constraints:

```
1 <= rains.length <= 10^5  
0 <= rains[i] <= 10^9
```

*Solution*

06/20/2020:

```
class Solution {  
public:  
    vector<int> avoidFlood(vector<int>& rains) {  
        int n = rains.size();  
        unordered_map<int, int> bad;  
        set<int> canfix;  
        vector<int> ret(n, -1);  
        for (int i = 0; i < n; ++i) {  
            if (rains[i] == 0) {  
                canfix.insert(i);  
                ret[i] = 1;  
            } else {  
                if (bad.count(rains[i]) == 0) {  
                    bad[rains[i]] = i;  
                } else {  
                    int must = bad[rains[i]];  
                    auto it = canfix.upper_bound(must);  
                    if (it == canfix.end()) return {};  
                    ret[*it] = rains[i];  
                    canfix.erase(it);  
                    bad[rains[i]] = i;  
                }  
            }  
        }  
        return ret;  
    }  
};
```